## Deductive Program Verification: Solutions to Exercise #3

Yu-Fang Chen

Dept. of Information Management, National Taiwan University

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## Note

We assume the binding powers of the various operators decrease in this order:  $(\cdot)^n$  (exponentiation),  $\{+, -\}$ ,  $\neg$ ,  $\{=, \geq, \leq\}$ ,  $\{\forall, \exists\}, \{\land, \lor\}, \rightarrow, \leftrightarrow, \equiv$ .

## Solution

1. Consider the following program skeleton of mutual exclusion by a semaphore.

Program MUX-SEM:

s: natural **initially** s = 1

 $\left[\begin{array}{c}l_0: \textbf{loop forever do}\\ \left[\begin{array}{c}l_1: \text{ request}(s);\\ l_2: \text{ release}(s);\end{array}\right]\right] \parallel \left[\begin{array}{c}m_0: \textbf{loop forever do}\\ \left[\begin{array}{c}m_1: \text{ request}(s);\\ m_2: \text{ release}(s);\end{array}\right]\right]$ 

where

- request(s) <sup>△</sup>= await s > 0 then s := s 1 end
  release(s) <sup>△</sup>= s := s + 1
- Please re-describe the program as a fair transition system (FTS) and specify its safety and response properties in LTL.

Solution.

•  $V \triangleq \{s: natural, \pi_0 : \{l_0, l_1, l_2\}, \pi_1 : \{m_0, m_1, m_2\}\}$ 

• 
$$\Theta \triangleq \pi_0 = l_0 \land \pi_1 = m_0 \land s = 1$$

•  $T \triangleq \{\tau_I, \tau_{l_0}, \tau_{l_1}, \tau_{l_2}, \tau_{m_0}, \tau_{m_1}, \tau_{m_2}\}$ , whose transition relations are  $\rho_I : \pi'_0 = \pi_0 \land \pi'_1 = \pi_1 \land s' = s$   $\rho_{l_0} : \pi_0 = l_0 \land \pi'_0 = l_1 \land s' = s \land \pi'_1 = \pi_1$   $\rho_{l_1} : \pi_0 = l_1 \land s > 0 \land \pi'_0 = l_2 \land s' = s - 1 \land \pi'_1 = \pi_1$   $\rho_{l_2} : \pi_0 = l_2 \land s' = s + 1 \land \pi'_1 = \pi_1$   $\rho_{m_0} : \pi_1 = m_0 \land \pi'_1 = m_1 \land s' = s \land \pi'_0 = \pi_0$  $\rho_{m_1} : \pi_1 = m_1 \land s > 0 \land \pi'_1 = m_2 \land s' = s - 1 \land \pi'_0 = \pi_0$   $\rho_{m_2}: \pi_1 = m_2 \wedge s' = s + 1 \wedge \pi'_0 = \pi_0$ 

- $\mathcal{J} = \{\tau_{l_0}, \tau_{l_2}, \tau_{m_0}, \tau_{m_2}\}$
- $\mathcal{C} = \{\tau_{l_1}, \tau_{m_1}\}$

The safety property satisfied by this model:  $\Box(\neg(\pi_0 = l_2 \land \pi_1 = m_2))$ 

The response property satisfied by this model:

 $\Box(\pi_0 = l_1 \to \Diamond(\pi_0 = l_2)) \land \Box(\pi_1 = m_1 \to \Diamond(\pi_1 = m_2))$