Automata Theory

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Outline

- Finite state automata
- Regular Expressions
- WS1S
- Transducers
- ω-Automata
- Linear temporal logic

Finite State Automata

- *Finite state automata* (*FSA*)
  - A self-operating machine with predetermined operations and a limited amount of memory
  - A language recognizer
  - A simplest computational model (abstract model of computers)

Applications

- Lexical analyzer
- String processing
- Spell checking
- Model checking

Finite State Machine

Graphics from pngtree.com

Components of FSA

- Inputs (words over a finite alphabet)
- States
- Starting states
- Transitions
- Final states

Alphabet

- An *alphabet* is a set of symbols
- Types of alphabet: *classical* and *propositional*
- Examples:
  - $\{a,\ b\}$
  - $\{send,\ receive,\ ack\}$
  - $\{(p\ \ q),\ (\neg p\ \ q),\ (p\ \neg q),\ (\neg p\ \neg q)\}$

Words

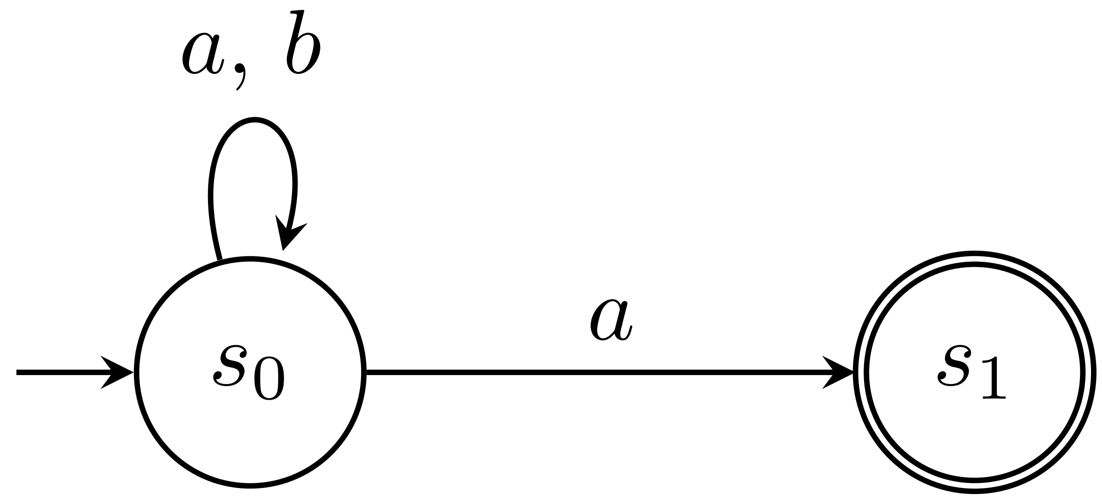
- Let Σ be a finite alphabet
- A *word* w over Σ ($w \in \Sigma^*$) is a sequence of symbols $w_0 w_1 w_2 \ldots w_{n-1}$ with $w_i \in \Sigma$
  - Length of w, denoted by $|w|$, is n
  - The empty word is denoted by ϵ
- Examples ($\Sigma = \{a, b\}$):
  - $a\ b\ b\ a$
  - $a\ b\ a\ b\ a\ b$

w^* : repeat w finitely many times

Formal Syntax of FSA

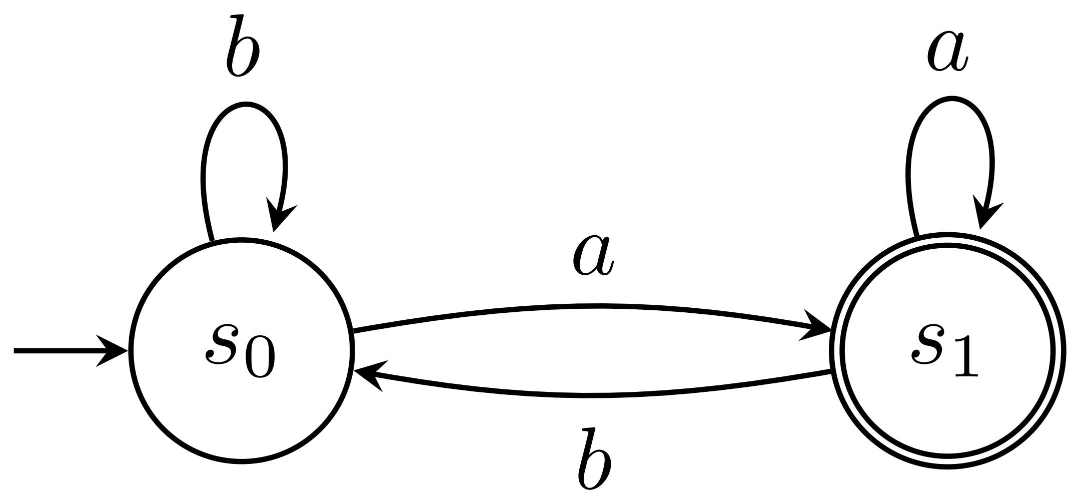
- A finite state automaton is a 5-tuple $(Q, \Sigma, \delta, I, F)$ where
  - Q is a finite set of *states*,
  - Σ is a finite *alphabet*,
  - $\delta : Q \times \Sigma \to 2^Q$ is the *transition function* (sometimes written as a relation $\delta : Q \times \Sigma \times Q$),
  - $I \subseteq Q$ is the set of *initial states*, and
  - $F \subseteq Q$ is the set of *accepting (final) states*

Automaton M_1



- Alphabet: $\{a,\ b\}$
- States: $\{s_0,\ s_1\}$
- Initial states: $\{s_0\}$
- Transitions: $\{(s_0,\ a,\ s_0),\ (s_0,\ a,\ s_1),\ (s_0,\ b,\ s_0)\}$
- Accepting states: $\{s_1\}$

Automaton M_2



$$A = (Q,\ \Sigma,\ \delta,\ I,\ F)$$

$$\Sigma = \{a,\ b\}$$

$Q = ?$

$I = ?$

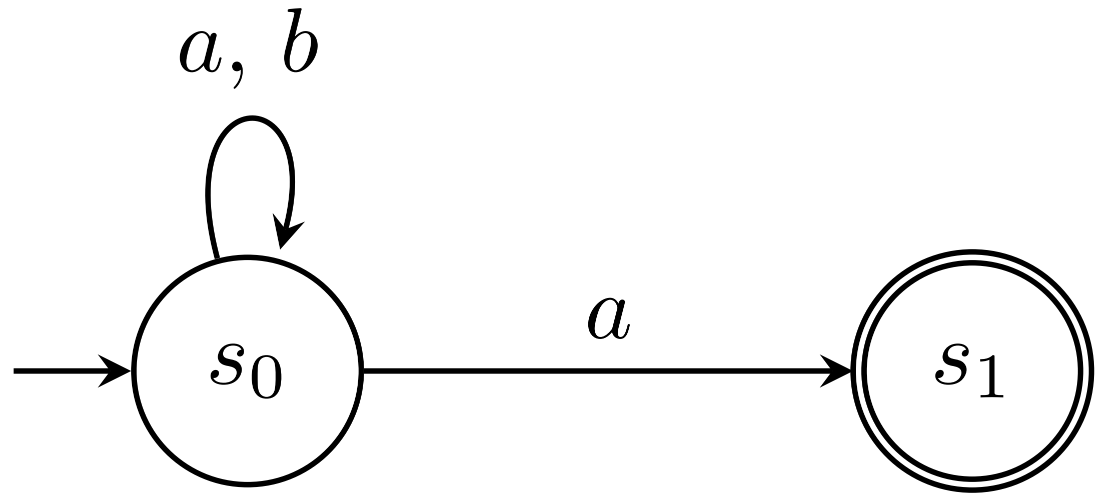
$\delta = ?$

$F = ?$

Formal Semantics of FSA

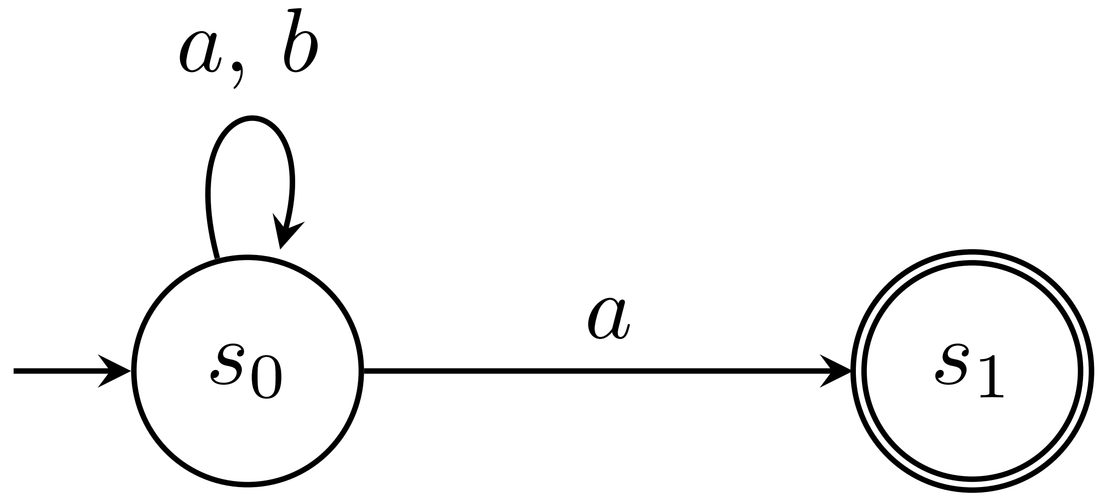
- Let $M = (Q, \Sigma, \delta, I, F)$ be a finite state automaton
- Let $w = w_0 w_1 w_2 \ldots w_{n\text{-}1}$ be a word over Σ
- A *run* of w on M is a sequence of states $s_0 s_1 s_2 \ldots s_n$ where
  - $s_0 \in I$
  - $(s_i, w_i, s_{i+1}) \in \delta$

Runs



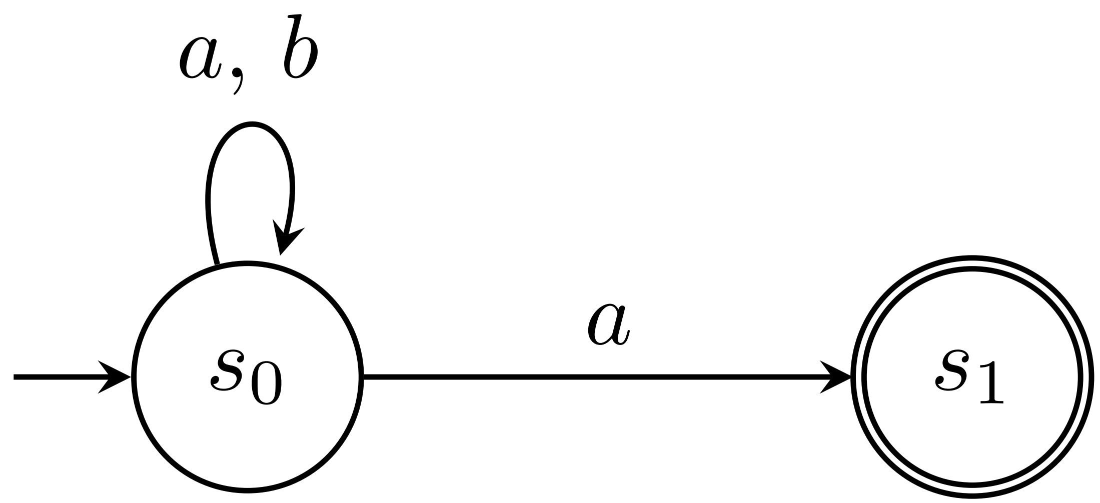
- What are the runs of the following words?
  - $a\ b\ a\ b$
  - $a\ b\ b\ a$

Runs



- What are the runs of the following words?
  - $a\ b\ a\ b$ $\qquad\qquad$ $s_0\ s_0\ s_0\ s_0\ s_0$
  - $a\ b\ b\ a$

Runs

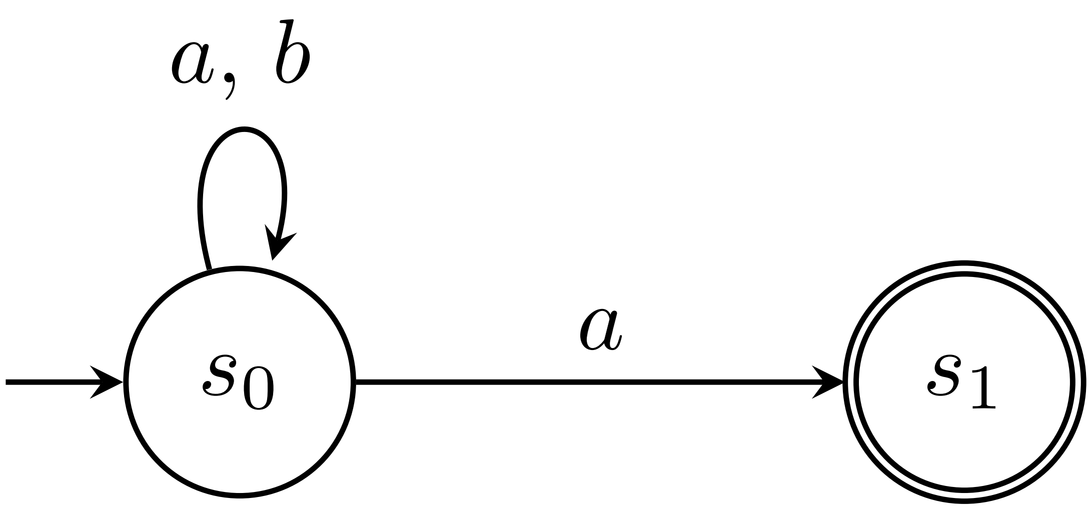


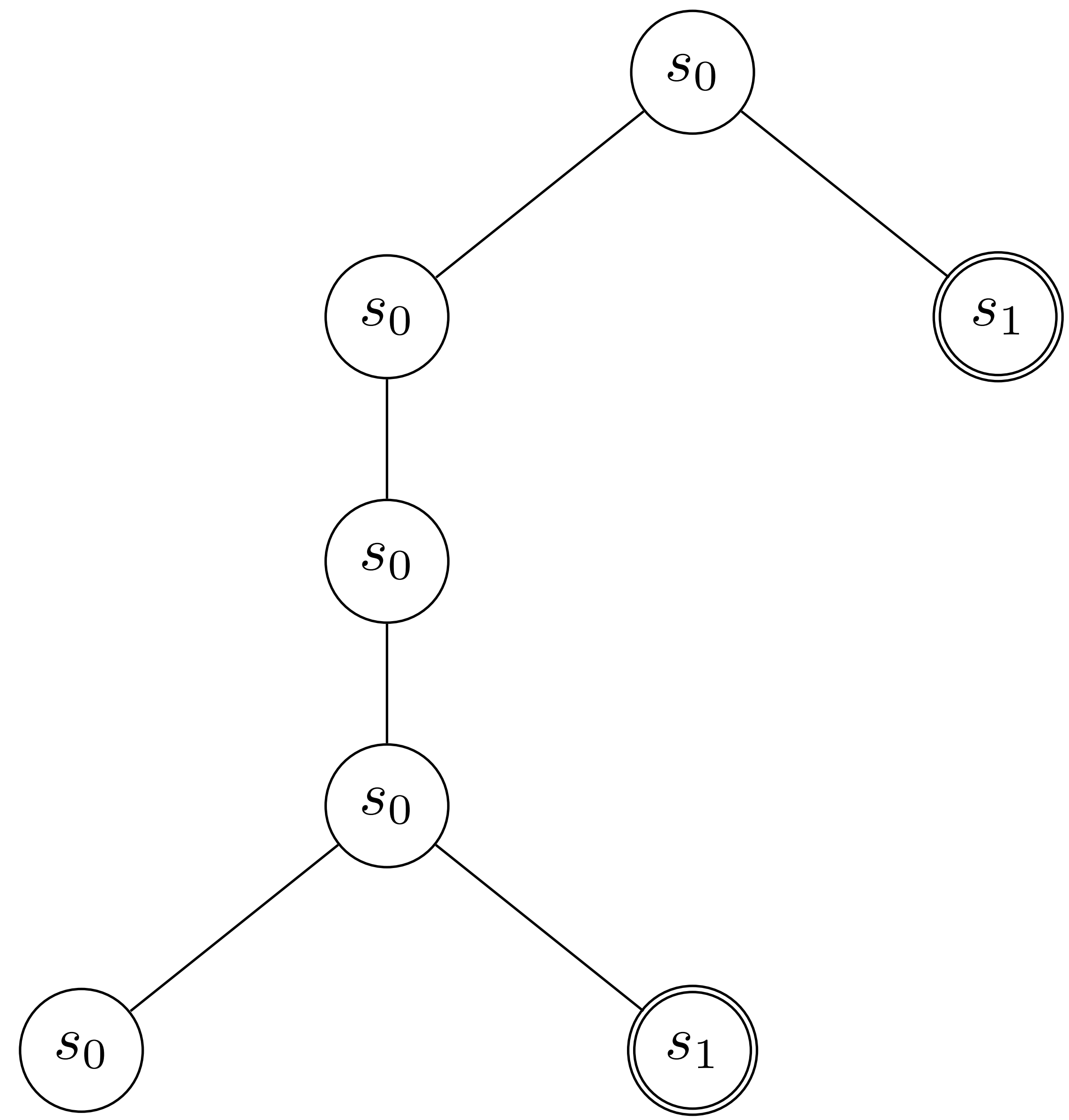
- What are the runs of the following words?

  - $a\ b\ a\ b$ $\qquad$ $s_0\ s_0\ s_0\ s_0\ s_0$

  - $a\ b\ b\ a$ $\qquad$ $s_0\ s_0\ s_0\ s_0\ s_0$ and $s_0\ s_0\ s_0\ s_0\ s_1$

Run Tree



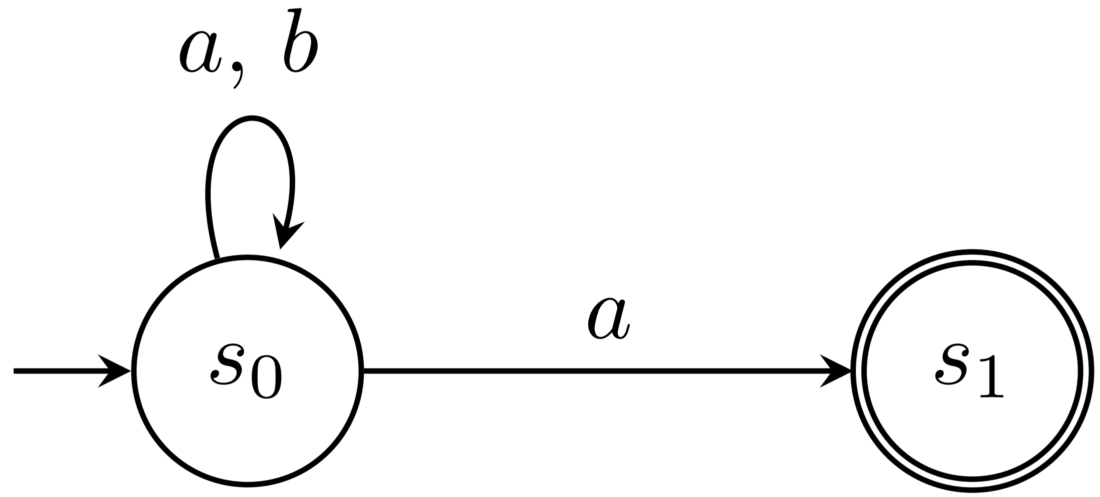


the run tree of $abba$ on M_1

Formal Semantics of FSA (cont'd)

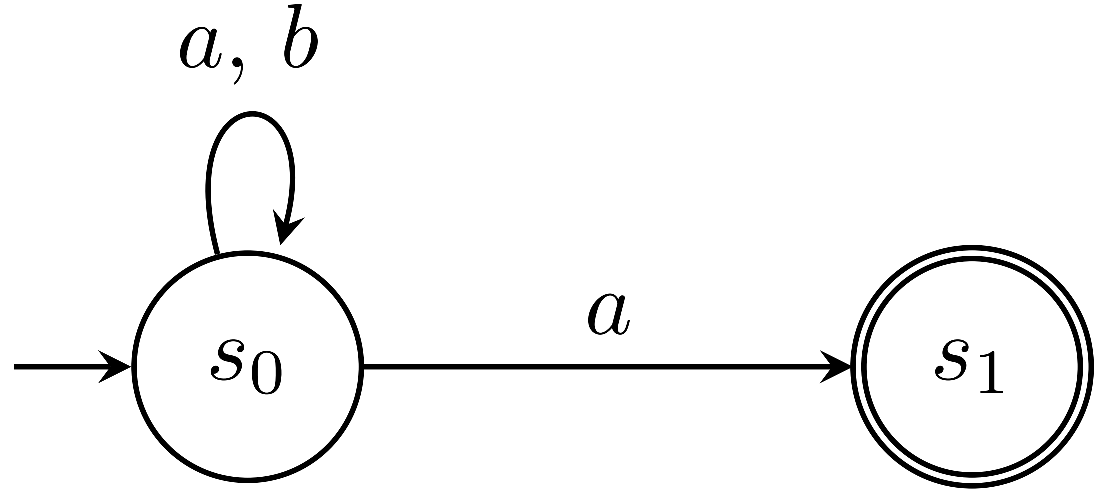
- $M = (Q, \Sigma, \delta, I, F)$
- A run $s_0 s_1 s_2 \ldots s_n$ is *accepting* if $s_n \in F$
- A word w is accepted by M if there is an accepting run of w on M
- The *language* of M is the set of strings accepted by M, denoted by $L(M)$

Accepting Runs



- Which run is accepting?
  - $s_0\ s_0\ s_0\ s_0\ s_0$
  - $s_0\ s_0\ s_0\ s_0\ s_1$

Automaton M_1



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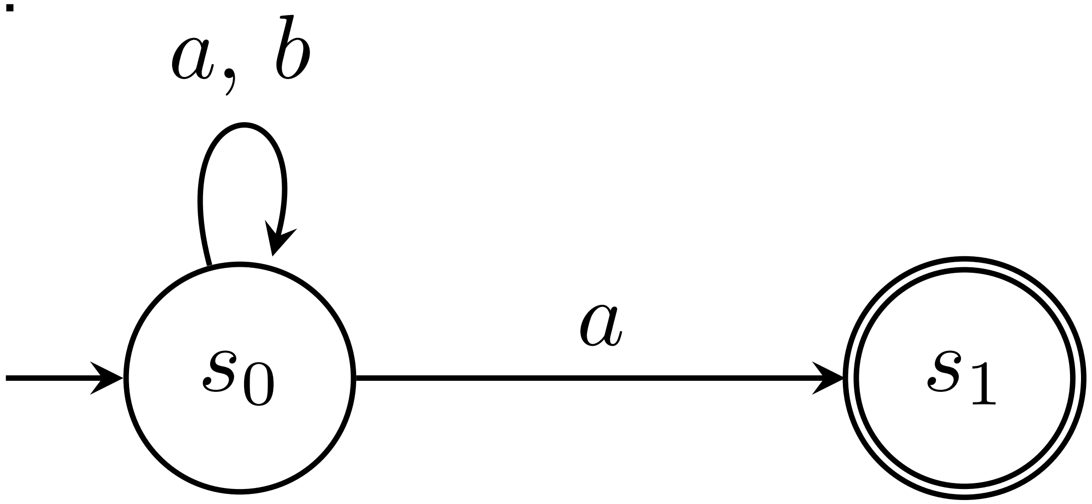
Languages

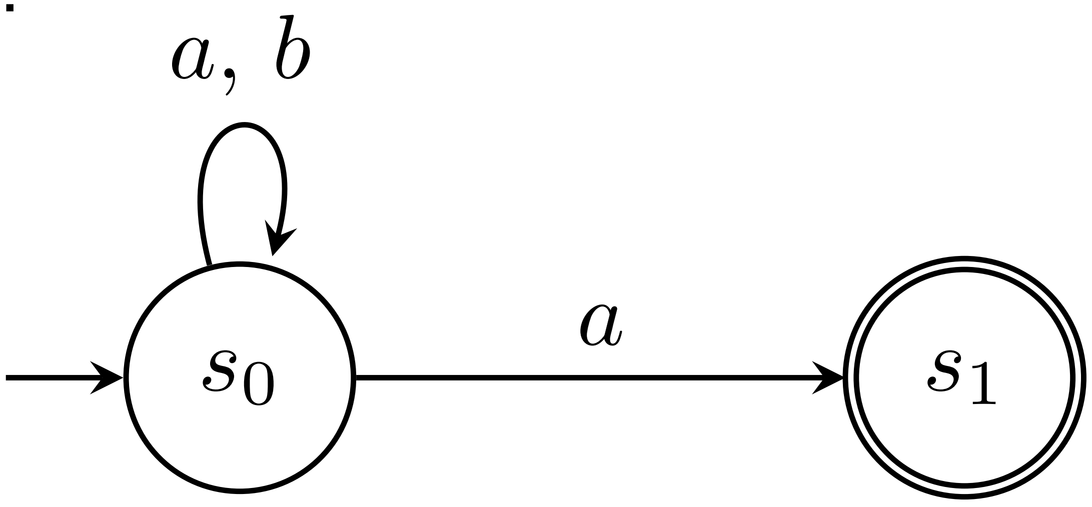
- What is the language of M_1?



- The language recognized by a finite state automaton is a *regular language*

Languages

- What is the language of M_1?

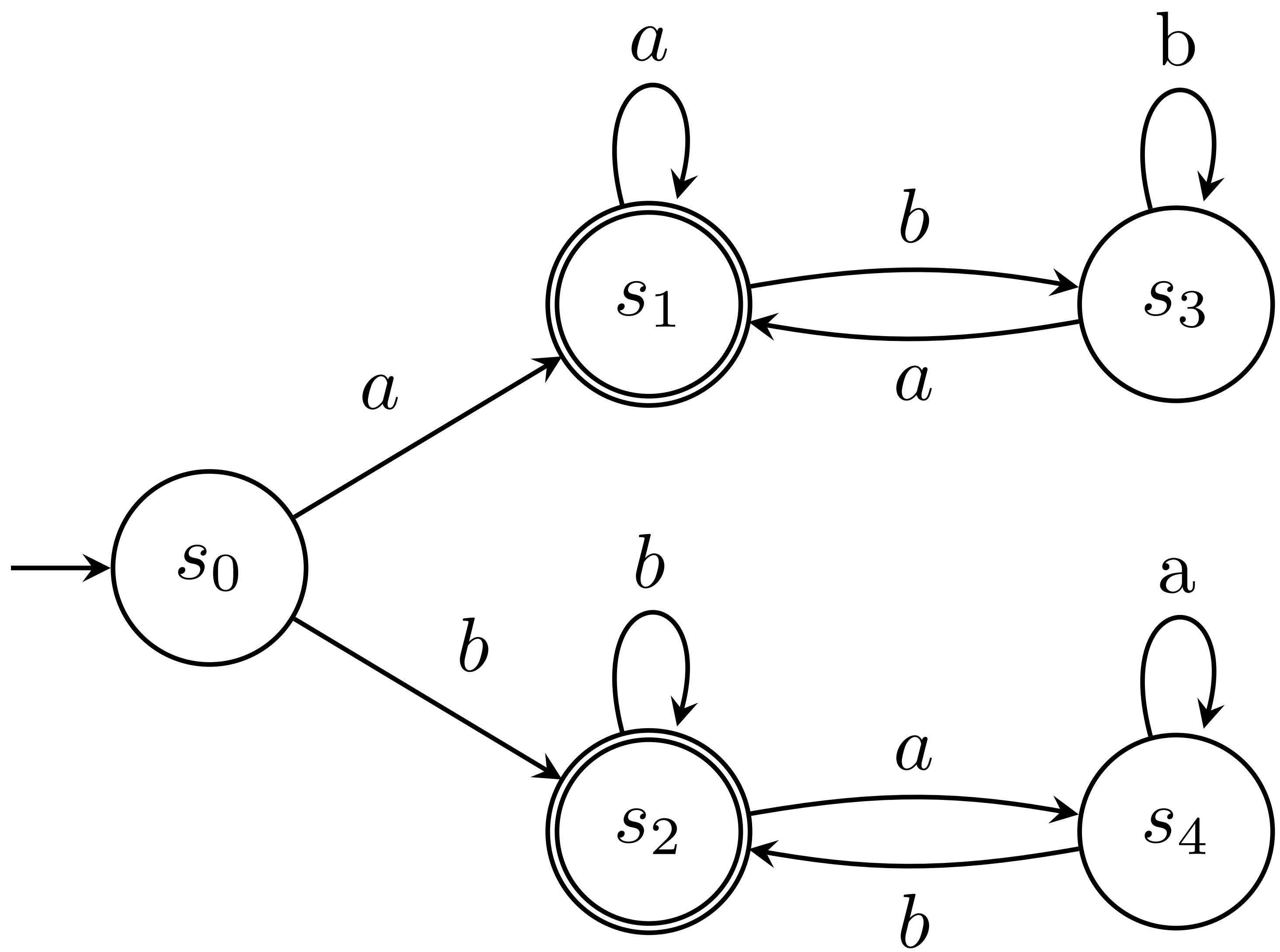


$L(M_1) = \{ w_0 w_1 \ldots w_n \mid n \in \mathbb{N} \textit{ and } w_n = a \}$

- The language recognized by a finite state automaton is a *regular language*

Languages (cont'd)

- What is the language of the following automaton?



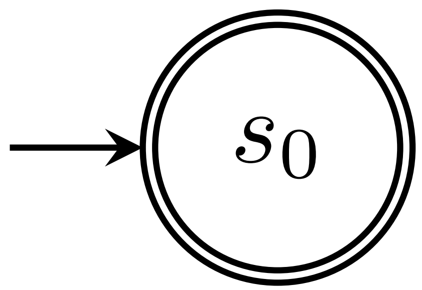
Exercise

- Given an alphabet $\{1, 2, +\}$, draw a finite state automaton such that the automaton accepts words evaluated to 3
- Given an alphabet $\{0, 1,\}$, draw a finite state automaton such that the automaton accepts words containing substring 001

Emptiness and Universality

- $M = (Q, \Sigma, \delta, I, F)$
- An automaton M is *empty* if $L(M) = \varnothing$
- An automaton M is *universal* if $L(M) = \Sigma^*$

Emptiness and Universality



is this automaton empty?

- $M = (Q, \Sigma, \delta, I, F)$
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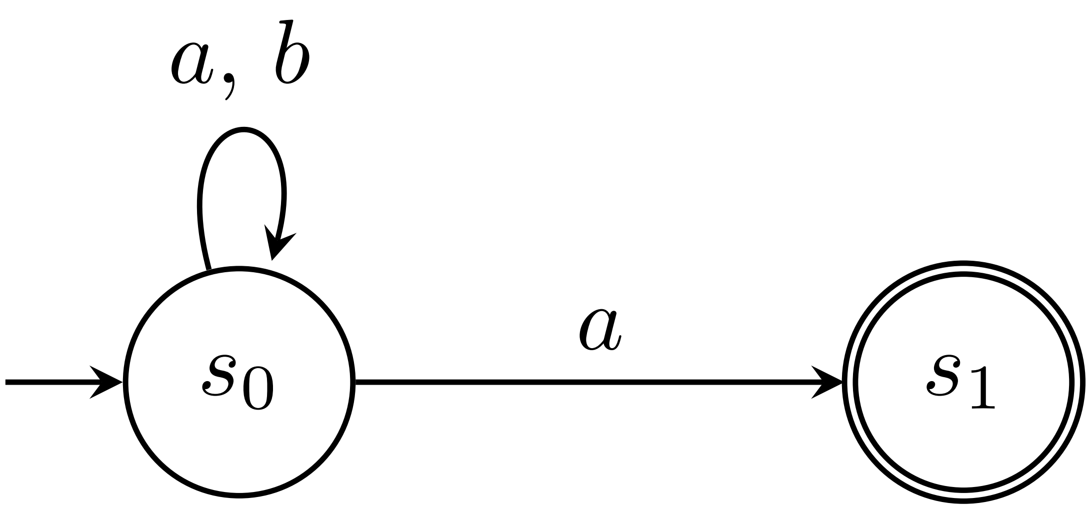
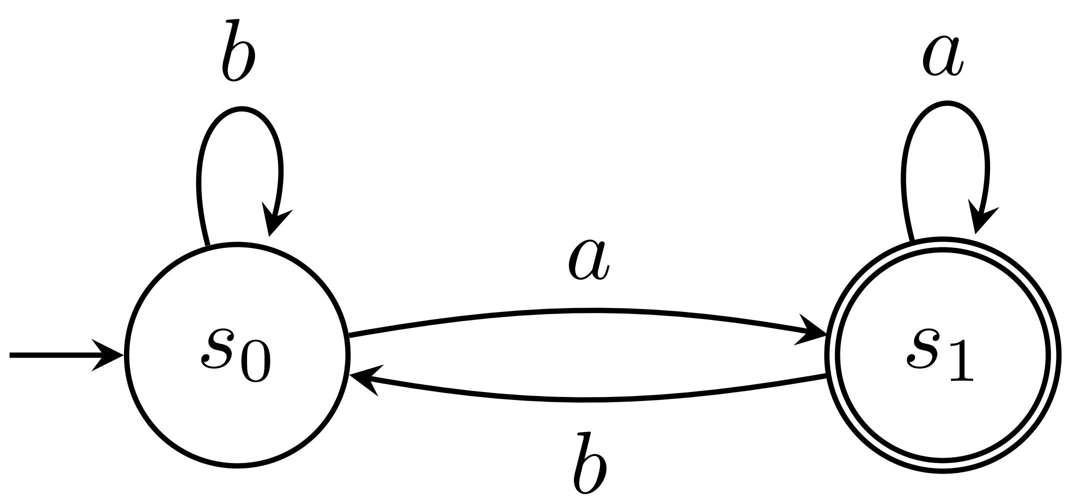
Equivalence

- Two automata are *equivalent* if they recognize the same language





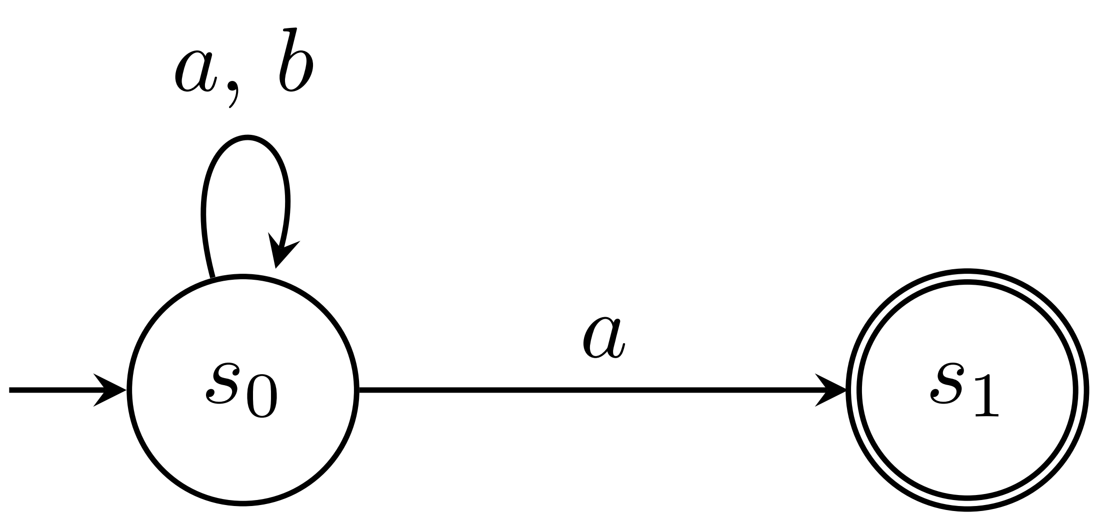
$$L(M_1) = L(M_2)?$$

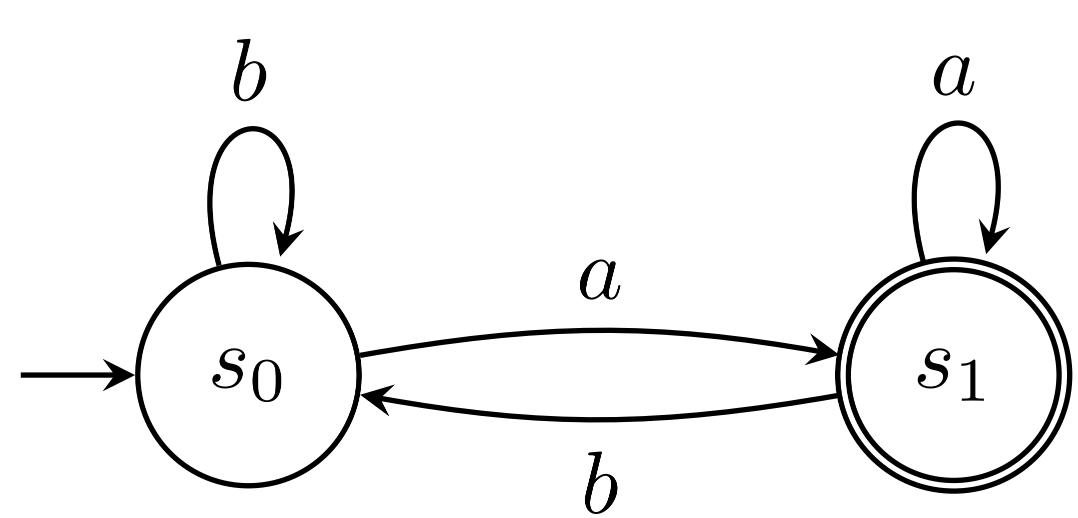
Deterministic Finite Automata (DFA)

- An automaton $M = (Q, \Sigma, \delta, I, F)$ is *deterministic* if
  - $|I| = 1$ and
  - $|\delta(s, a)| = 1$ for all $s \in Q$ and $a \in \Sigma$

(is complete if $|\delta(s, a)| \geq 1$)

- Which one is deterministic?





Determinism VS Nondeterminism

- The language $L(D)$ of a DFA D is accepted by the NFA D (A DFA is also an NFA)

- Given an NFA N, Can we construct a DFA D such that $L(D) = L(N)$?

Determinism VS Nondeterminism

- The language $L(D)$ of a DFA D is accepted by the NFA D (A DFA is also an NFA)

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Determinism VS Nondeterminism

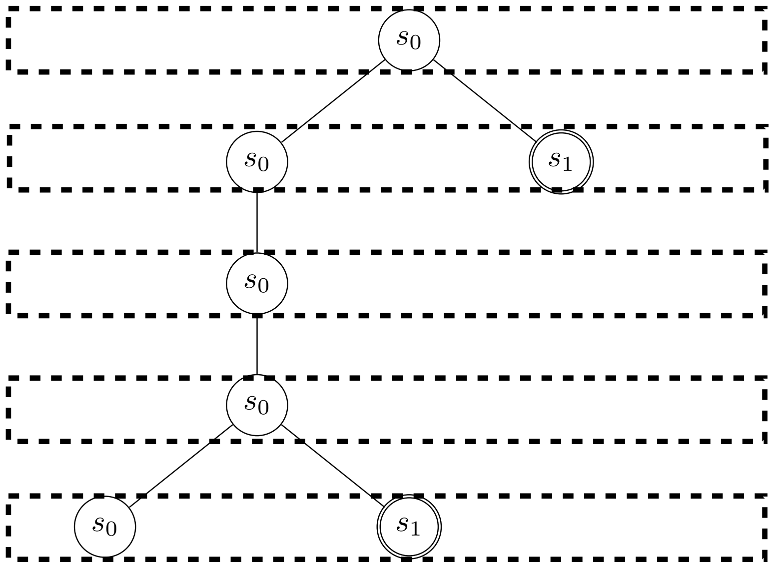
- The language $L(D)$ of a DFA D is accepted by the NFA D (A DFA is also an NFA)

- Given an NFA N, Can we construct a DFA D such that $L(D) = L(N)$?

- DFA and NFA have the same expressive power

Determinization

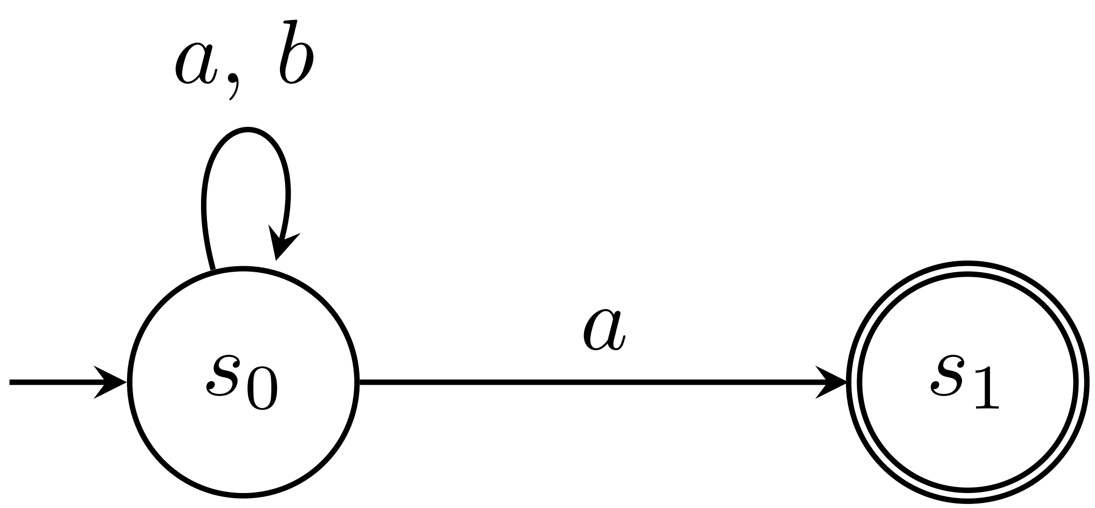
- Let $N = (Q, \Sigma, \delta, I, F)$.
- By *subset construction*, define $D = (2^Q, \Sigma, \Delta, \{ I \}, G)$ where
  - $\Delta(S, a) = \cup_{s \in S} \delta(s, a)$, and
  - $G = \{ S \in 2^Q \mid S \cap F \neq \varnothing \}$

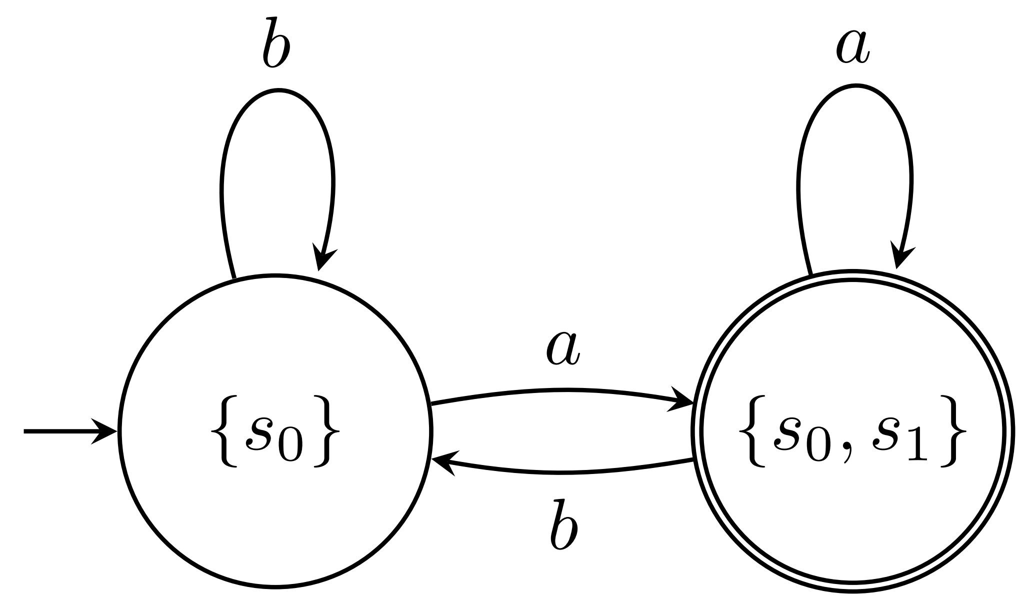


- We can show that $L(N) = L(D)$ by induction on the length of input words

Subset Construction

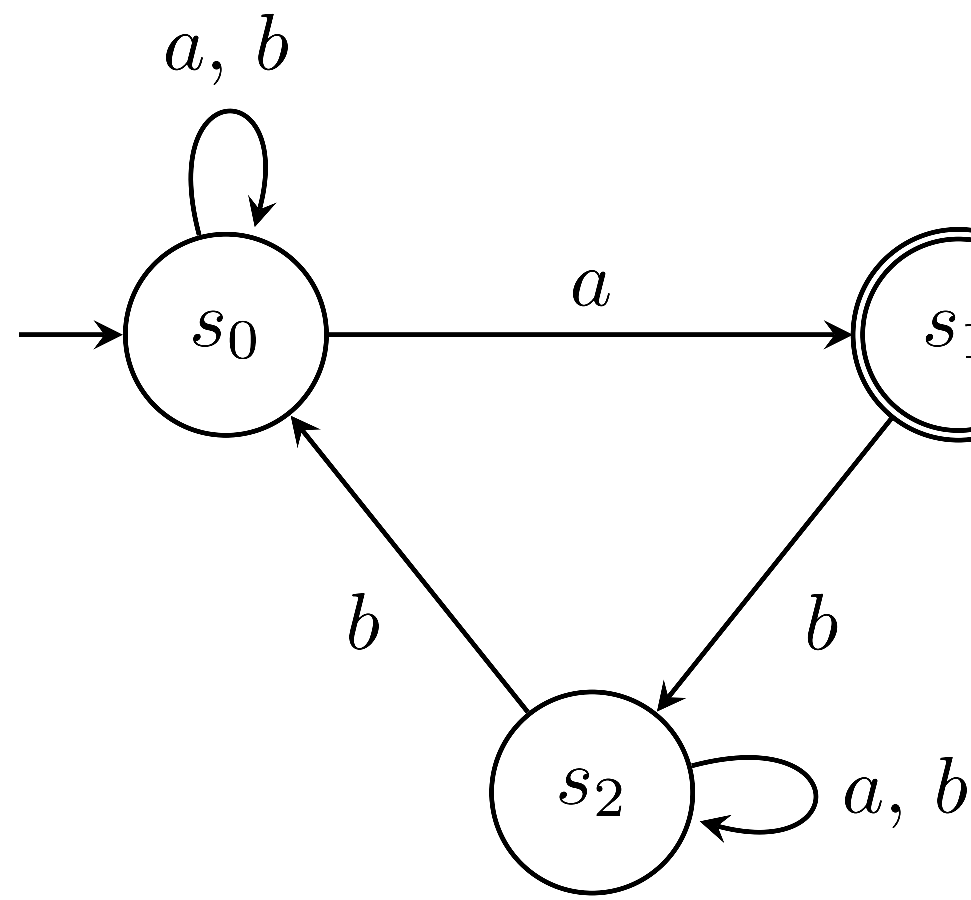
- What is the determinization of M_1?





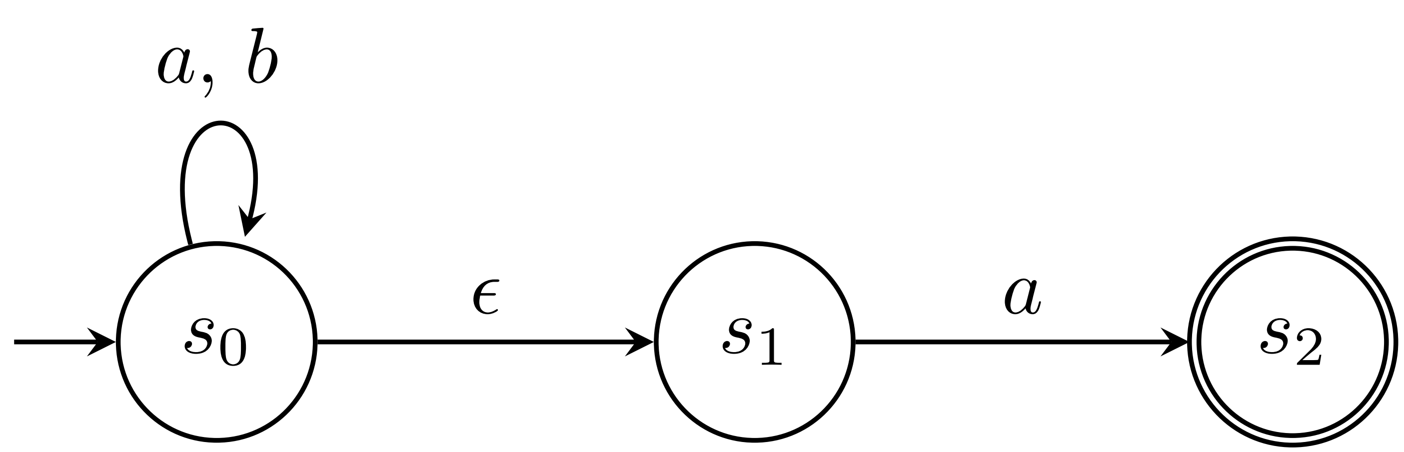
Exercise

- Apply subset construction to determinize the following automaton



ϵ-Transitions

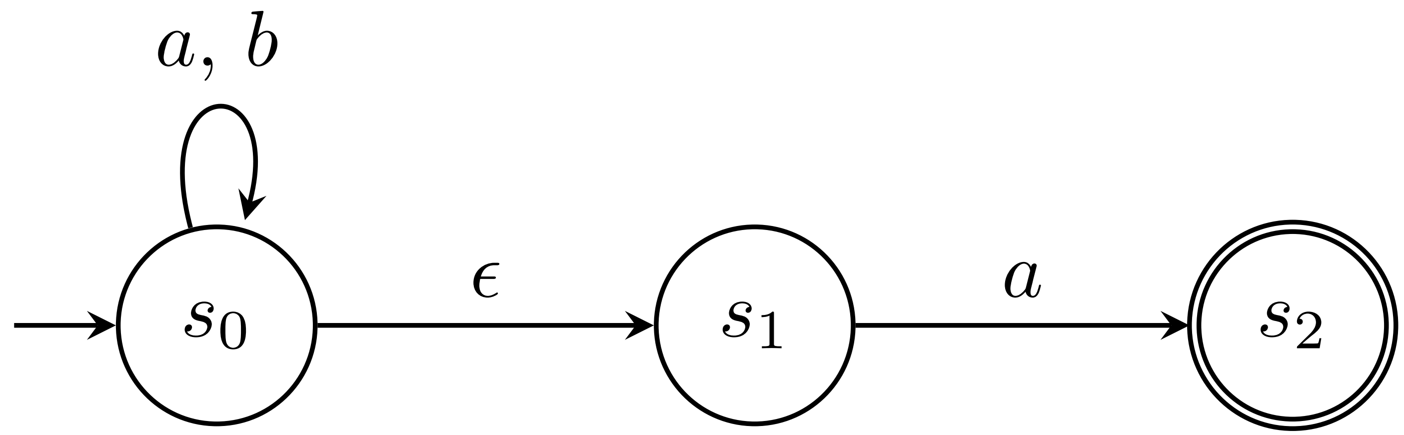
- Assume ϵ does not belong to the alphabet
- An *ϵ-transition* is a transition that does not need to consume any symbol
- ϵ-transitions are only allowed in NFA
- DFA and NFA with ϵ-transitions have the same expressive power



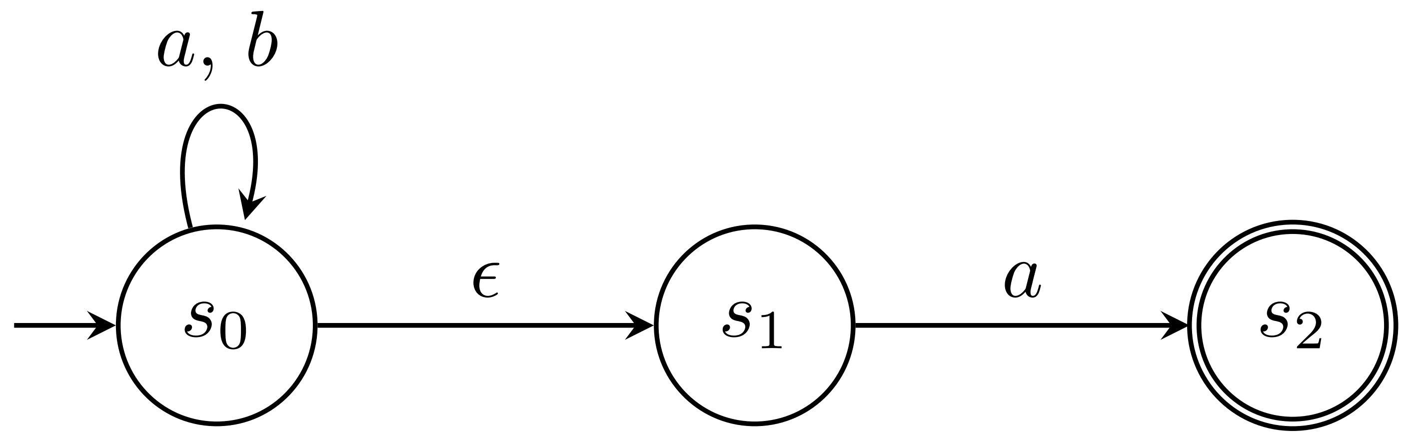
Elimination of ϵ-Transitions

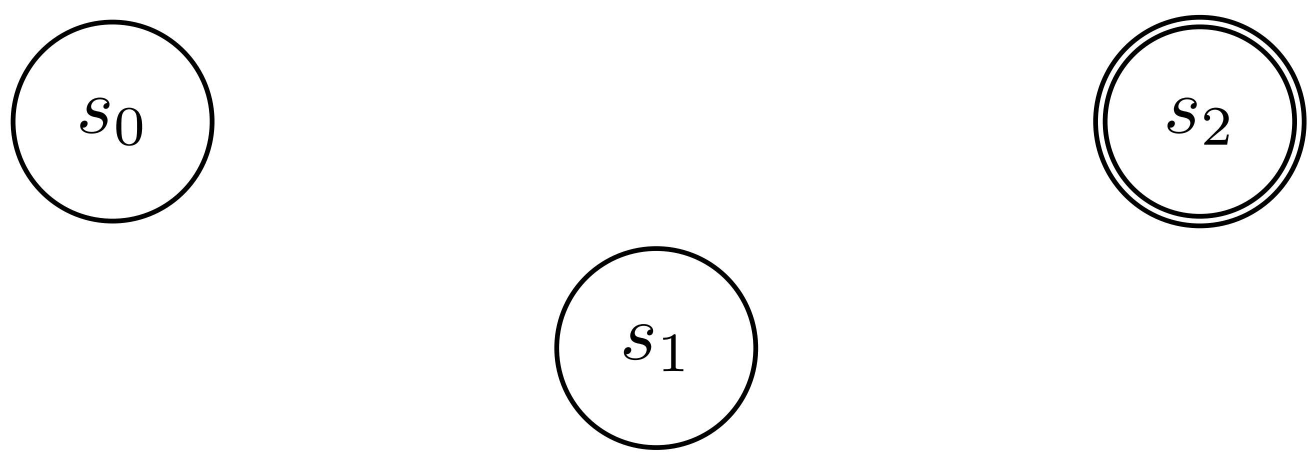
- $M = (Q, \Sigma \cup \{\epsilon\}, \delta, I, F)$ is an NFA with ϵ-transitions
- Let $E(X)$ denote the ϵ-closure of $X \subseteq Q$
  - $E(X) = \{ s \mid s \in X$ *or* s *is reachable from a state in* X *through* ϵ-*transitions* $\}$
- Construct an NFA $N = (Q, \Sigma, \Delta, J, F)$ *where*
  - $\Delta(s, a) = E(\delta(s, a))$, and
  - $J = E(I)$

Elimination of ϵ-Transitions Example

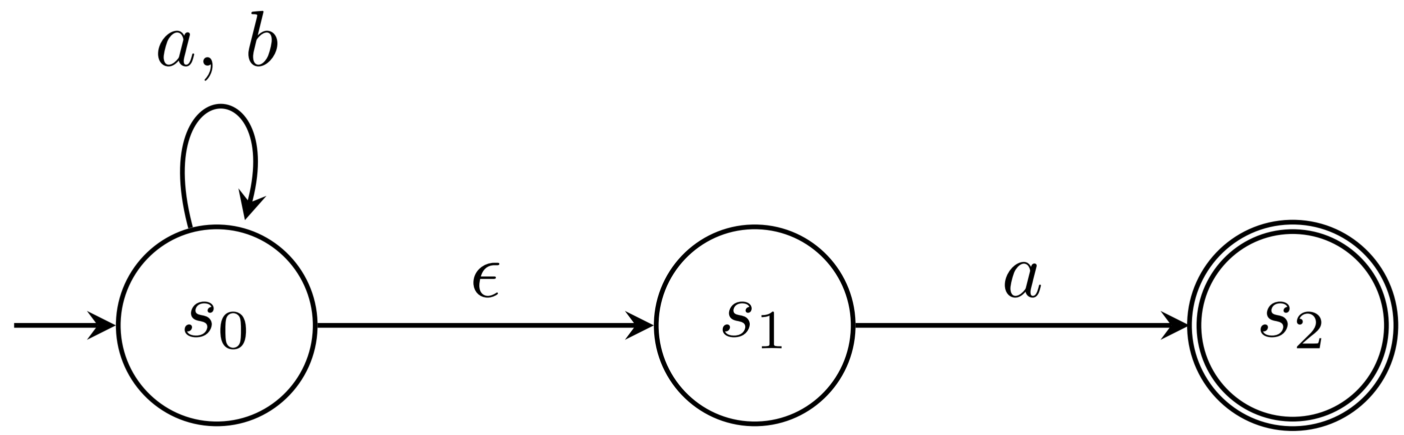


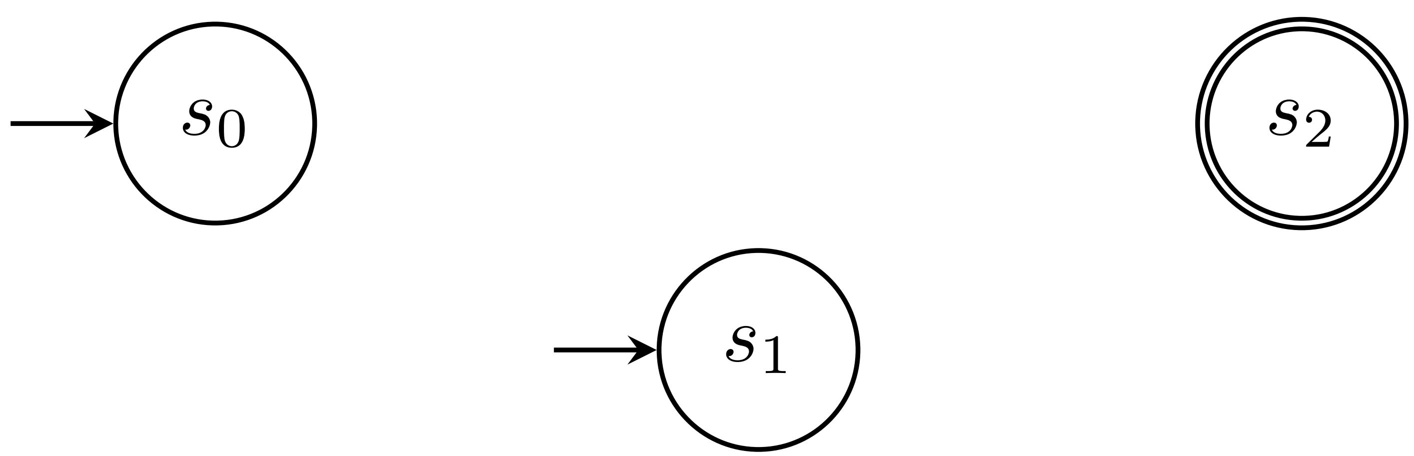
Elimination of ϵ-Transitions
Example





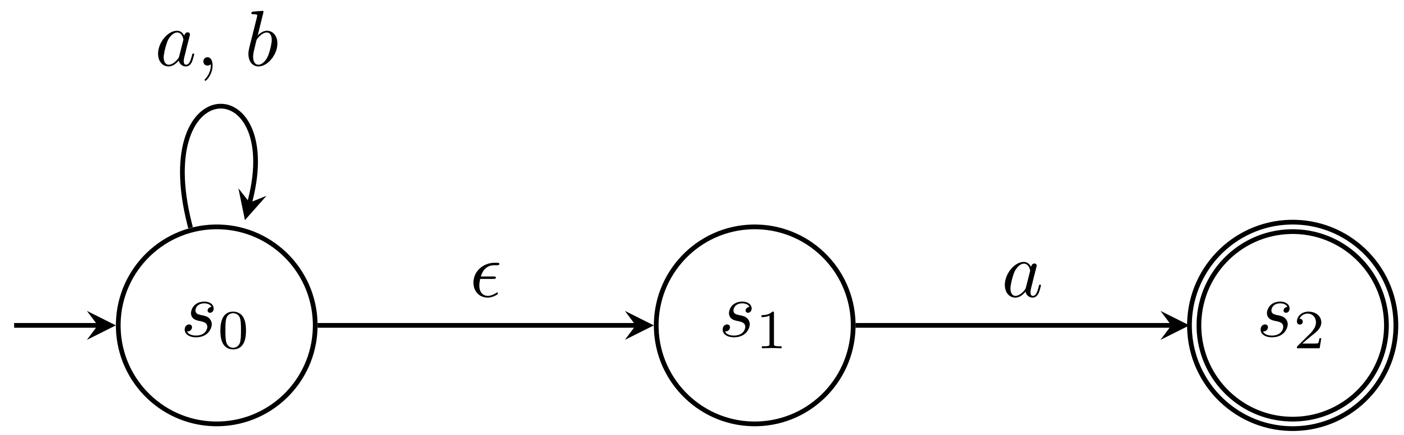
Elimination of ϵ-Transitions Example

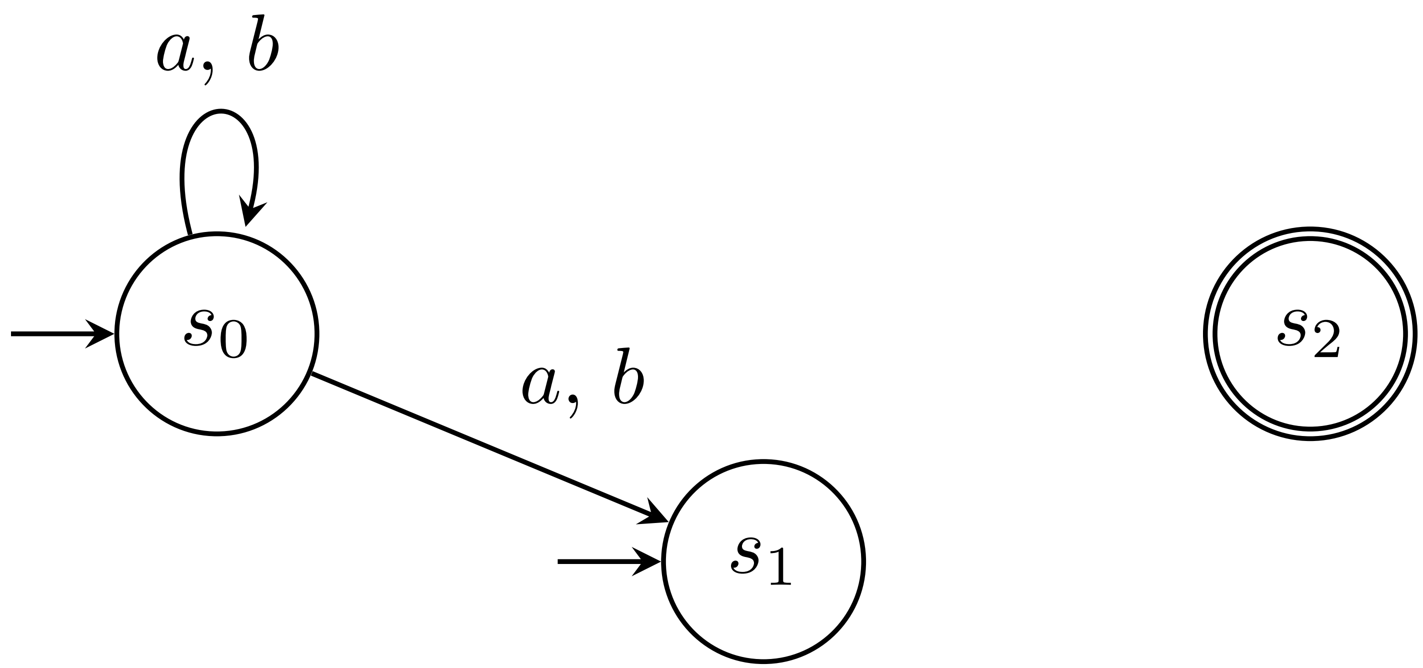




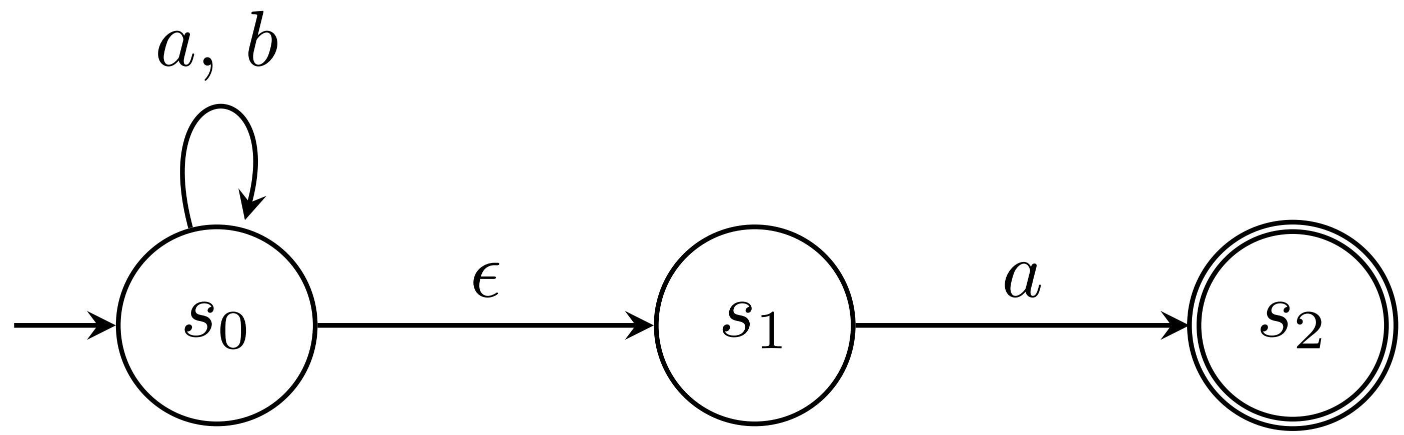
Elimination of ϵ-Transitions

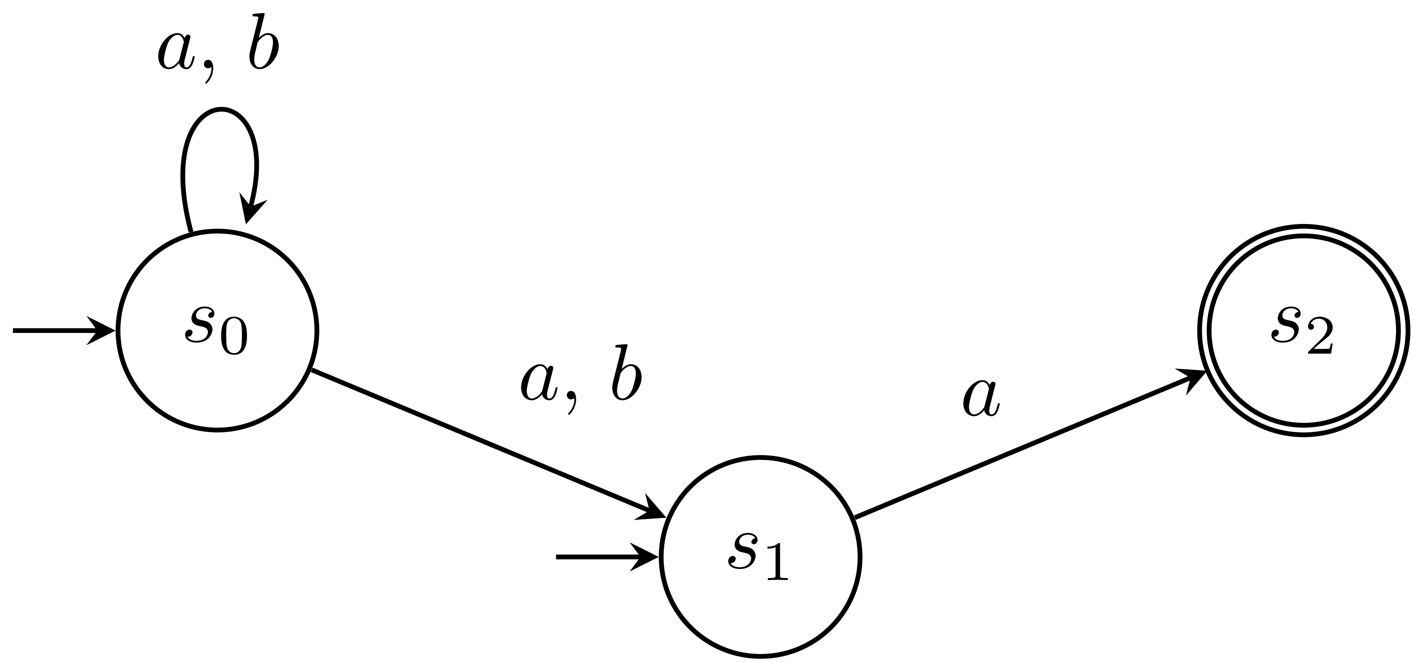
Example





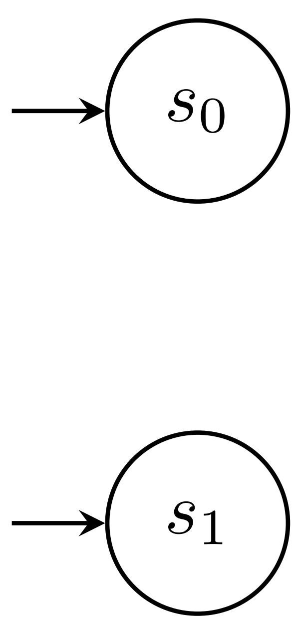
Elimination of ϵ-Transitions Example

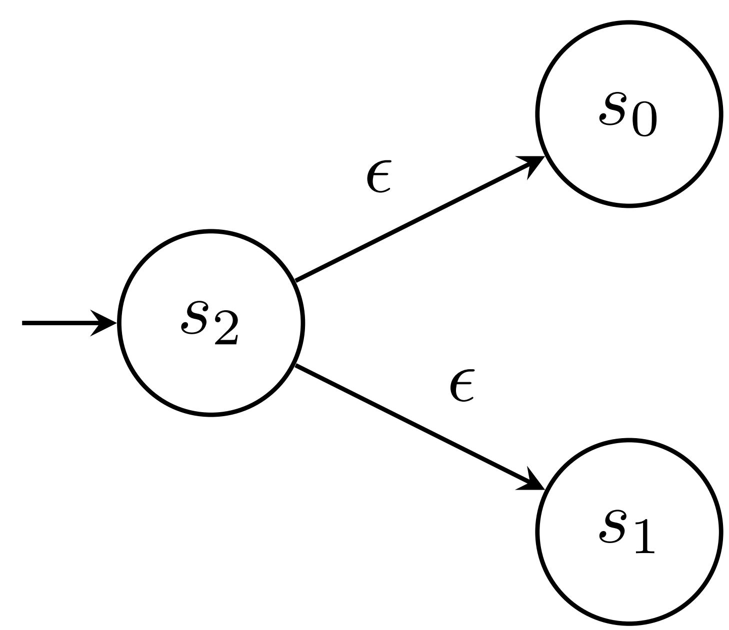




Single Initial State

- NFA may be defined as automata with single initial state
- NFA with multiple initial states does not have more expressive power





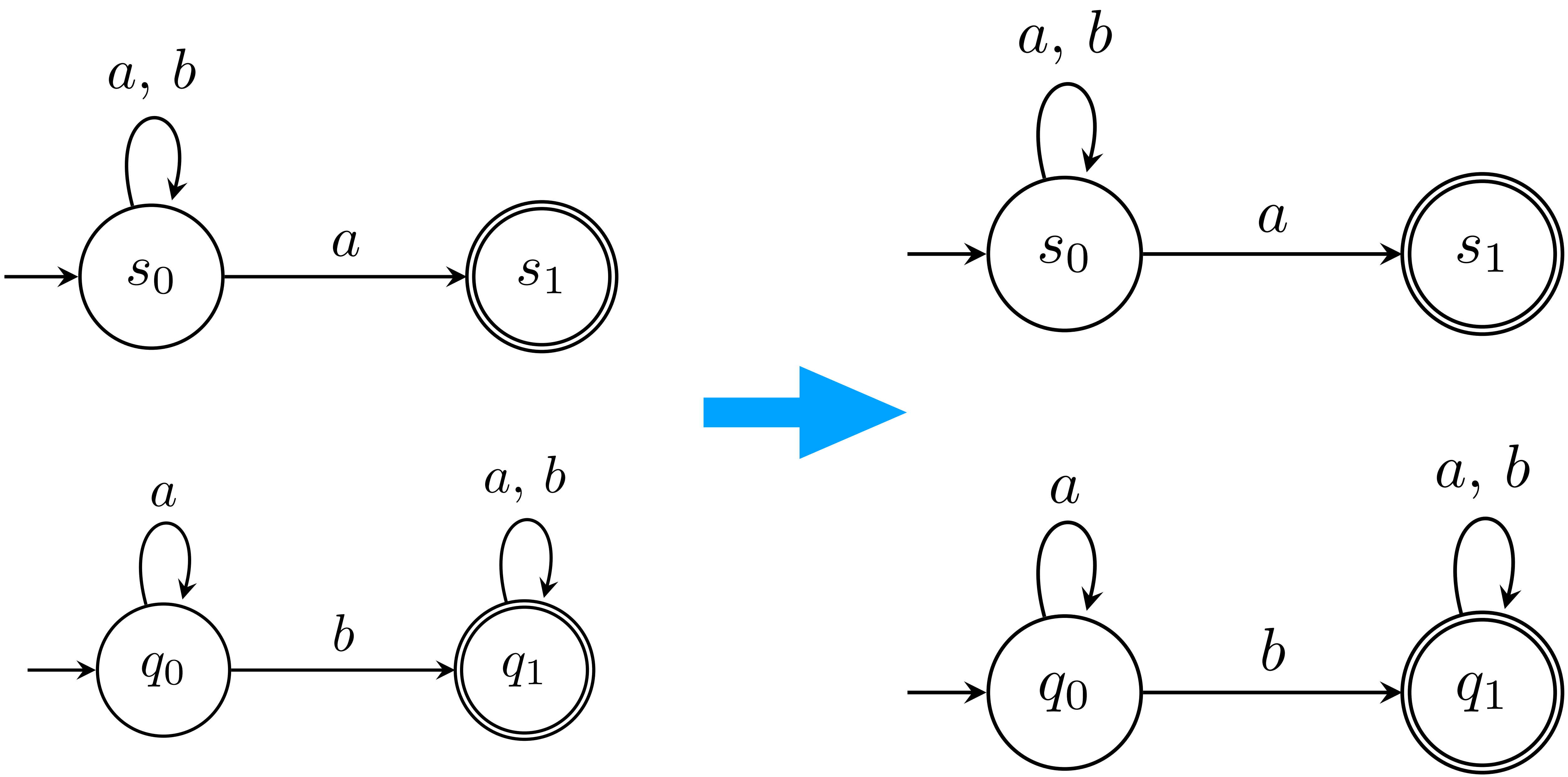
Closure Properties

- Regular languages are *closed* under the following operations.
  - union,
  - intersection,
  - concatenation,
  - Kleene closure, and
  - complementation

Union

- $M_1 = (Q_1, \Sigma, \delta_1, I_1, F_1)$, $M_2 = (Q_2, \Sigma, \delta_2, I_2, F_2)$
- Assume $Q_1 \cap Q_2 = \emptyset$
- $M_3 = (Q_1 \cup Q_2, \Sigma, \delta_3, I_1 \cup I_2, F_1 \cup F_2)$ where $(s, a, t) \in \delta_3$ if
  - $(s, a, t) \in \delta_1$, *or*
  - $(s, a, t) \in \delta_2$
- $L(M_3) = L(M_1) \cup L(M_2)$

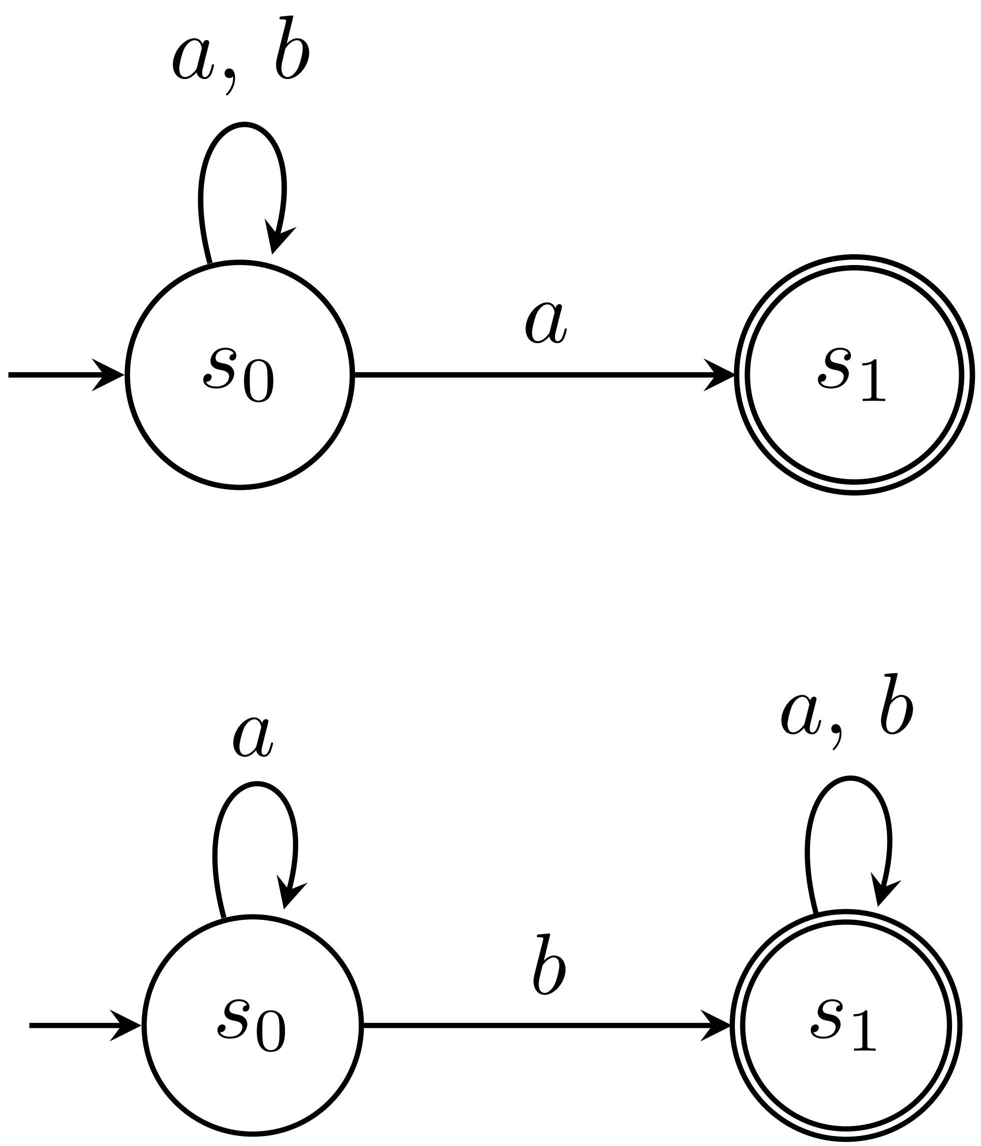
Union
Example

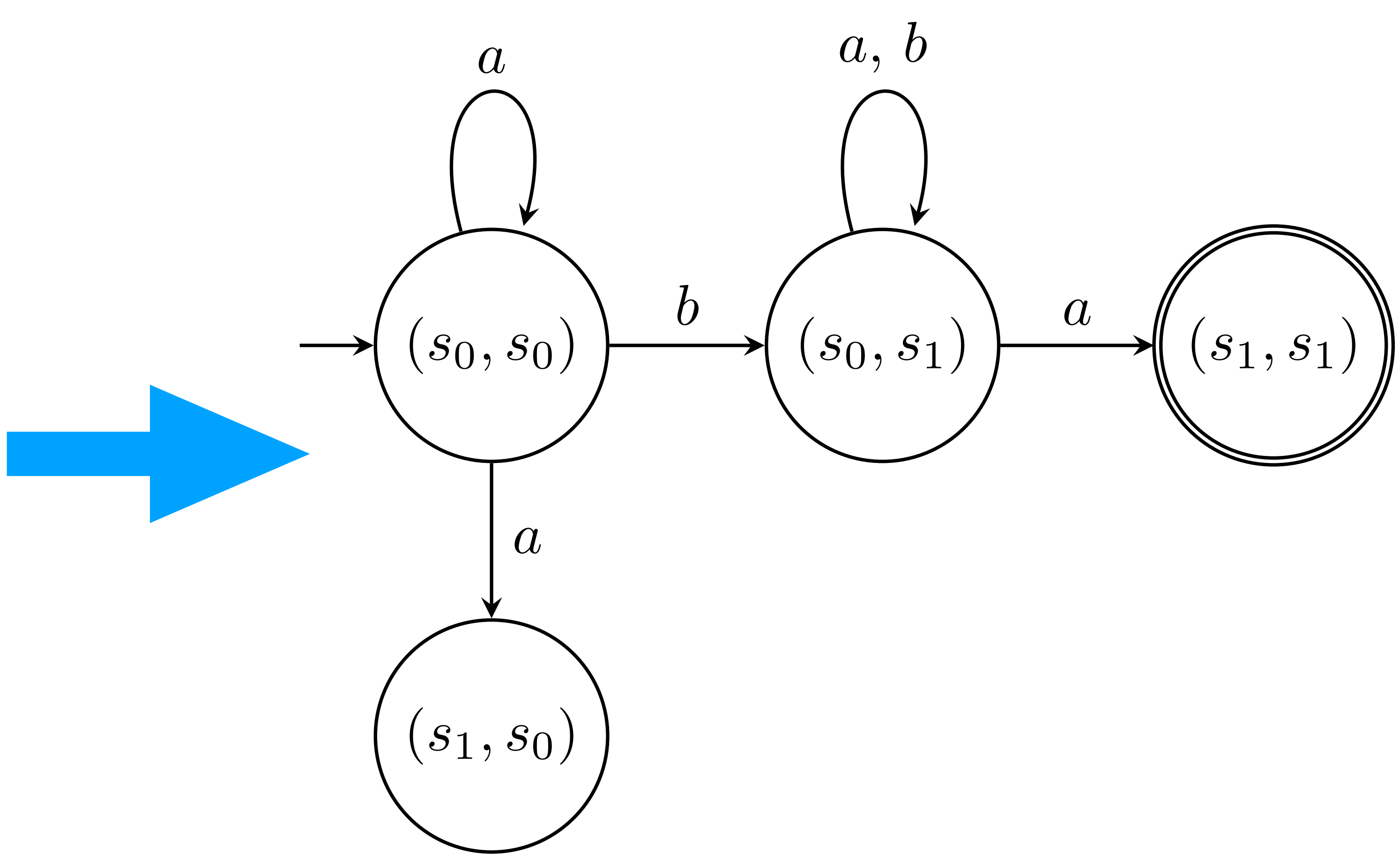


Intersection

- $M_1 = (Q_1, \Sigma, \delta_1, I_1, F_1)$, $M_2 = (Q_2, \Sigma, \delta_2, I_2, F_2)$
- $M_3 = (Q_1 \times Q_2, \Sigma, \delta_3, I_1 \times I_2, F_1 \times F_2)$ where $((s_1, s_2), a, (t_1, t_2)) \in \delta_3$ if
  - $(s_1, a, t_1) \in \delta_1$, and
  - $(s_2, a, t_2) \in \delta_2$
- $L(M_3) = L(M_1) \cap L(M_2)$

Intersection Example

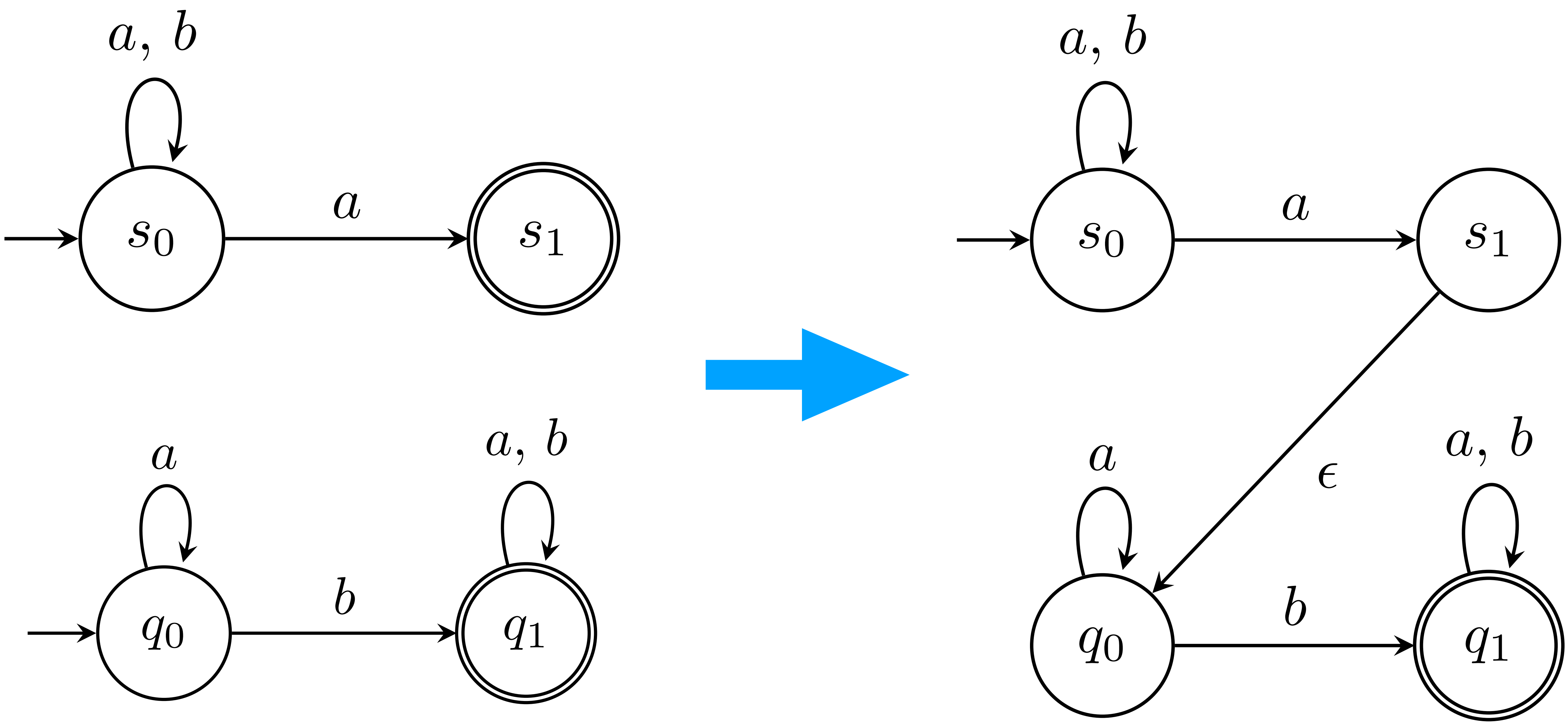




Concatenation

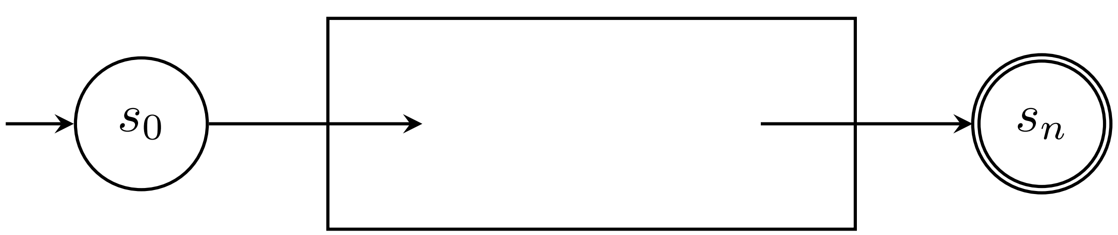
- $M_1 = (Q_1, \Sigma, \delta_1, I_1, F_1)$, $M_2 = (Q_2, \Sigma, \delta_2, I_2, F_2)$
- Assume $Q_1 \cap Q_2 = \varnothing$ and $\epsilon \notin \Sigma$
- $M_3 = (Q_1 \cup Q_2, \Sigma \cup \{\epsilon\}, \delta_3, I_1, F_2)$ where $(s, a, t) \in \delta_3$ if
  - $(s, a, t) \in \delta_1$,
  - $(s, a, t) \in \delta_2$, *or*
  - $a = \epsilon$, $s \in F_1$, and $t \in I_2$.
- $L(M_3) = L(M_1)L(M_2) = \{ uv \mid u \in L(M_1) \textit{ and } v \in L(M_2) \}$

Concatenation Example



Kleene Closure

- An operation that repeat strings accepted by a FSA arbitrary number of times (including zero time)



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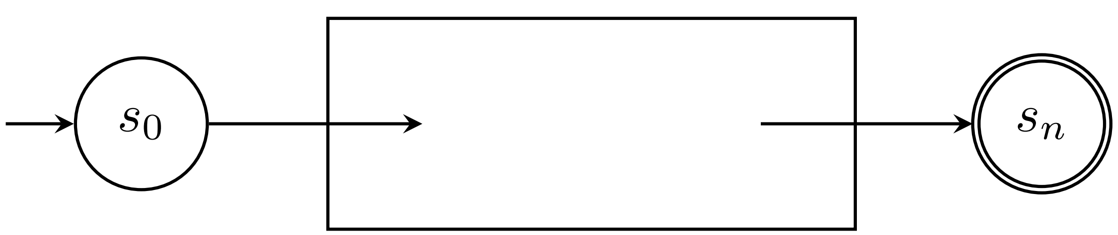

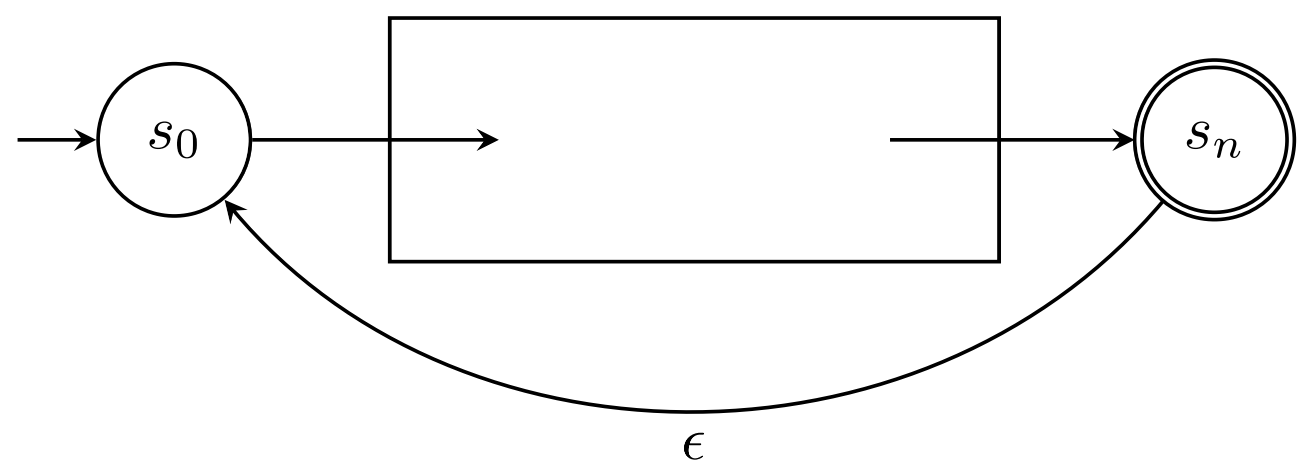

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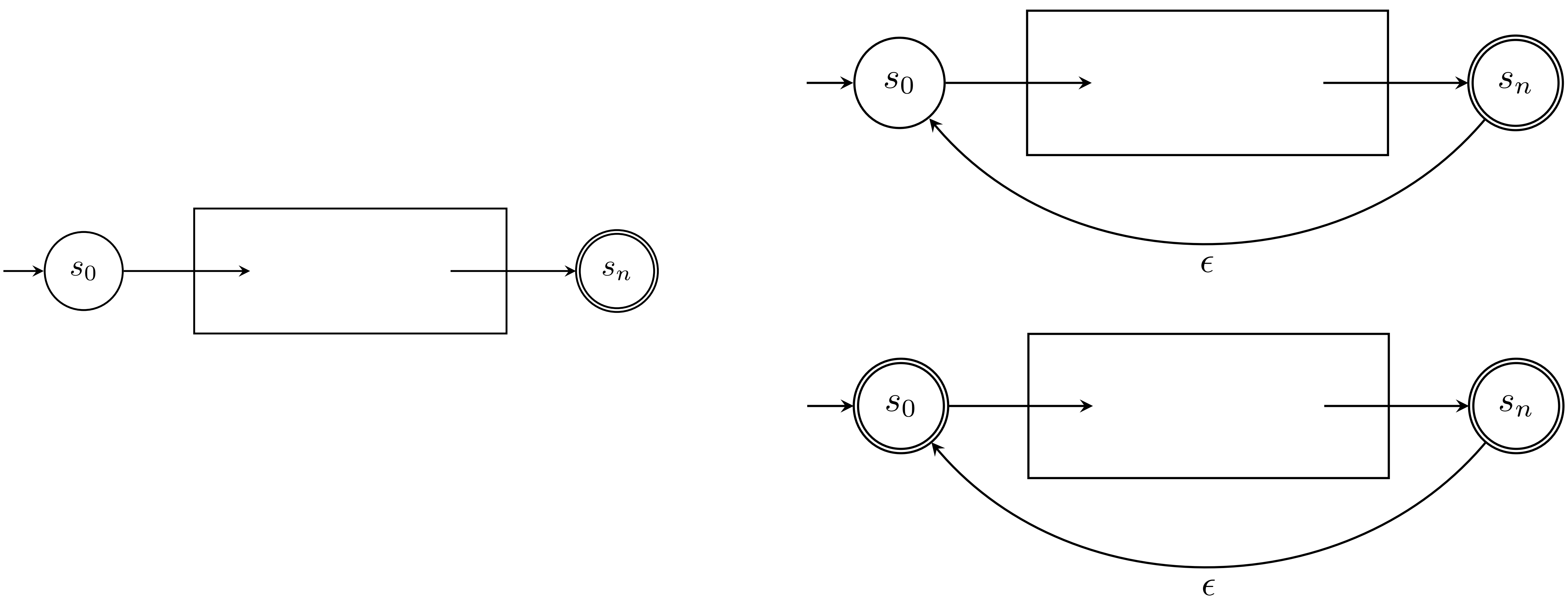
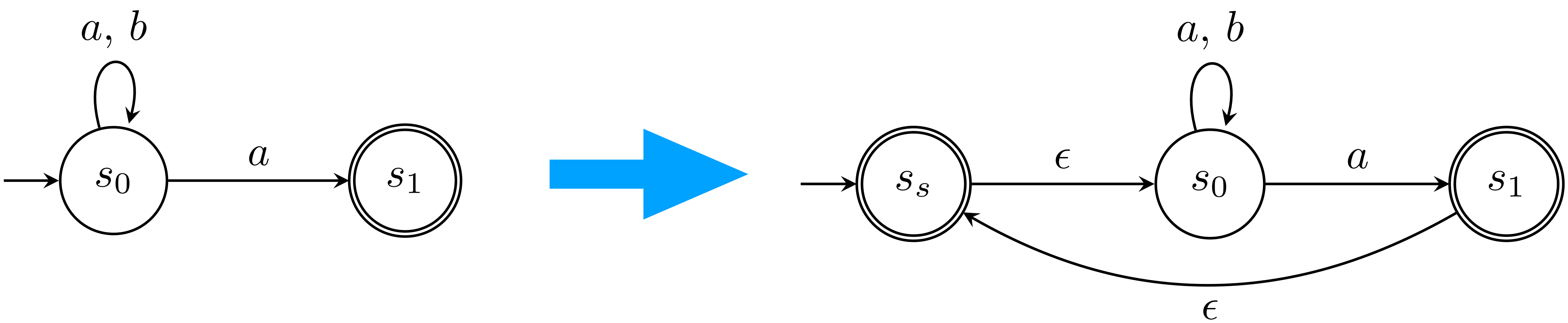


Kleene Closure (cont'd)

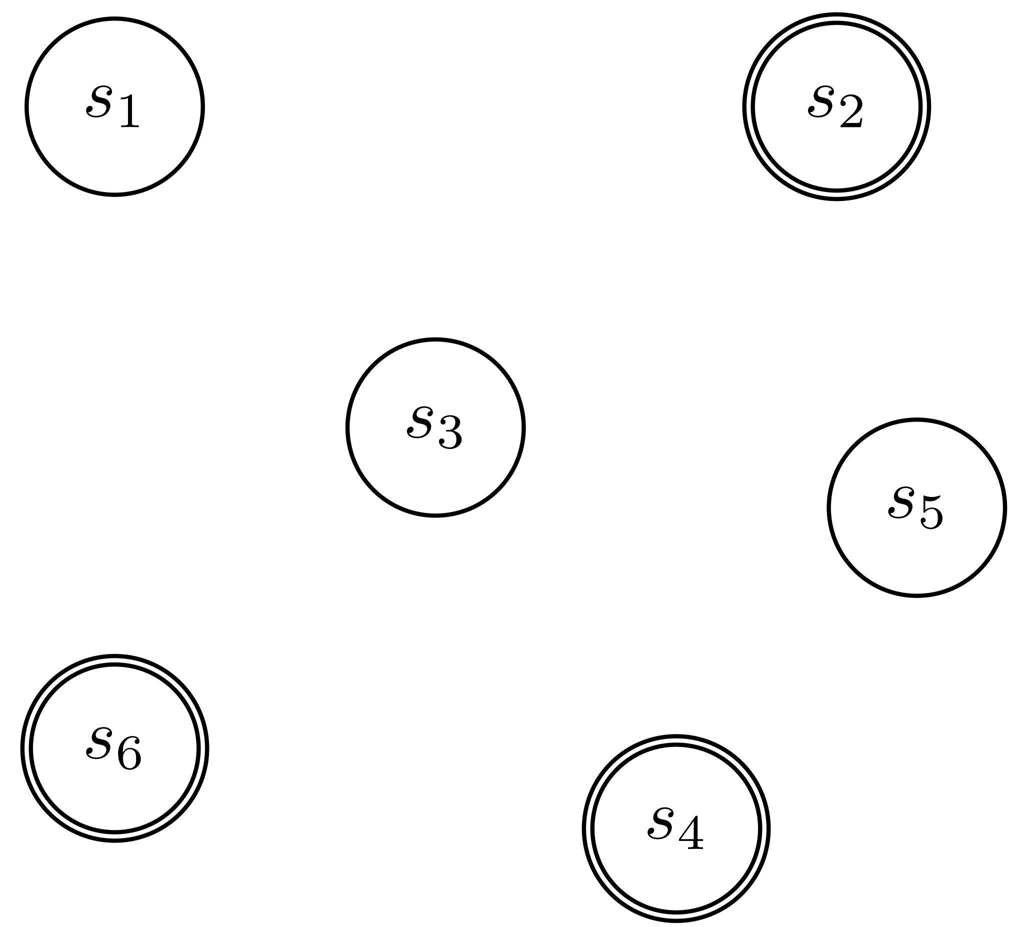
- $M = (Q, \Sigma, \delta, I, F)$
- Assume $\epsilon \notin \Sigma$ and $s_s \notin Q$
- $M' = (Q \cup \{s_s\}, \Sigma \cup \{\epsilon\}, \Delta, \{s_s\}, F \cup \{s_s\})$ where $(s, a, t) \in \Delta$ if
  - $s = s_s$, $t \in I$, and $a = \epsilon$,
  - $(s, a, t) \in \delta$, or
  - $s \in F$, $t \in I$, and $a = \epsilon$
- $L(M') = L(M)^*$

Kleene Closure Example



Complementation
DFA

- $M = (Q, \Sigma, \delta, I, F)$ is a DFA
- $M' = (Q, \Sigma, \delta, I, Q \setminus F)$
- $L(M') = \Sigma^* \setminus L(M)$



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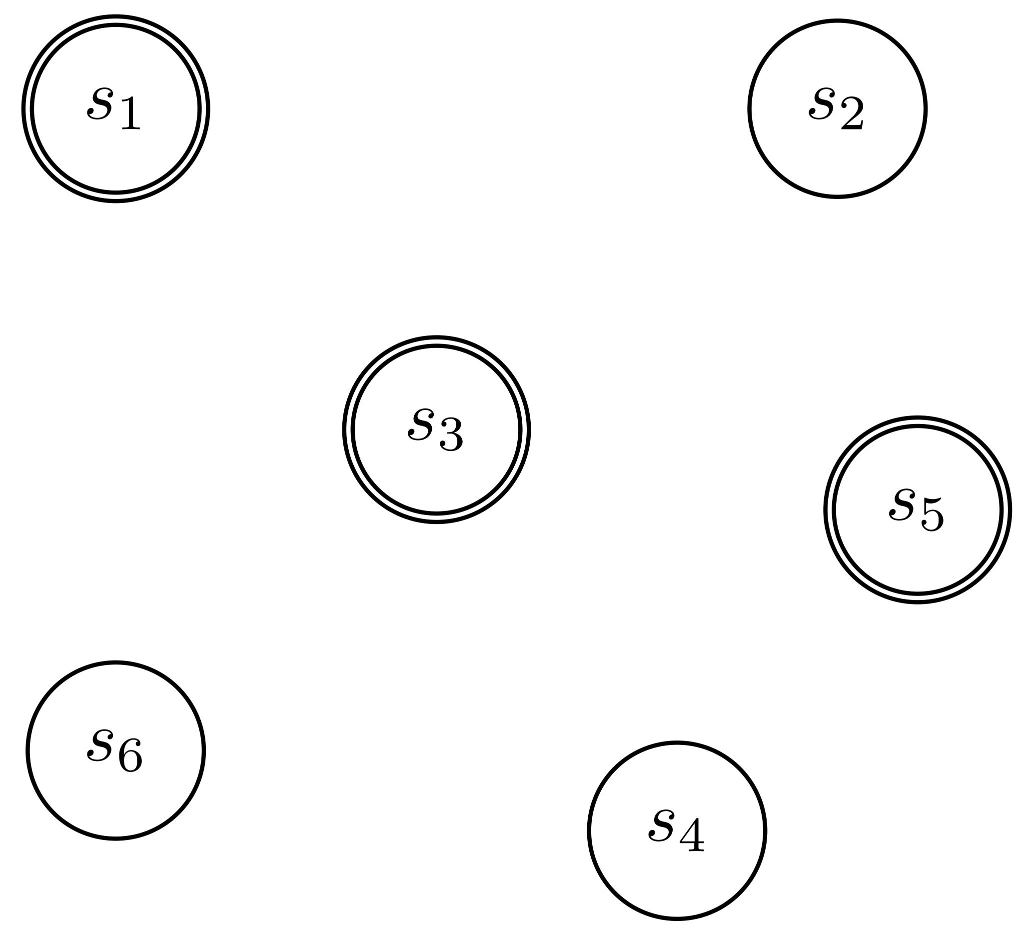
DFA

- $M = (Q,\ \Sigma,\ \delta,\ I,\ F)$ is a DFA
- $M' = (Q,\ \Sigma,\ \delta,\ I,\ Q \setminus F)$
- $L(M') = \Sigma^* \setminus L(M)$



Complementation
NFA

- $M = (Q, \Sigma, \delta, I, F)$ is an NFA.
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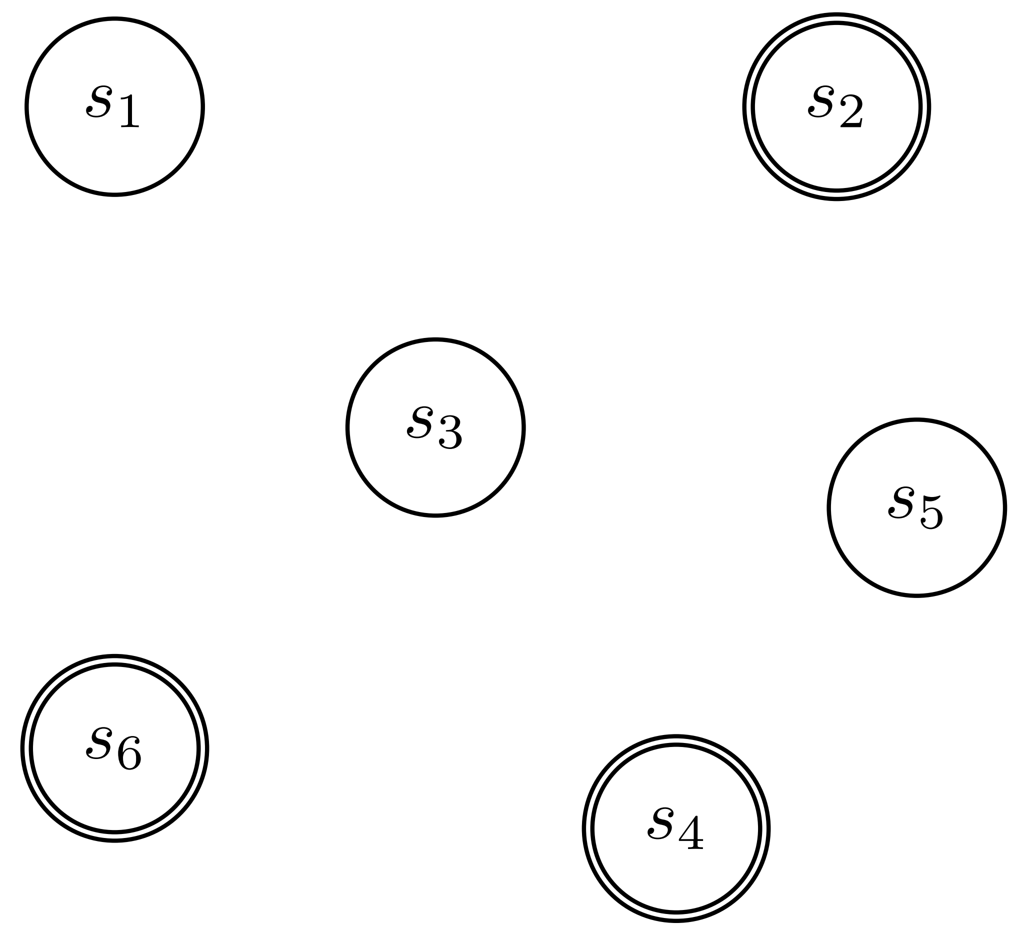
- $L(M') = \Sigma^* \setminus L(M)$?



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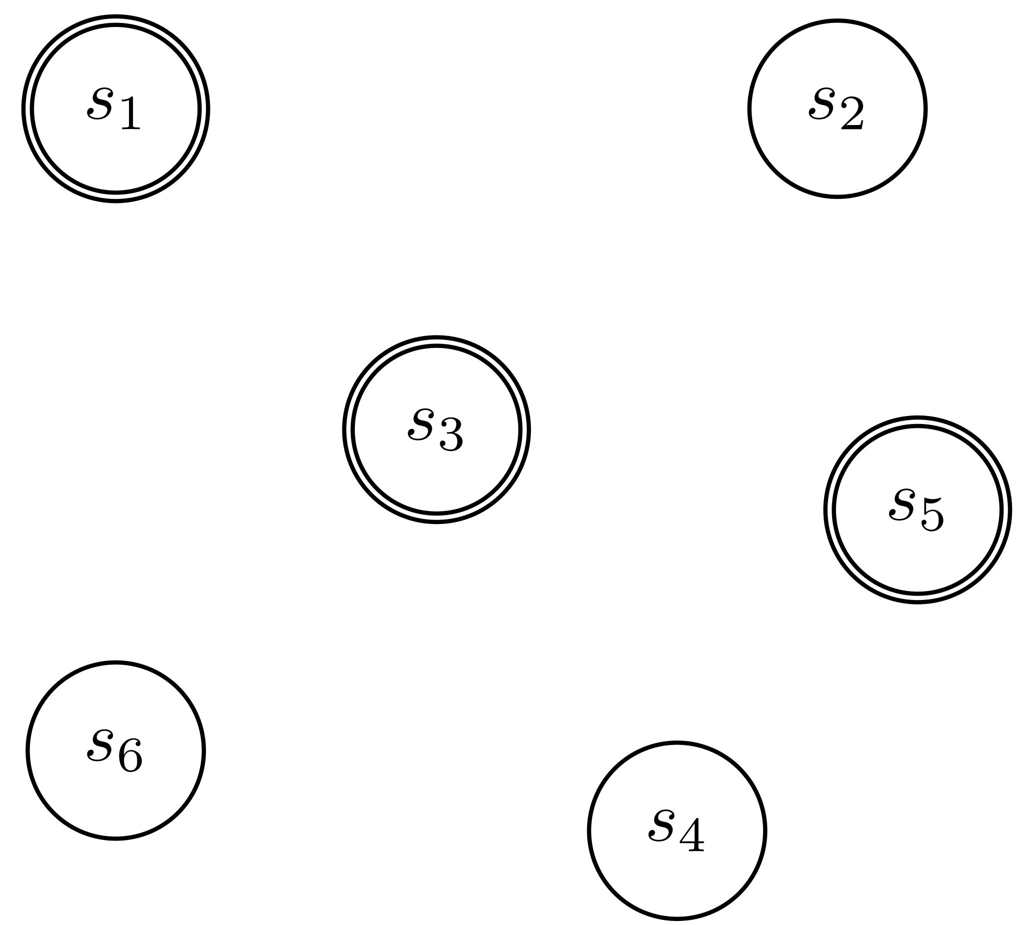

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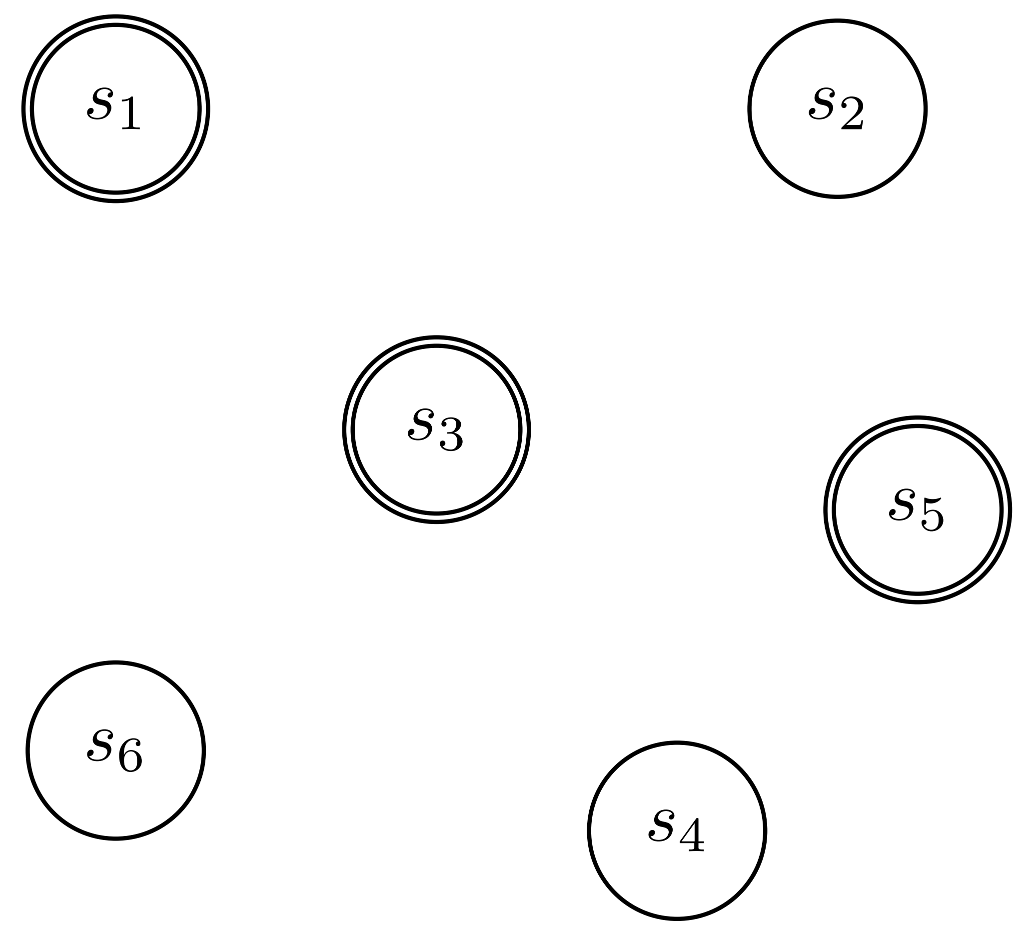

Complementation NFA

- $M = (Q, \Sigma, \delta, I, F)$ is an NFA.
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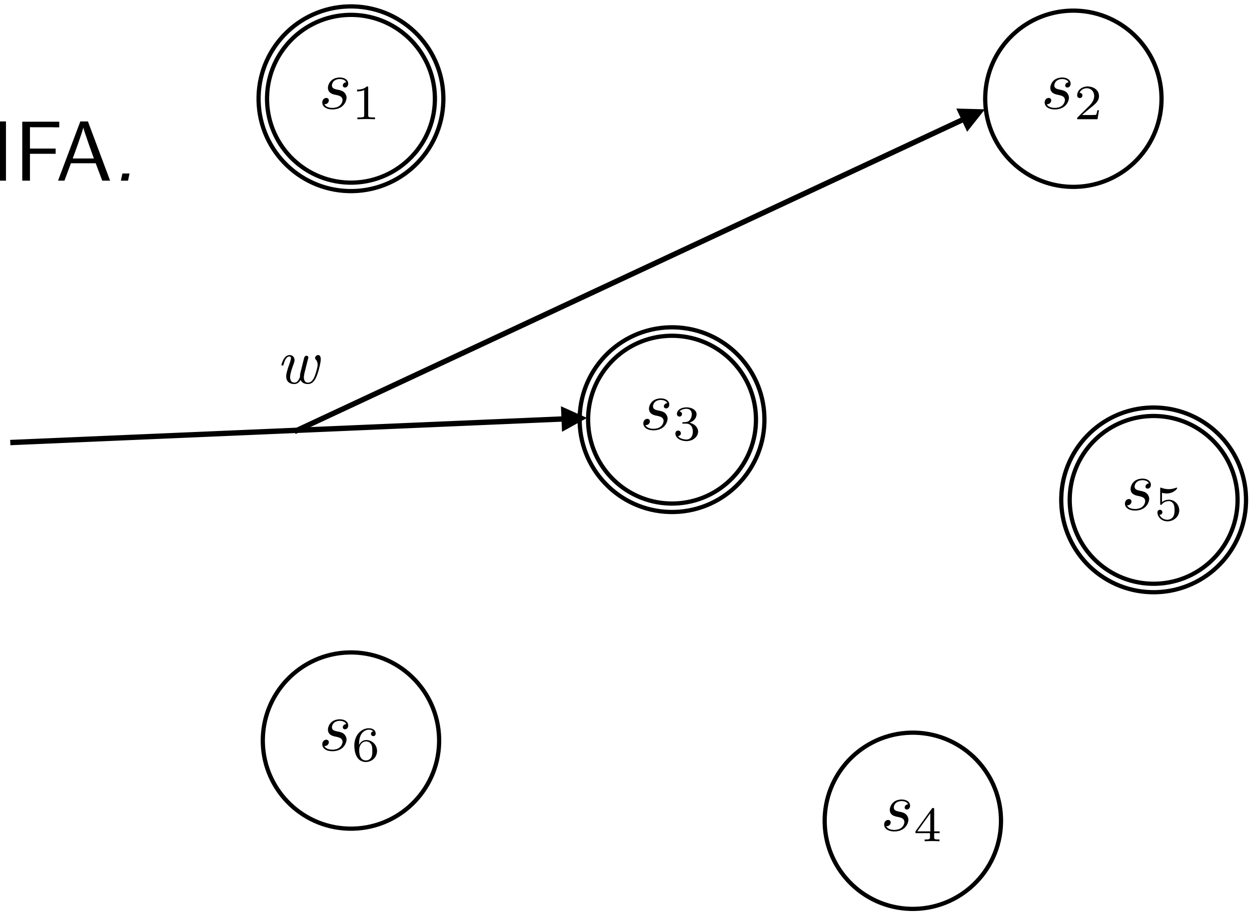


Exercise

- Let $M_1 = (Q_1,\ \Sigma,\ \delta_1,\ I_1,\ F_1)$ and $M_2 = (Q_2,\ \Sigma,\ \delta_2,\ I_2,\ F_2)$ be two NFAs

- Construct an NFA M_3 such that $L(M_3) = L(M_1) \setminus L(M_2)$

- Please describe the components of M_3 in detail

Minimization

- Given a DFA M_1, can we construct a minimal DFA M_2 such that $L(M_1) = L(M_2)$?

- Given an NFA M_1, can we construct a minimal NFA M_2 such that $L(M_1) = L(M_2)$?

Minimization

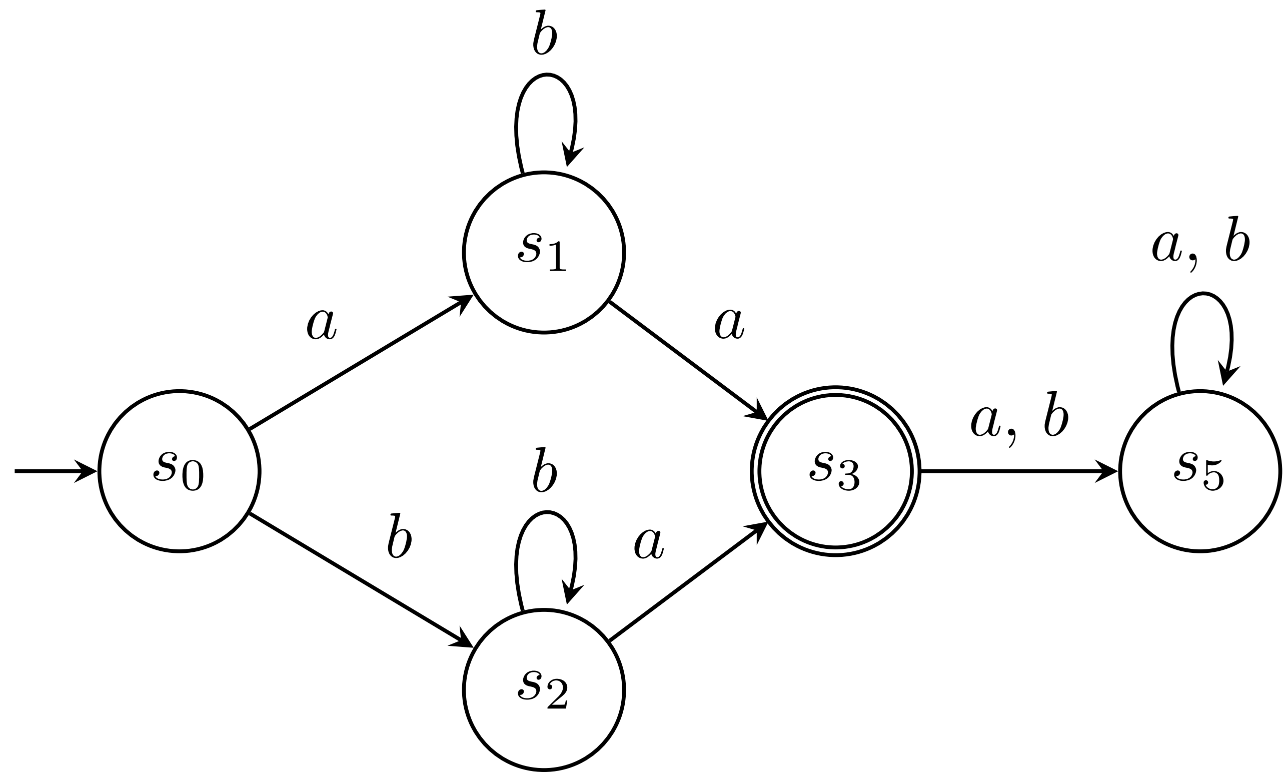
- Given a DFA M_1, can we construct a minimal DFA M_2 such that $L(M_1) = L(M_2)$?
- Given an NFA M_1, can we construct a minimal NFA M_2 such that $L(M_1) = L(M_2)$?

Minimization

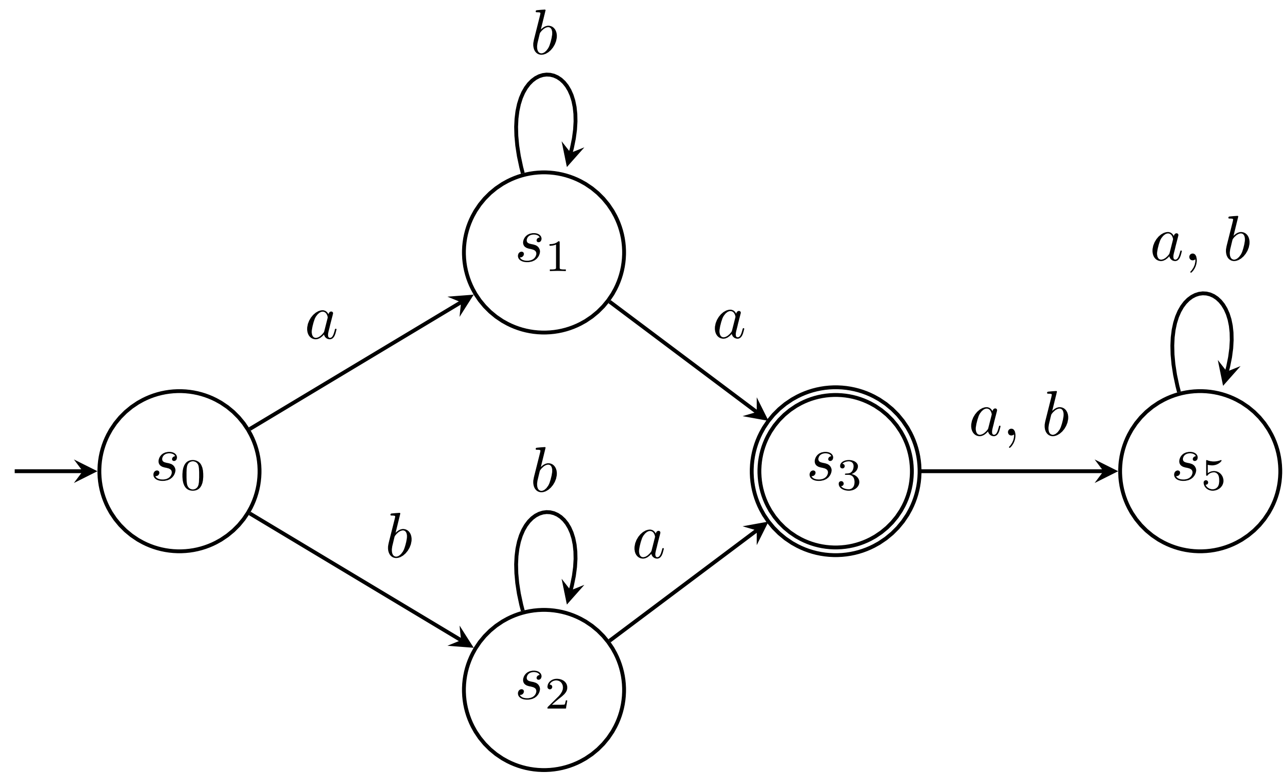
- Given a DFA M_1, can we construct a minimal DFA M_2 such that $L(M_1) = L(M_2)$? ○

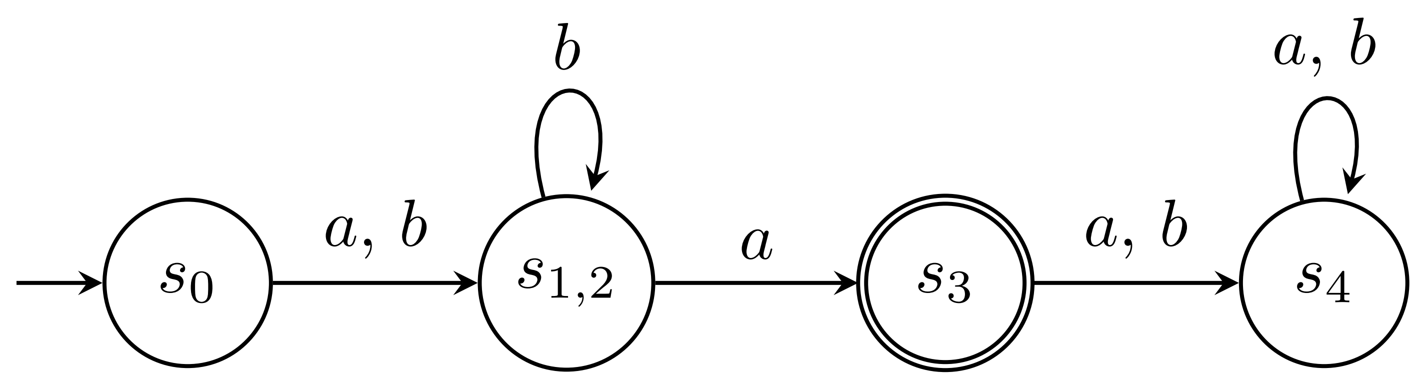
- Given an NFA M_1, can we construct a minimal NFA M_2 such that $L(M_1) = L(M_2)$? ○ **but harder**

Intuition



Intuition





Myhill-Nerode Theorem

- Given a language $L \subseteq \Sigma^*$, define a binary relation R_L over Σ^* as follows
  - $x R_L y$ *iff* $\forall z \in \Sigma^* \, (xz \in L \leftrightarrow yz \in L)$
- R_L can be shown to be an equivalence relation
- R_L divide the set of string into *equivalence classes*
- L is regular iff R_L has a finite number of equivalence classes
- The number of states in the minimal DFA recognizing L is equal to the number of equivalence classes in R_L

Minimization
Idea

- For a language $L \subseteq \Sigma^*$, compute the equivalence classes of L

- Construct a state for each equivalence class

- A equivalence class C_1 can take an a-transition to another equivalence class C_2 if there is a string $x \in C_1$ such that $xa \in C_2$

- How to find the equivalence classes?

Minimization

Hopcroft's Algorithm

```
P := {F, Q \ F};
W := {F};
while (W is not empty) do
     choose and remove a set A from W
     for each c in Σ do
          let X be the set of states for which a transition on c leads to a state in A
          for each set Y in P for which X ∩ Y is nonempty and Y \ X is nonempty do
               replace Y in P by the two sets X ∩ Y and Y \ X
               if Y is in W
                    replace Y in W by the same two sets
               else
                    if |X ∩ Y| <= |Y \ X|
                         add X ∩ Y to W
                    else
                         add Y \ X to W
          end;
     end;
end;
```

the pseudocode is taken from https://en.wikipedia.org/wiki/DFA_minimization

Language Expressions

- So far we know that a regular language can be accepted by a finite state automaton
- Can we represent a regular language in other forms?

Language Expressions

- So far we know that a regular language can be accepted by a finite state automaton
- Can we represent a regular language in other forms?

regular expressions

Regular Expressions (RE)

- Let Σ be an alphabet
- The regular expressions over Σ are defined as follows
  - $\varnothing$ is a regular expression denoting the empty set;
  - ϵ is a regular expression denoting the set $\{\epsilon\}$;
  - for each $a \in \Sigma$, a is a regular expression denoting the set $\{a\}$;
  - if r and s are regular expressions denoting the sets R and S respectively, then $r+s$, rs, and r^* are regular expressions denoting $R \cup S$, RS, and R^* respectively
- The language of a regular expression e is denoted by $L(e)$

Regular Expressions Examples

- Let $\Sigma = \{a,\ b\}$
- $a^*ba^* = \{w \mid w \textit{ has exactly a single } b\}$
- $\Sigma^*b\Sigma^* = \{w \mid w \textit{ has at least one } b\}$
- $\Sigma^*aba\Sigma^* = \{w \mid w \textit{ has a substring } aba\}$
- $a{+}b{+}a\Sigma^*a{+}b\Sigma^*b = \{w \mid w$ *starts and ends with the same symbol*$\}$

Regular Expressions Examples (cont'd)

- $r + \emptyset = ?$

- $r + \epsilon = ?$

- $r\emptyset = ?$

- $r\epsilon = ?$

Regular Expressions Examples (cont'd)

- $r + \emptyset = ?$ $\quad r$

- $r + \epsilon = ?$

- $r\emptyset = ?$

- $r\epsilon = ?$

Regular Expressions Examples (cont'd)

- $r{+}\emptyset = ?$ r
- $r{+}\epsilon = ?$ $r{+}\epsilon$
- $r\emptyset = ?$
- $r\epsilon = ?$

Regular Expressions Examples (cont'd)

- $r + \emptyset = ?$ r
- $r + \epsilon = ?$ $r + \epsilon$
- $r\emptyset = ?$ $\emptyset$
- $r\epsilon = ?$

Regular Expressions Examples (cont'd)

- $r+\emptyset = ?$ $\quad r$
- $r+\epsilon = ?$ $\quad r+\epsilon$
- $r\emptyset = ?$ $\quad \emptyset$
- $r\epsilon = ?$ $\quad r$

Exercise

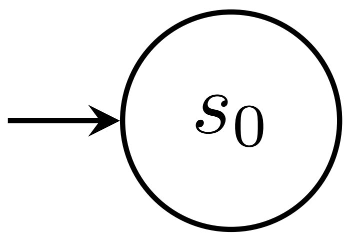
- Write regular expressions to describe the following languages. ($\Sigma = \{a, b\}$)
  - $\{w \mid \text{the length of } w \text{ is even}\}$
  - $\{w \mid w \text{ has at most two } b\text{'s}\}$
  - $\{w \mid \text{every } a \text{ in } w \text{ is followed by } b\}$

Regular Expressions
VS
Finite State Automata

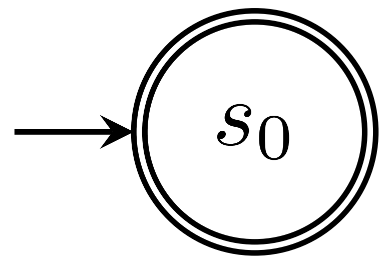
- A language is recognized by an NFA if and only if some regular expression describes it

- A language is regular if and only if some regular expression describes it

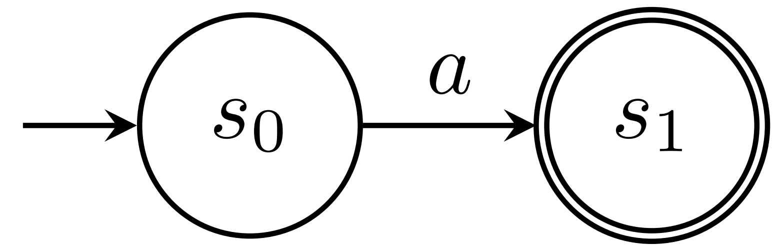
From RE to NFA



$\varnothing$



ϵ



a

Let A_r be an NFA recognizing the language of a regular expression r

$r+s$: union of A_r and A_s

rs: concatenation of A_r and A_s

r^*: the Kleene closure of A_r

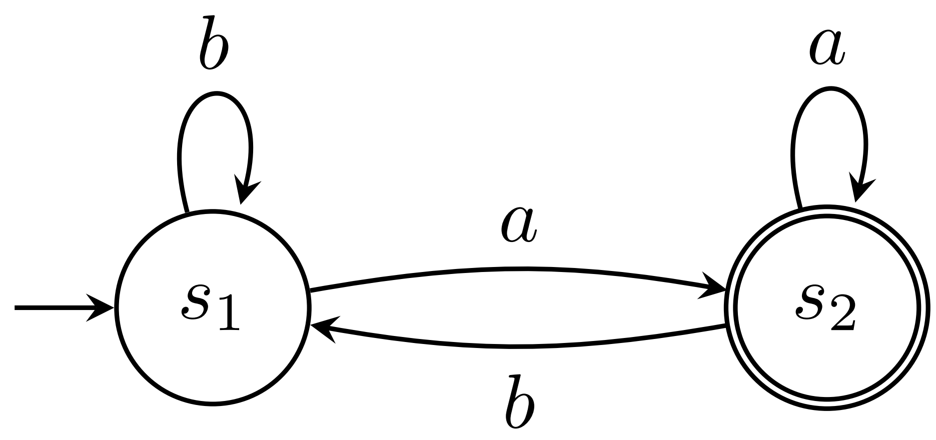
From NFA to RE

- Transitive Closure Method
- State Removal Method
- Brzozowski Algebraic Method

Transitive Closure Method

- Let $D = (\{s_1, \ldots, s_n\}, \Sigma, \delta, \{s_1\}, F)$ be a DFA
- Define
  - $R_{ij}{}^{0} = \{a \mid (s_i, a, s_j) \in \delta\}$ if $i \neq j$
  - $R_{ij}{}^{0} = \{a \mid (s_i, a, s_j) \in \delta\} \cup \{\epsilon\}$ if $i = j$
  - $R_{ij}{}^{k} = R_{ik}{}^{k\text{-}1}(R_{kk}{}^{k\text{-}1})^{*}R_{kj}{}^{k\text{-}1} \cup R_{ij}{}^{k\text{-}1}$
- $R_{ij}{}^{k}$ represents the inputs that cause D to go from s_i to s_j without passing through a state higher than s_k
- $R_{ij}{}^{k}$ can be denoted by regular expressions
- $L(D) = \cup_{S_j \in F} R_{1j}{}^{n}$

Transitive Closure Method Example



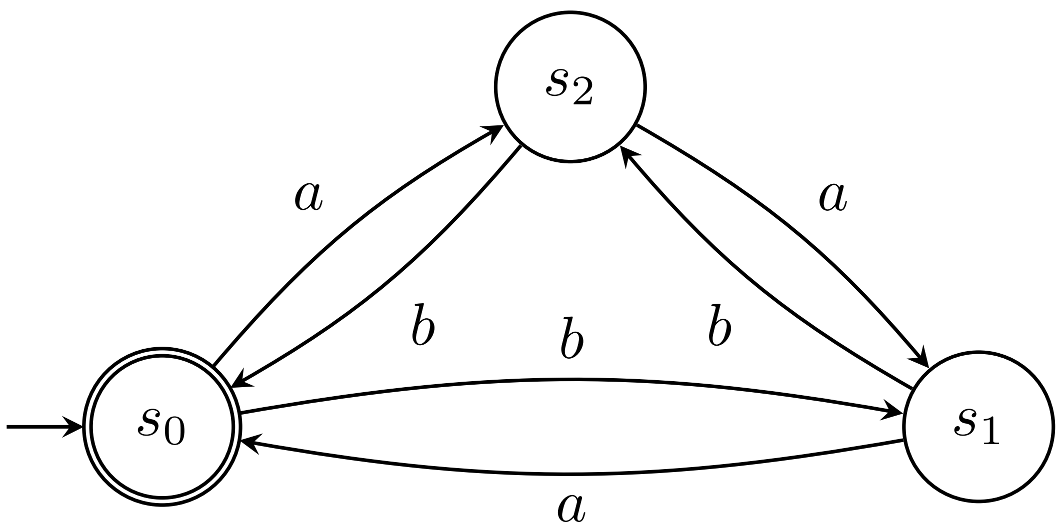
	$k = 0$	$k = 1$	$k = 2$
${R_{11}}^k$	$b+\epsilon$	$(b+\epsilon)(b+\epsilon)^*(b+\epsilon)+(b+\epsilon)$ $= b^*$	
${R_{12}}^k$	a	$(b+\epsilon)(b+\epsilon)^*a+a$ $= b^*a$	$b^*a(b^*a+\epsilon)^*(b^*a+\epsilon)+b^*a$ $= (a+b)^*a$
${R_{21}}^k$	b	$b(b+\epsilon)^*(b+\epsilon)+b$ $= b^+$	
${R_{22}}^k$	$a+\epsilon$	$b(b+\epsilon)^*a+(a+\epsilon)$ $= b^*a+\epsilon$	

$b^+ = bb^*$

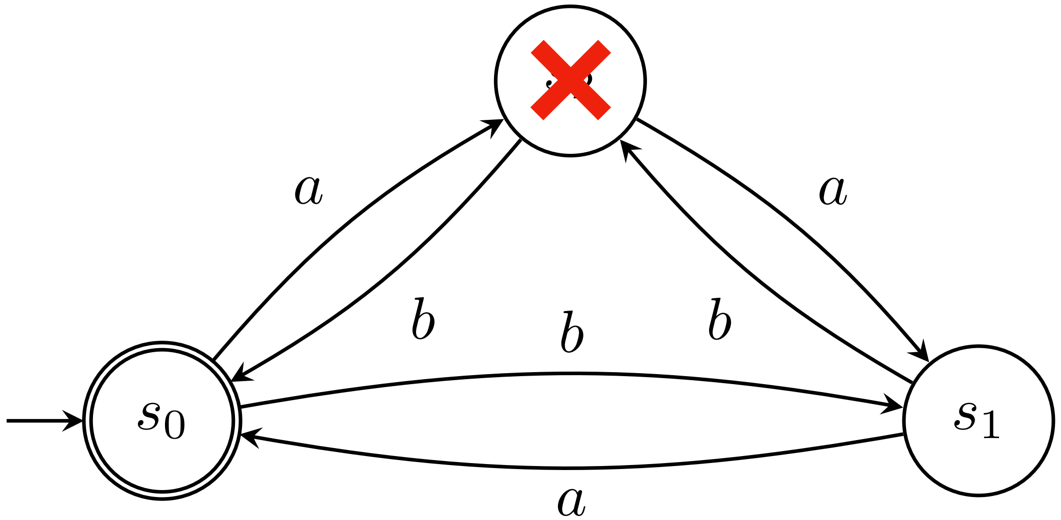
State Removal Method

- Make the NFA has a single accepting state
- Make the NFA has a single initial state
- Remove states and change transition labels (may be regular expressions) until there is only the initial state and the accepting state
- Compute the regular expression

State Removal Method Example



State Removal Method Example



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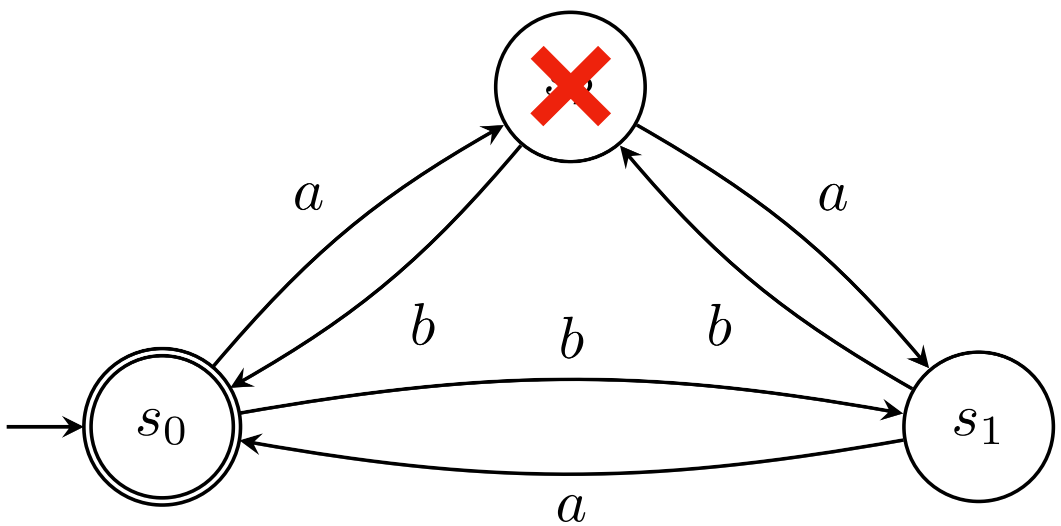

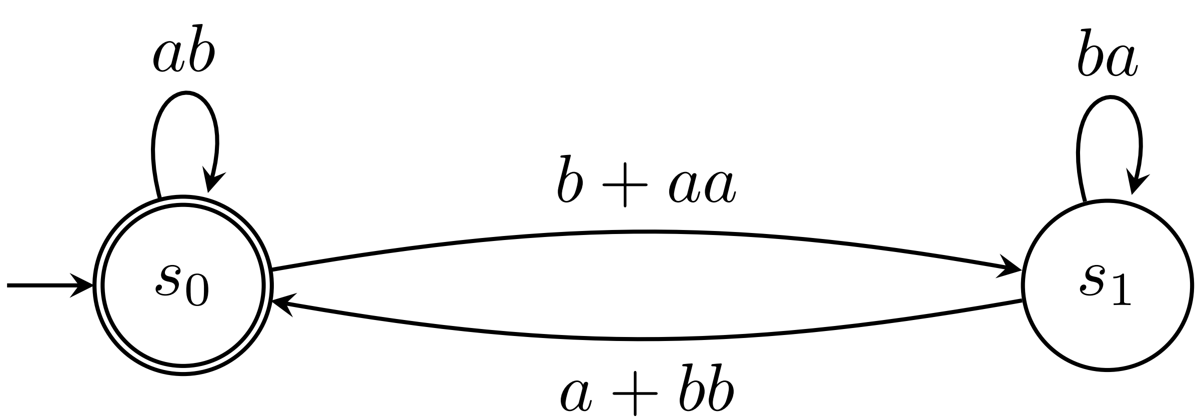

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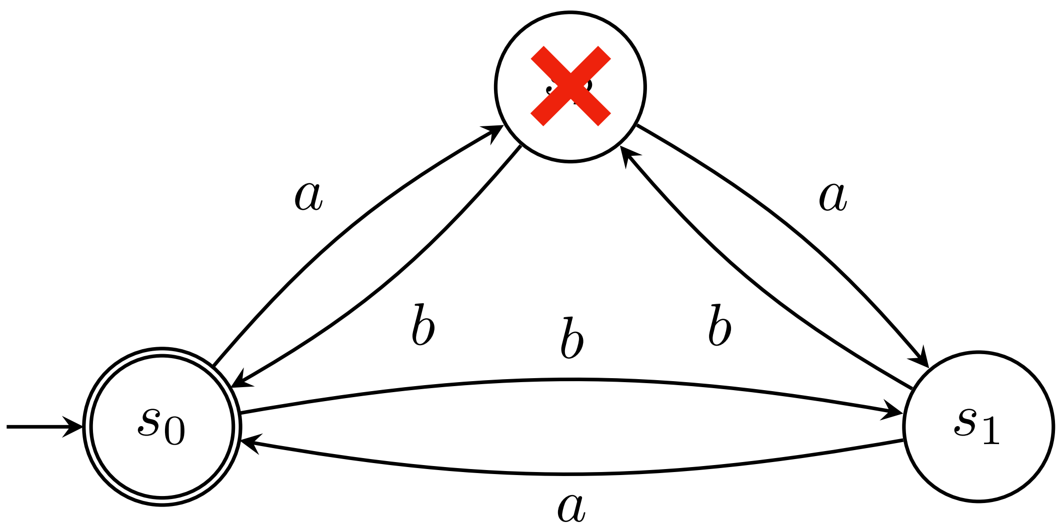
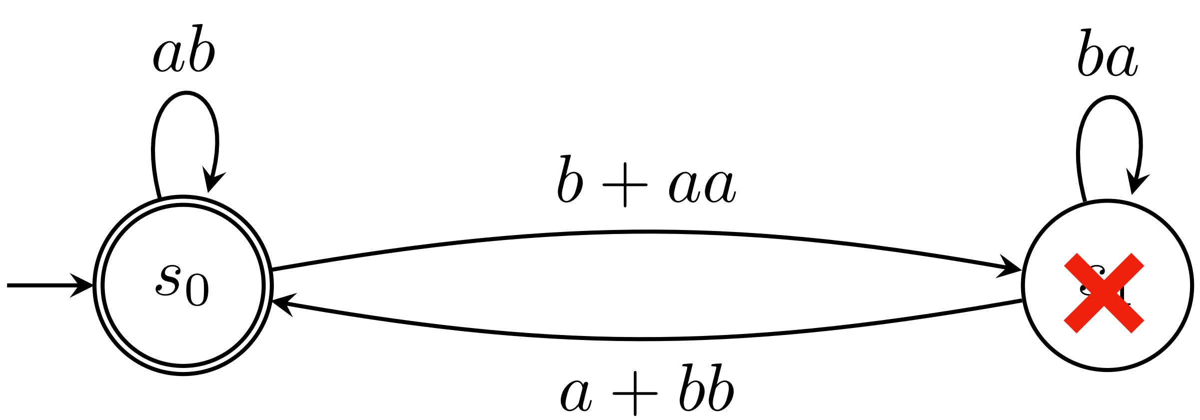
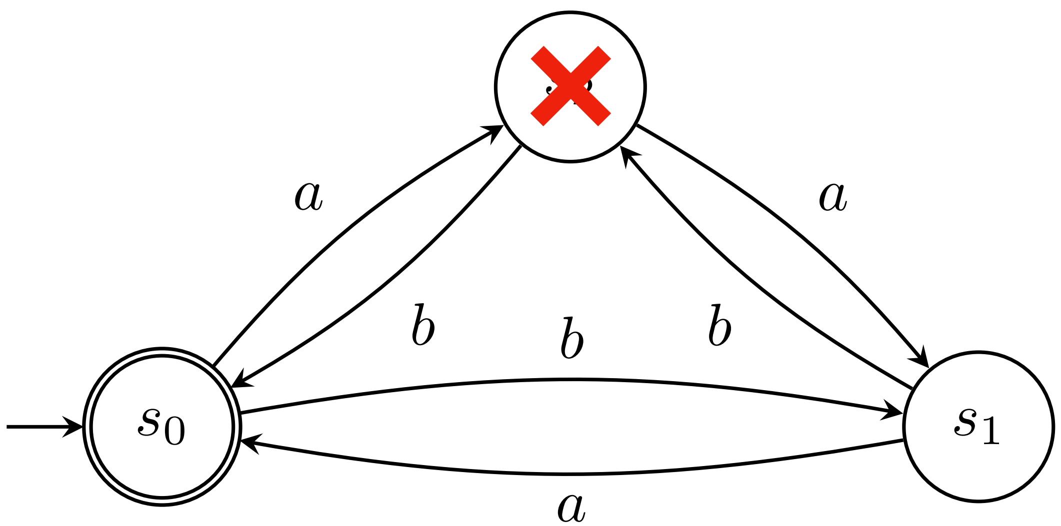
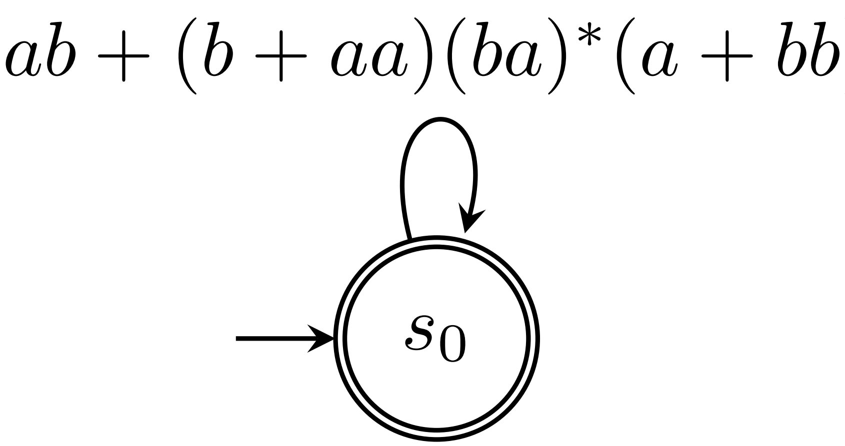
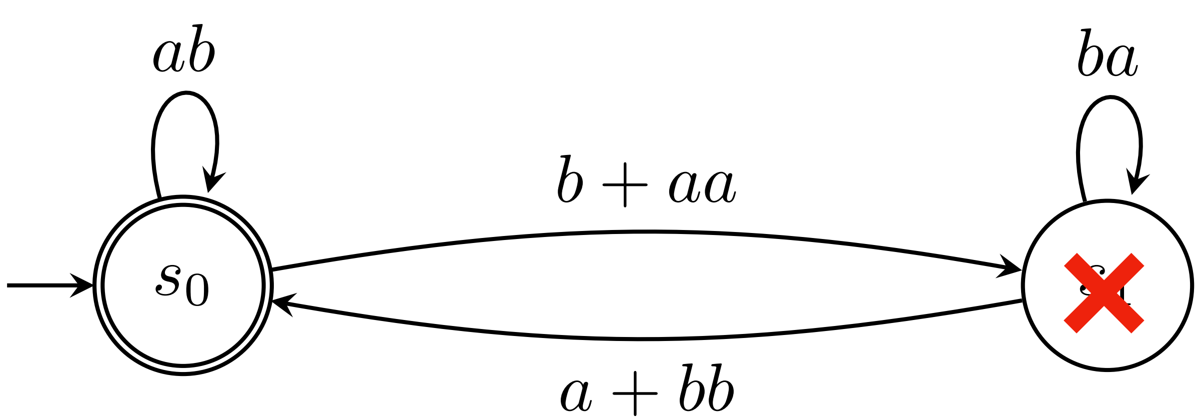




State Removal Method Example







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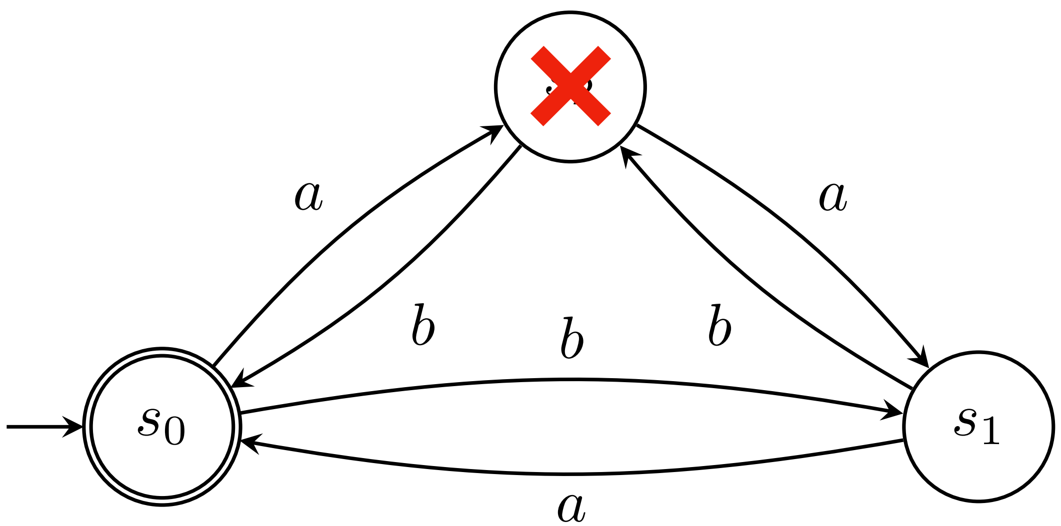
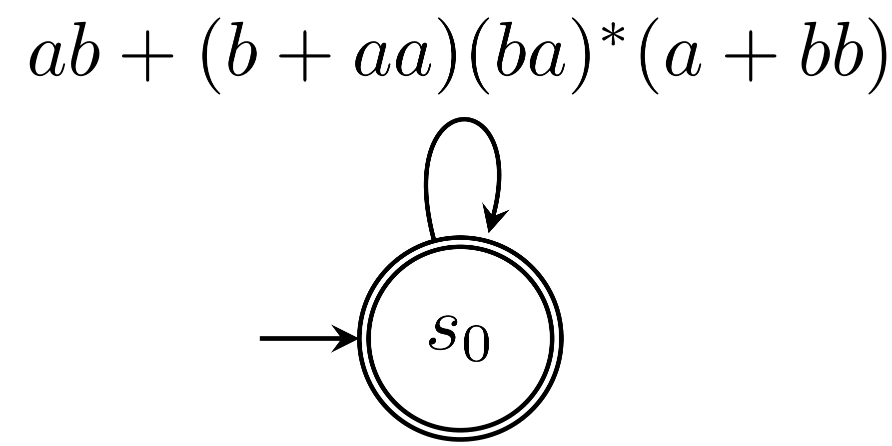
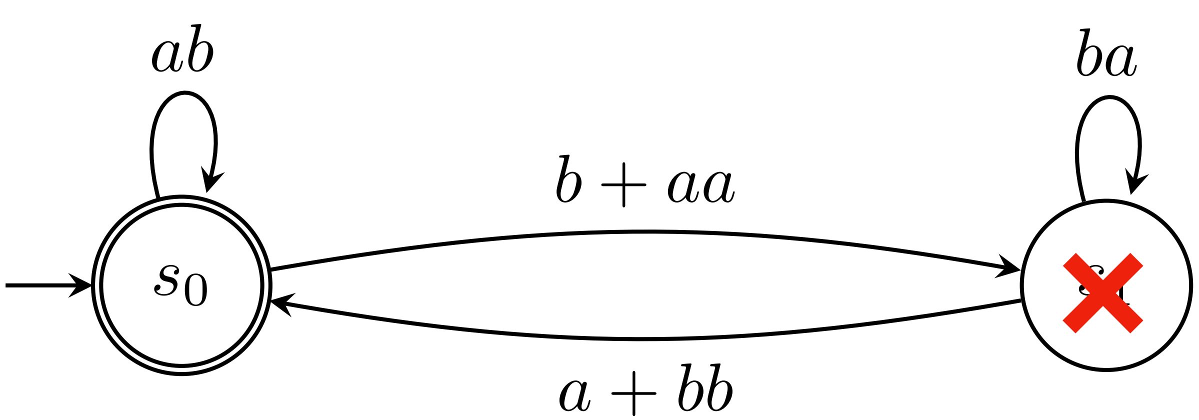






$$(ab+(b+aa)(ba)^*(a+bb))^*$$

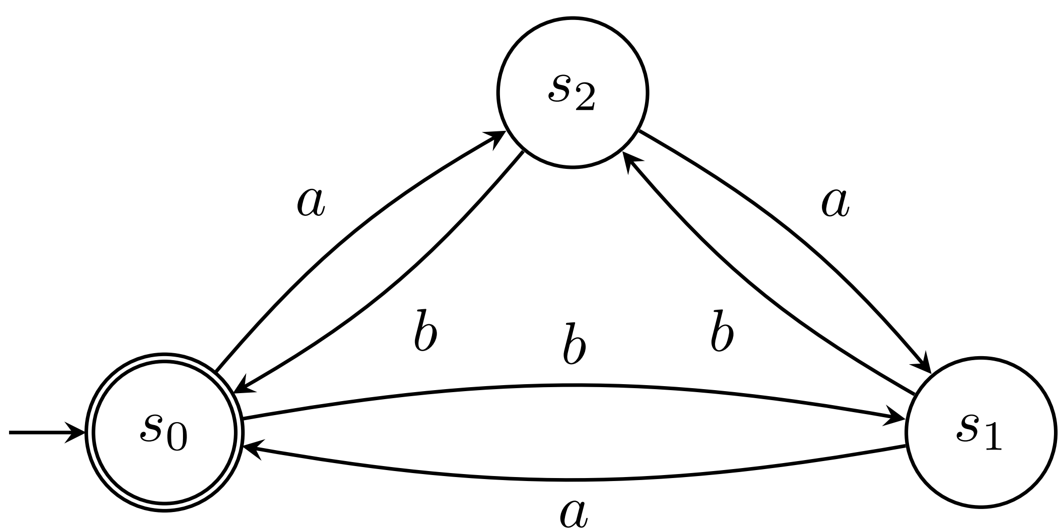
Brzozowski Algebraic Method

- $M = (Q, \Sigma, \delta, \{q_0\}, F)$ is an NFA containing no ϵ-transitions
- For every q_i, create the equation

$$Q_i = +_{q_i \xrightarrow{a} q_j} aQ_j + \begin{cases} \{\epsilon\}, \text{if } q_i \in F \\ \varnothing, \text{else} \end{cases}$$

- Solve the equation system and find Q_0

Brzozowski Algebraic Method Example



$$Q_0 = bQ_1 + aQ_2 + \epsilon$$
$$Q_1 = aQ_0 + bQ_2$$
$$Q_2 = bQ_0 + aQ_1$$

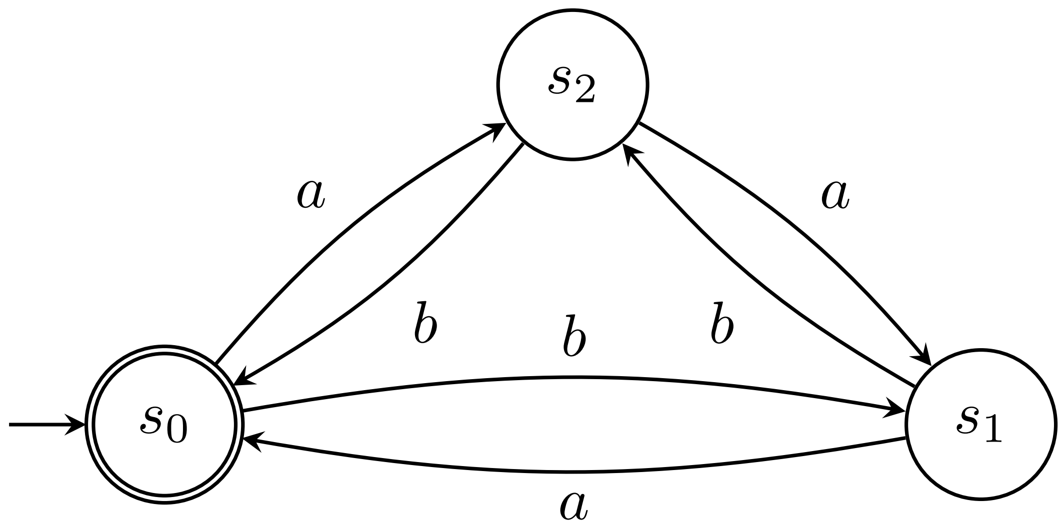
$$\begin{aligned} Q_2 &= bQ_0 + aQ_1 \\ &= bQ_0 + a(aQ_0 + bQ_2) \\ &= abQ_2 + (b+aa)Q_0 \end{aligned}$$

by Arden's Lemma:

$L = UL+V$ iff $L = U^*V$ where $L, U, V \subseteq \Sigma^*$ with $\epsilon \notin U$

$$Q_2 = (ab)^*(b+aa)Q_0$$

Brzozowski Algebraic Method Example (cont'd)



$$Q_0 = bQ_1 + aQ_2 + \epsilon$$
$$Q_1 = aQ_0 + bQ_2$$
$$Q_2 = bQ_0 + aQ_1$$

$$Q_2 = (ab)^*(b+aa)Q_0$$

$$\begin{aligned} Q_0 &= bQ_1 + aQ_2 + \epsilon \\ &= b(aQ_0 + bQ_2) + aQ_2 + \epsilon \\ &= baQ_0 + (bb+a)Q_2 + \epsilon \\ &= (ba+(bb+a)(ab)^*(b+aa))Q_0 + \epsilon \end{aligned}$$

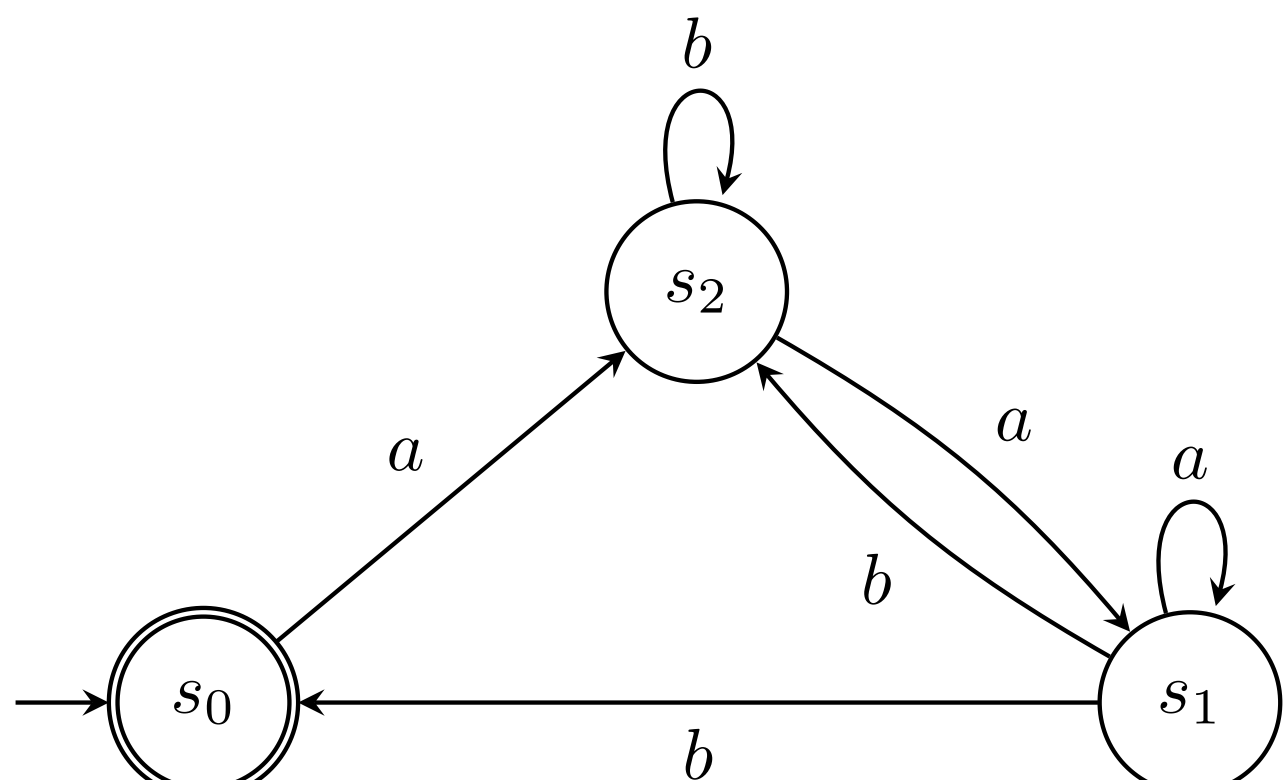
by Arden's Lemma:

$L = UL+V$ iff $L = U^*V$ where $L, U, V \subseteq \Sigma^*$ with $\epsilon \notin U$

$$Q_0 = (ba+(bb+a)(ab)^*(b+aa))^*$$

Exercise

- Express the language of the following automaton by a regular expression



WS1S

- Syntax of S1S (monadic second-order logic of one successor)
  - First-order variable set: $V = \{x_1,\ x_2,\ ...\}$
  - Second-order variable set: $X = \{X_1,\ X_2,\ ...\}$
  - Terms: $t ::= 0 \mid x_i$
  - Formulas: $\varphi ::= S(t,\ t) \mid X_i(t) \mid \neg\ \varphi \mid \varphi \wedge \varphi \mid \exists x_i.\varphi \mid \exists X_i.\varphi$
- S is the successor predicate
- WS1S: fragment of S1S which allows only quantification over finite sets

Semantics of S1S

- Signature $\langle \mathbb{N}, S \rangle$
- Interpretation $\sigma = \langle \sigma_1, \sigma_2 \rangle, \sigma_1 : V \to \mathbb{N}, \sigma_2 : X \to 2^{\mathbb{N}}$
- *Satisfiability*

$$\begin{array}{lll}
\sigma \models X(t) & \textit{iff} & \sigma(t) \in \sigma(X) \\
\sigma \models S(t,t') & \textit{iff} & \sigma(t) + 1 = \sigma(t') \\
\sigma \models \neg\varphi & \textit{iff} & \sigma \not\models \varphi \\
\sigma \models \varphi_1 \wedge \varphi_2 & \textit{iff} & \sigma \models \varphi_1 \text{ and } \sigma \models \varphi_2 \\
\sigma \models \exists x.\varphi & \textit{iff} & \sigma[n/x] \models \varphi \text{ for some } n \in \mathbb{N} \\
\sigma \models \exists X.\varphi & \textit{iff} & \sigma[N/X] \models \varphi \text{ for some } N \in 2^{\mathbb{N}}
\end{array}$$

- *Validity*

$$\models \varphi \quad \textit{iff} \quad \sigma \models \varphi \text{ for all interpretations } \sigma$$

Abbreviations

$$
\begin{array}{lcl}
\varphi_1 \vee \varphi_2 & := & \neg(\neg\varphi_1 \wedge \neg\varphi_2) \\
\varphi_1 \rightarrow \varphi_2 & := & \neg\varphi_1 \vee \varphi_2 \\
\forall x.\varphi & := & \neg\exists x.\neg\varphi \\
\forall X.\varphi & := & \neg\exists X.\neg\varphi \\
x \leq y & := & \forall X.(y \in X \wedge \forall z.\forall z'.(z \in X \wedge S(z', z) \rightarrow z' \in X) \rightarrow X(x)) \\
x < y & := & x \leq y \wedge \neg(y \leq x) \\
\mathit{first}(x) & := & \neg\exists y.S(y, x) \\
\mathit{last}(x) & := & \neg\exists y.S(x, y) \\
X \subseteq Y & := & \forall x.(x \in X \rightarrow x \in Y) \\
X = Y & := & X \subseteq Y \wedge Y \subseteq X \\
X = \emptyset & := & \forall Z, X \subseteq Z \\
\mathit{sing}(X) & := & X \neq \emptyset \wedge \forall Y.(Y \subseteq X \rightarrow (X \subseteq Y \vee Y = \emptyset))
\end{array}
$$

WS1S on Words

- Let Σ be a finite set of alphabet
- A word is defined as $w = w_0 w_1 ... w_{n\text{-}1}$
- A unary predicate P_a is defined for every $a \in \Sigma$ such that $P_a(i)$ if and only if $w_i = a$
- Domain of w: $dom(w) = \{0, ..., n - 1\}$
- Signature of w: $\langle dom(w), S^w, (P_a)_{a \in \Sigma} \rangle$
- Büchi Theorem: a language $L \subseteq \Sigma^*$ is regular if and only if L is expressible in WS1S

Signatures of Words

- Given an alphabet $\Sigma = \{a, b\}$, the signature of $w = abba$ is $\langle \{0, 1, 2, 3\}, S^w, P_a, P_b \rangle$ with the following interpretation
  - $S^w = \{(0, 1), (1, 2), (2, 3)\}$
  - $P_a = \{0, 3\}$
  - $P_b = \{1, 2\}$

WS1S Examples

- the last symbol is a
  - $\exists x.(P_a(x) \wedge \neg \exists y.(x < y))$
- contains substring ab
  - $\exists x.\exists y.(P_a(x) \wedge P_b(y) \wedge S(x,y))$

WS1S Examples (cont'd)

- has substring ba^*b
  - $\exists x.\exists y.(x < y \wedge P_b(x) \wedge P_b(y) \wedge \forall z((x < z \wedge z < y) \rightarrow P_a(z)))$
- non-empty word with a even length
  - $\exists f.\exists l.\exists X.(\mathit{first}(f) \wedge \mathit{last}(l) \wedge X(f) \wedge \neg X(l) \wedge \forall y.\forall z.(S(y,z) \rightarrow (X(y) \leftrightarrow \neg X(z))))$

Exercises

- Write WS1S formulas to describe the following words
  - Only a's can occur between any two occurrences of b's
  - Has an odd length (please start with $\exists$)

From NFA to WS1S

- Let $M = (Q,\ \Sigma,\ \delta,\ \{s_0\},\ F)$ be an NFA
- Assume $Q = \{s_0,\ s_1,\ ...,\ s_n\}$
- Non-empty accepting words will satisfy the following formula

$$\begin{array}{lll}\exists X_0 \ldots X_n. \quad (& \wedge_{i \neq j} & \forall .x \neg (x \in X_i \wedge x \in X_j) \\ & \wedge & \forall x.(\mathit{first}(x) \rightarrow x \in X_0) \\ & \wedge & \forall x.\forall y.(S(x,y) \rightarrow \vee_{(s_i,a,s_j)\in\delta}(x \in X_i \wedge x \in P_a \wedge y \in X_j)) \\ & \wedge & \forall x.(\mathit{last}(x) \rightarrow \vee_{(s_i,a,s_f)\in\delta; s_f\in F}(x \in X_i \wedge x \in P_a))) \end{array}$$

A Better Encoding

- Assume $|\Sigma| = 2^m$
- A symbol is binary encoded as $(t_0,\ t_1,\ ...,\ t_{m\text{-}1})$
- A word is defined as $w = w_0 w_1 ... w_{n\text{-}1}$
- A unary predicate P_i is defined for every $i \in \{0,...,m\text{-}1\}$ such that $P_i(j)$ if and only if the i-th track of w_j is 1
- Example:
  - $m = 2,\ \Sigma = \{a,\ b,\ c,\ d\},\ a = (00),\ b = (01),\ c = (10),\ d = (11)$
  - $P_0 = \{0,\ 3,\ 4\},\ P_1 = \{1,\ 4\}$
  - $w = (10)(01)(00)(10)(11) = cbacd$

Non-regular Languages

- Examples of non-regular languages:
  - $\{\ a^n b^n \mid n \in \mathbb{N}\ \}$
  - $\{\ w\#w \mid w \in \{a,\ b\}^*\ \}$
- How to prove that a language is non-regular?

Pumping Lemma

- If L is a regular language, then there is a number $p \geq 1$ (the pumping length) such that, if s is any string in L and $|s| \geq p$, then s may be divided as $s = xyz$ satisfying
  - for each $i \geq 0$, $xy^iz \in L$,
  - $|y| > 0$, and
  - $|xy| \leq p$.

Pumping Lemma Example

- Let's show that $L = \{ a^n b^n \mid n \in \mathbb{N} \}$ is non-regular
- Assume L is regular and let $w = a^p b^p$
- By pumping lemma, there are x, y, and z such that $w = xyz$,
  - $xy^i z \in L$ for each $i \geq 0$,
  - $|y| > 0$, and
  - $|xy| \leq p$
- With $|xy| \leq p$, we know that y contains only a
- But $xy^2 z \notin L$

Formal Languages

Chomsky Hierarchy	Grammar	Language	Computation Model
Type-0	Unrestricted	Recursively enumerable	Turing machine
Type-1	Context-sensitive	Context-sensitive	Linear-bounded
Type-2	Context-free	Context-free	Pushdown
Type-3	Regular	Regular	Finite

the list of formal languages in this table is not complete

Transducers

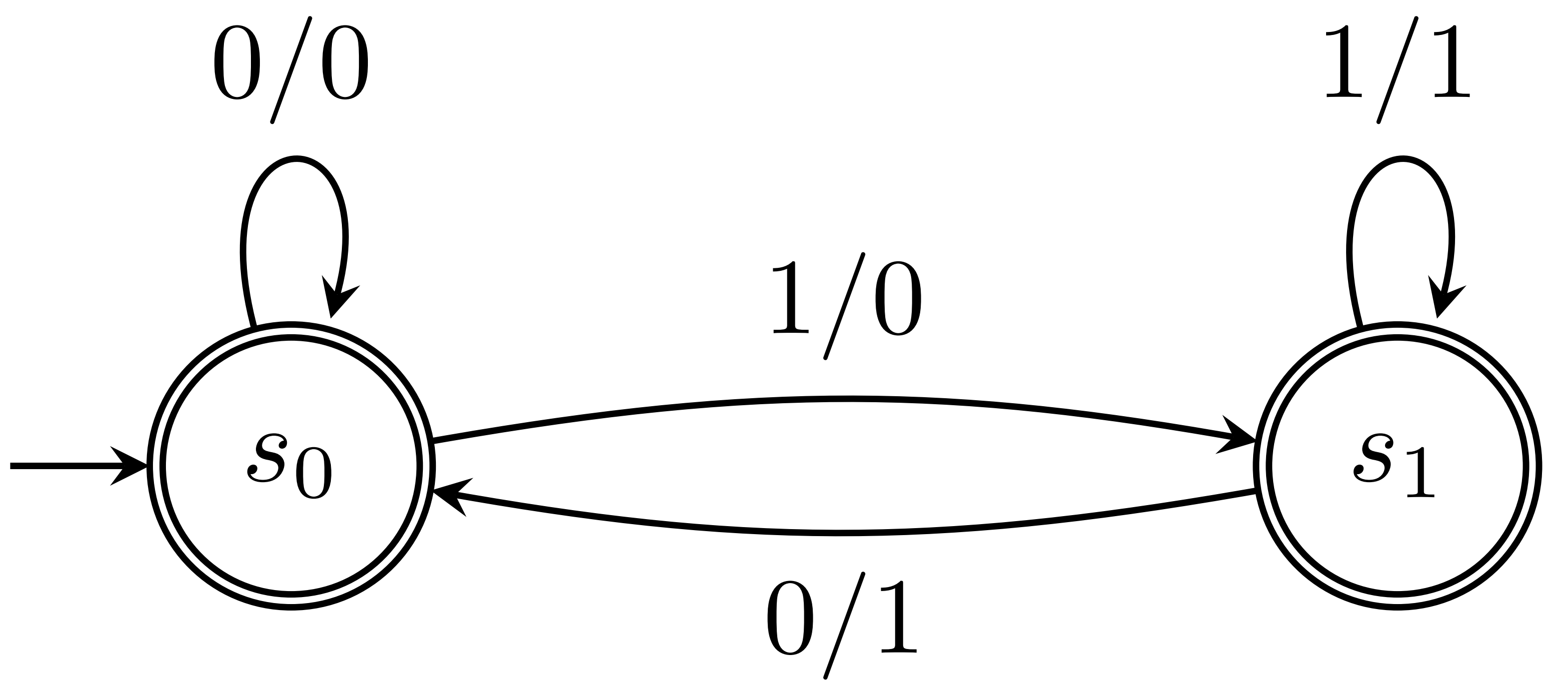
- *Finite state transducers* (*FST*)
  - Finite state automata with outputs
  - Model the relation between inputs and outputs

Formal Syntax of FST

- A finite state transducer is a 6-tuple $(Q, \Sigma, \Gamma, \delta, I, F)$ where
  - Q is a finite set of states,
  - Σ is a finite input alphabet,
  - Γ is a finite output alphabet,
  - $\delta : Q \times (\Sigma \cup \{\varepsilon\}) \times (\Gamma \cup \{\varepsilon\}) \rightarrow 2^Q$ is the transition function (sometimes written as a relation $\delta : Q \times (\Sigma \cup \{\varepsilon\}) \times (\Gamma \cup \{\varepsilon\}) \times Q$),
  - $I \subseteq Q$ is the set of initial states, and
  - $F \subseteq Q$ is the set of accepting (final) states

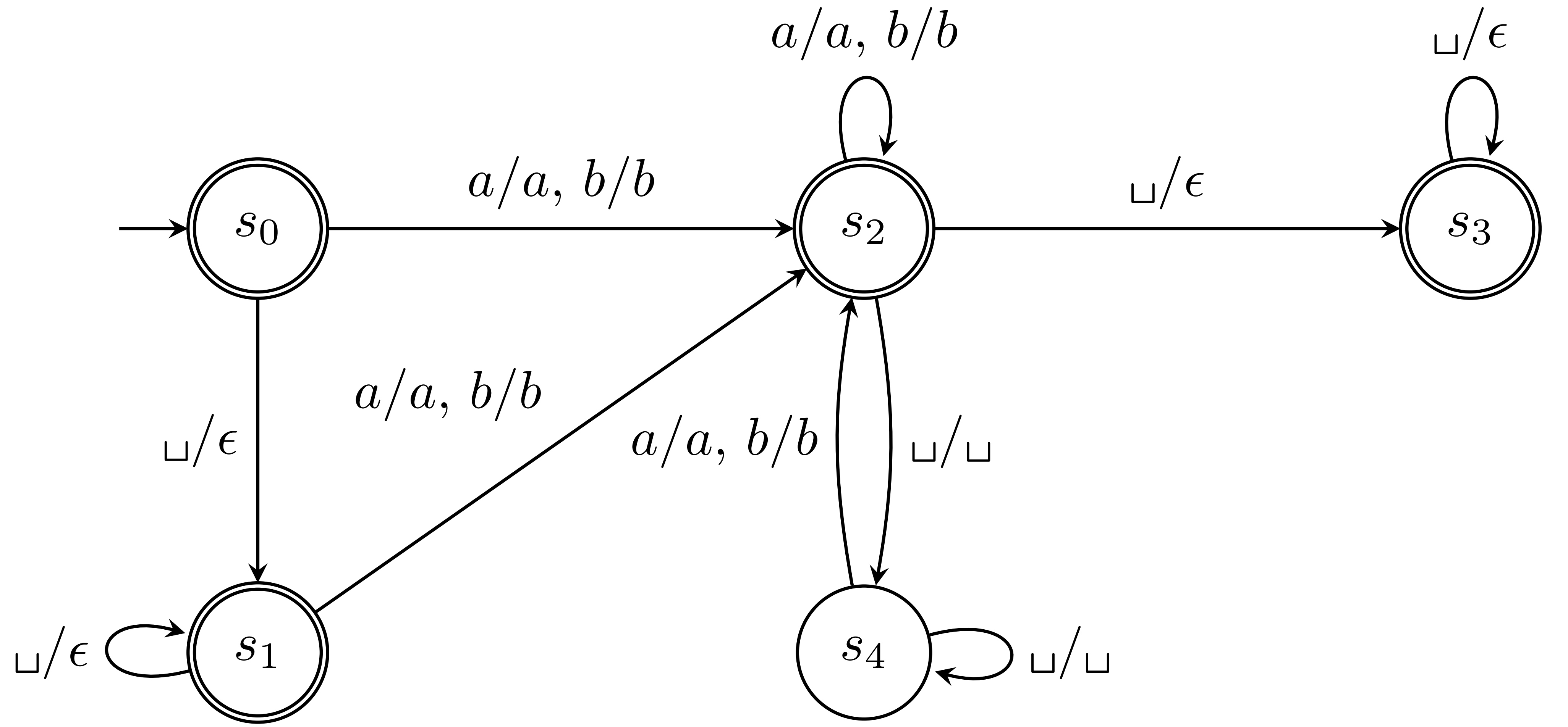
Example: Divide by 2

- Assume the alphabet is {0, 1}



Example: Trim

- Assume the alphabet is {a, b, ␣}, where ␣ denotes a space



Tools

- MONA (http://www.brics.dk/mona/)
- JFLAP (http://www.jflap.org)
- GOAL (http://goal.im.ntu.edu.tw/wiki/doku.php)

Infinite Computations

- A *reactive system* is a system that continuously interacts with its environment
- Computations of a reactive system are infinite
- How to model such infinite computations?
  - Automata on infinite words

Infinite Words

- Let Σ be a finite alphabet
- An infinite word *w* over Σ ($w \in \Sigma^\omega$) is a sequence of symbols $w_0 w_1 w_2 \ldots$ with $w_i \in \Sigma$
  - Length of w is $\boldsymbol{\omega}$
- Examples ($\Sigma = \{a,\ b\}$):
  - $a\ b\ (b\ a)^{\boldsymbol{\omega}}$
  - $a\ b\ a\ (b\ a\ b)^{\boldsymbol{\omega}}$

ω-Automata

Syntax

- An *ω-automaton* is a tuple $(Q, \Sigma, \delta, q_0, Acc)$ where
  - Q is a finite set of states,
  - Σ is a finite alphabet,
  - δ: $Q \times \Sigma \rightarrow 2^Q$ is the transition function,
  - q_0 is the initial state, and
  - Acc is the *acceptance condition*
- Different ω-automata can be defined by different acceptance conditions

ω-Automata
Semantics

- Let $M = (Q,\ \Sigma,\ \delta,\ q_0,\ Acc)$ be an ω-automaton
- Let $w = w_0 w_1 w_2 ...$ be an infinite word over Σ
- A *run* of w on M is a sequence of states $q_0 q_1 q_2 ...$ where $(q_i,\ w_i,\ q_{i+1}) \in \delta$

ω-Automata
Semantics (cont'd)

- A run is accepting if the run satisfies the acceptance condition Acc

- A word is accepted if there is a run of M on the word

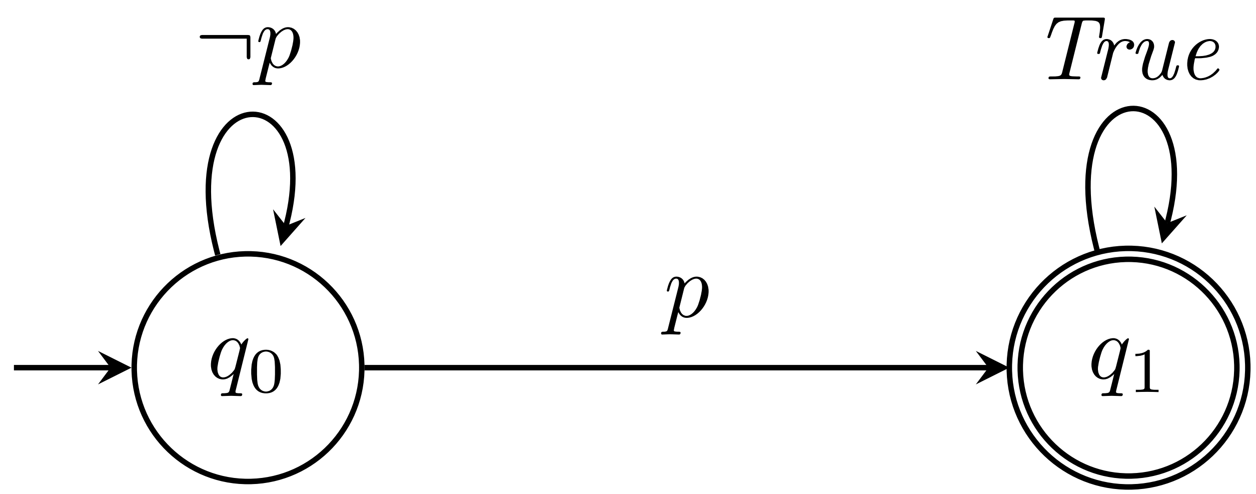
- The language of M, denoted by $L(M)$, is the set of words accepted by M

- Define $\mathit{Inf}(\rho) = \{s \mid s \textit{ occurs in } \rho \textit{ infinitely many times}\}$

Acceptance Conditions

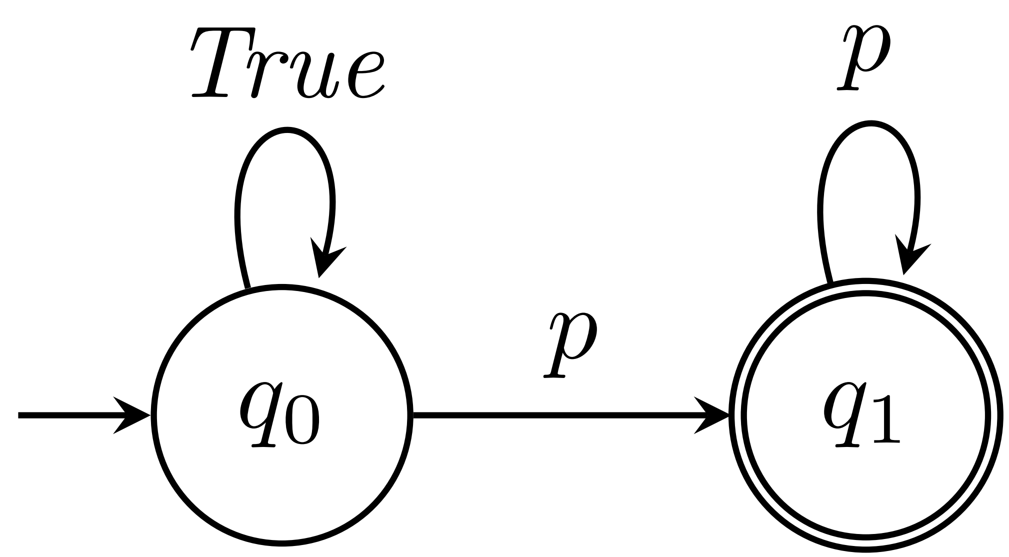
Acceptance Condition	Acc	Satisfaction	Abbrev.	Note
Büchi	$Acc = F \subseteq Q$	$Inf(\rho) \cap F \neq \varnothing$	NBW	
co-Büchi	$Acc = F \subseteq Q$	$Inf(\rho) \cap F = \varnothing$	NCW	
Generalized Büchi	$Acc = \{F_1, \ldots, F_n\}$, $F_i \subseteq Q$	$Inf(\rho) \cap F_i \neq \varnothing$ for all $F_i \in F$	NGW	
Rabin	$Acc = \{(E_1, F_1), \ldots, (E_n, F_n)\}$, $F_i \subseteq Q$, $E_i \subseteq Q$	$Inf(\rho) \cap E_i = \varnothing$ and $Inf(\rho) \cap F_i \neq \varnothing$ for some i	NRW	
Streett	$Acc = \{(E_1, F_1), \ldots, (E_n, F_n)\}$, $F_i \subseteq Q$, $E_i \subseteq Q$	$Inf(\rho) \cap F_i \neq \varnothing$ implies $Inf(\rho) \cap E_i \neq \varnothing$ for all i	NSW	
Muller	$Acc = \{F_1, \ldots, F_n\}$, $F_i \subseteq Q$	$Inf(\rho) = F_i$ for some i	NMW	
Parity	$Acc\colon Q \rightarrow \mathbb{N}$	min parity in ρ is even	NPW	$Acc(q)$ is the parity of q

Büchi Automata
Example 1



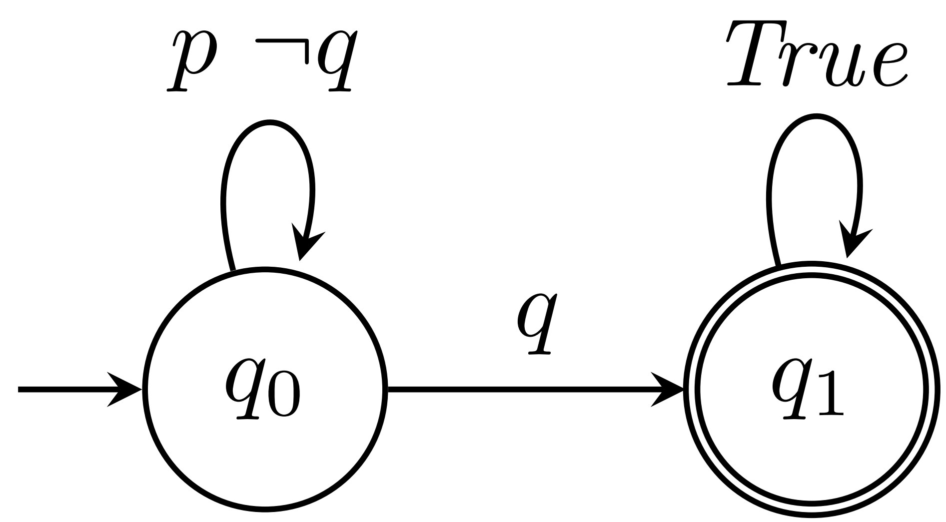
accepts infinite words where p holds eventually

Büchi Automata
Example 2



accepts infinite words where eventually p will always hold

Büchi Automata
Example 3



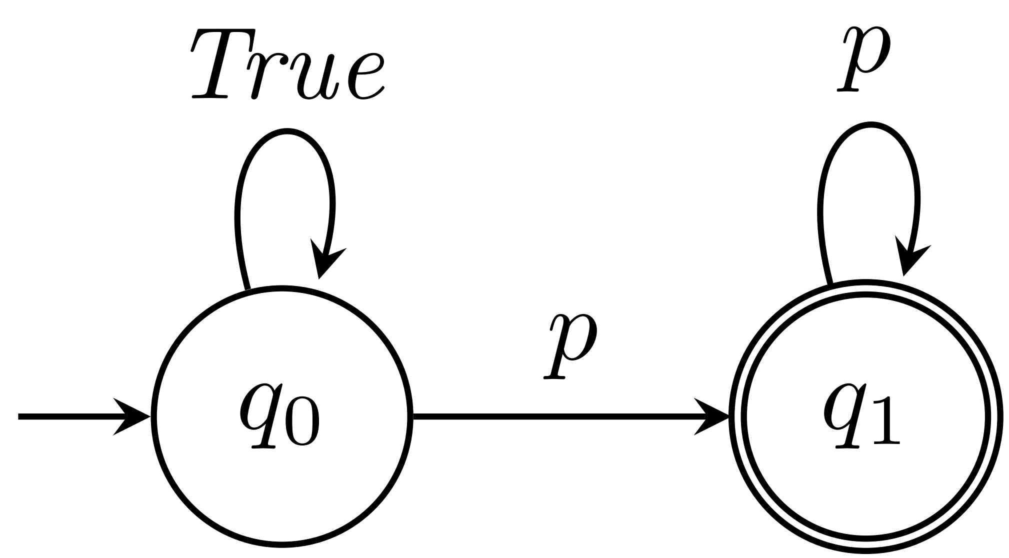
accepts infinite words where p holds until q holds

Exercise

- Draw a Büchi automaton that accepts infinite words where p holds infinitely many times. ($\Sigma = \{p, \neg p\}$)

Deterministic VS Nondeterministic

- Can you find a deterministic Büchi automaton (DBW) that accepts the same language?



Deterministic VS Nondeterministic

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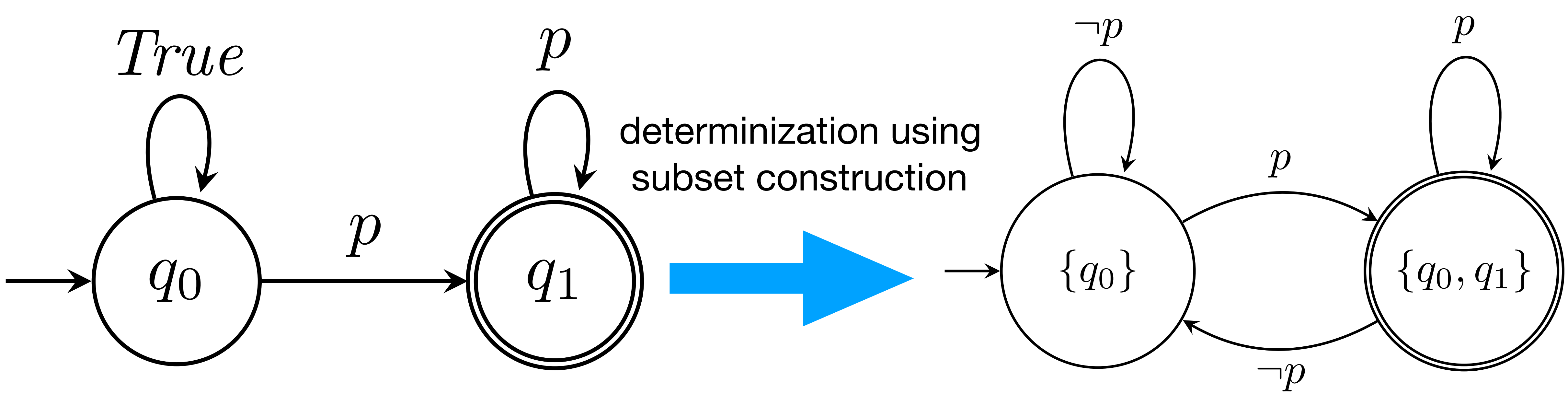

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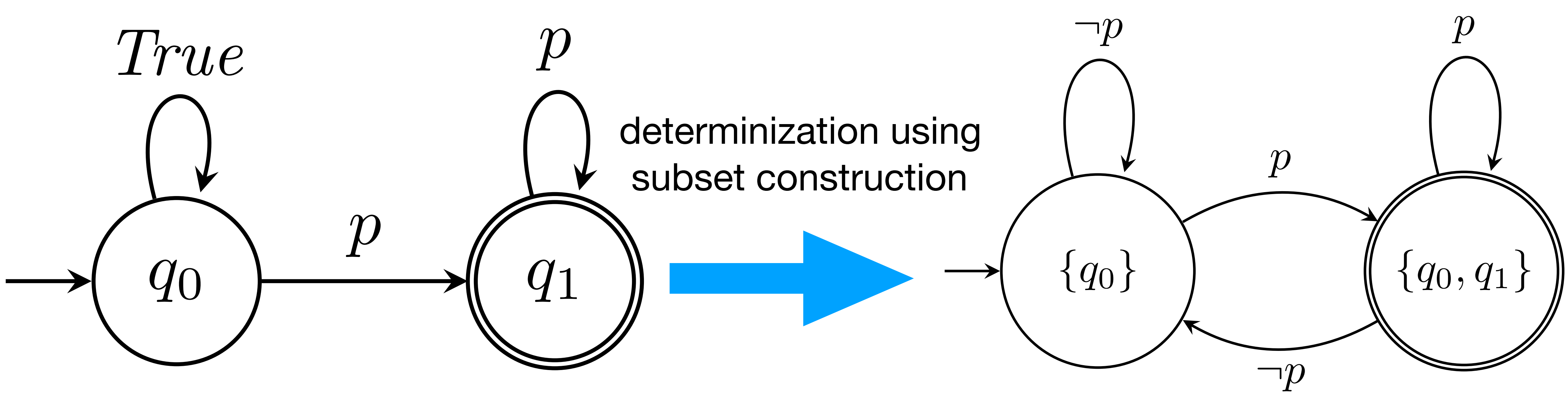


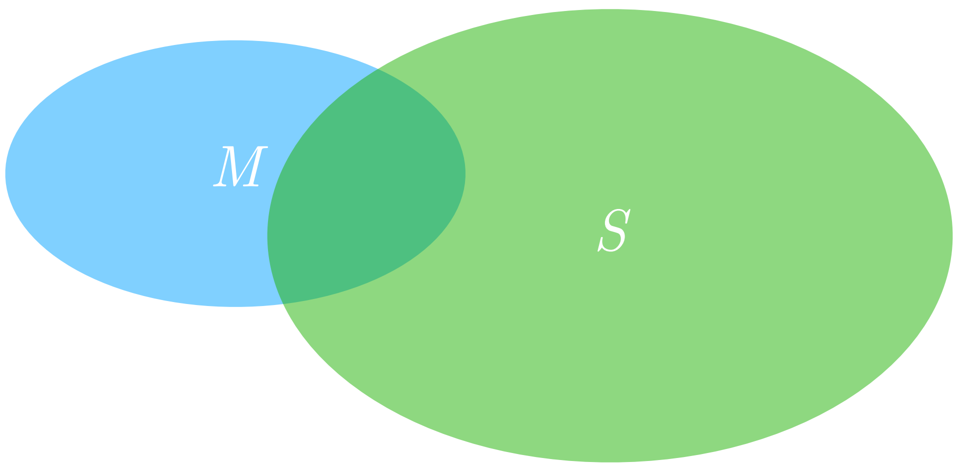
NBW is more expressive than DBW

Model VS Specification

- So far we already learnt some abstract machines as models of computations
- We may require that the computations must satisfy some properties
- How do we check?

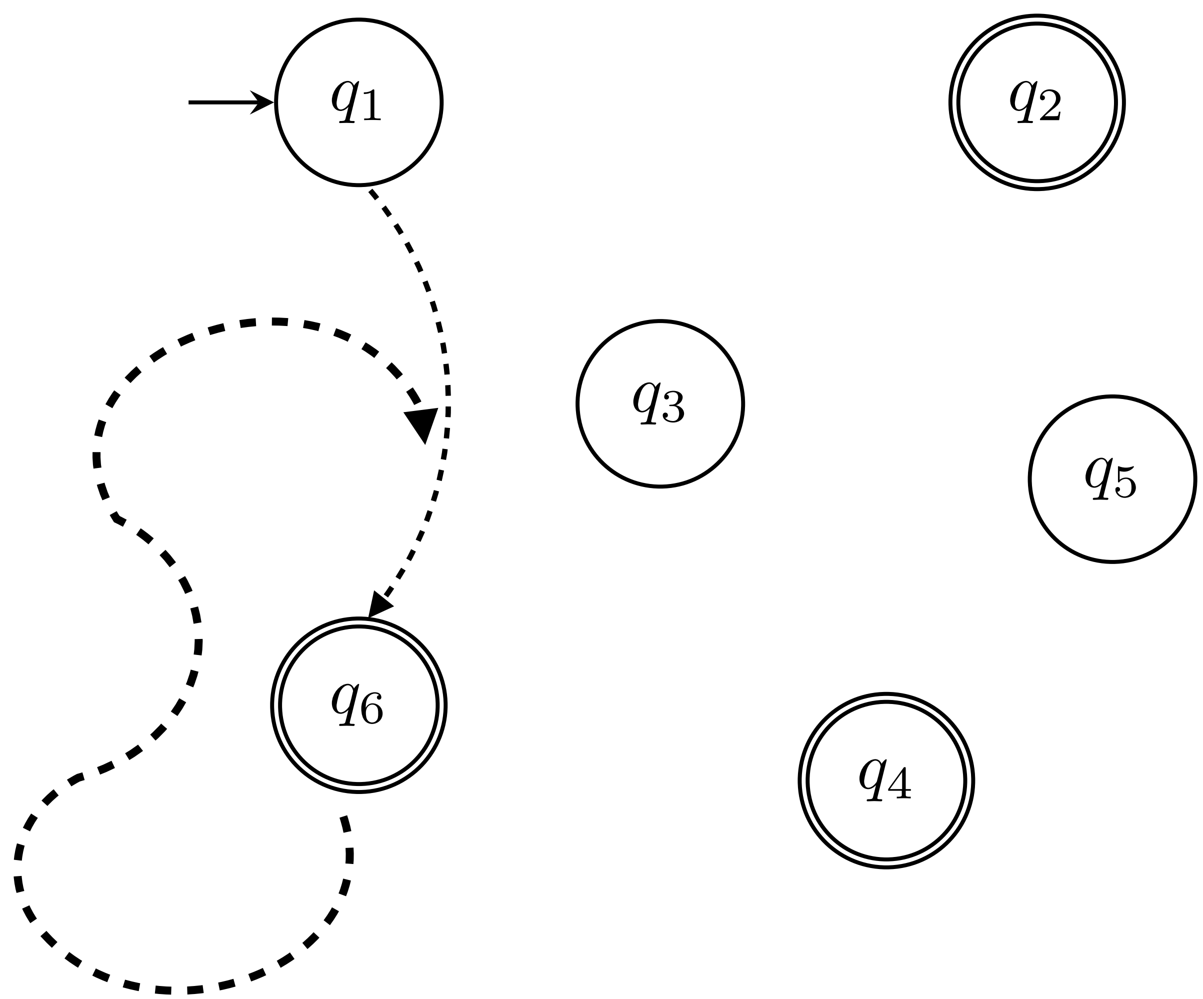
Model Checking

- Model the computations of a system as an automaton M
- Model the computations allowed by the specification as an automaton S
- Check if the system satisfies the specification by checking if $L(M) \subseteq L(S)$
- Or equivalently checking if P is empty where P is the intersection of
  - M and
  - the complement of S.



Emptiness Test

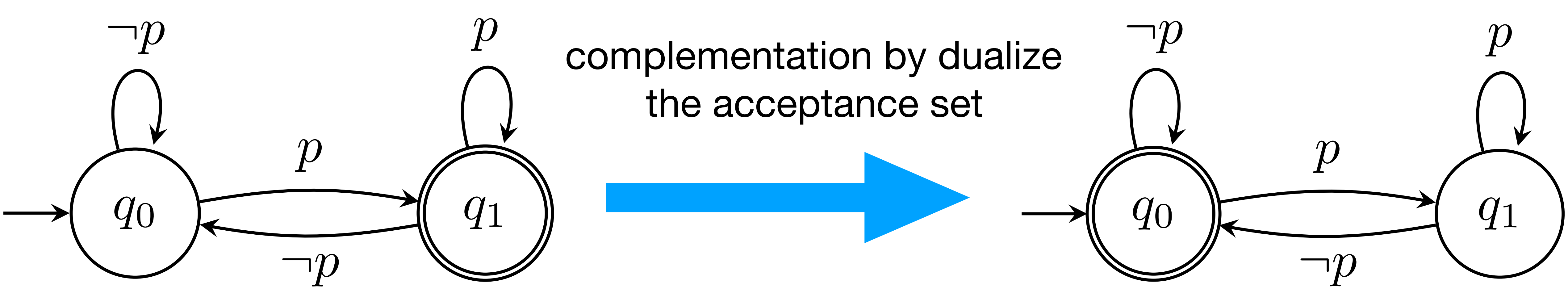
- Use double depth-first search to find an accepting lasso



Büchi Automata Intersection

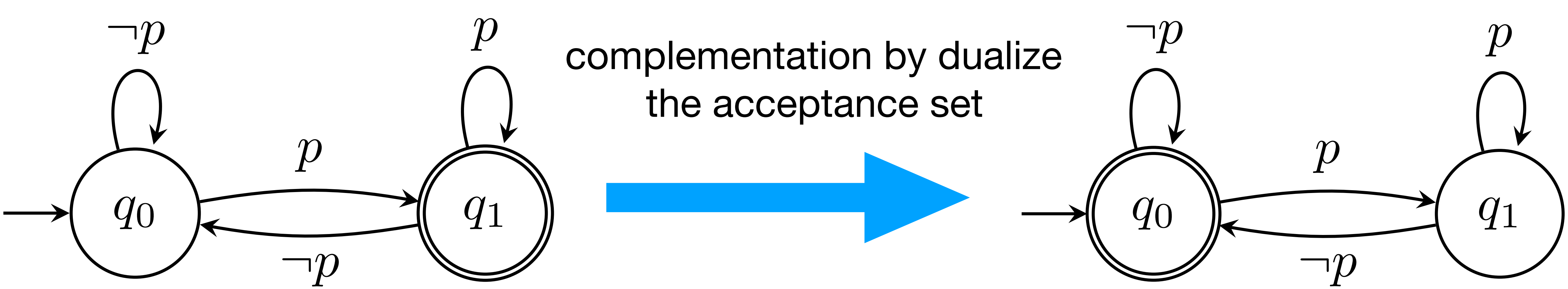
- $M_1 = (Q_1, \Sigma, \delta_1, q_{01}, F_1)$, $M_2 = (Q_2, \Sigma, \delta_2, q_{02}, F_2)$
- Construct $M = (Q_1 \times Q_2 \times \{0,1,2\}, \Sigma, \delta, (q_{01}, q_{02}, 0), Q_1 \times Q_2 \times \{0\})$ where $((q_1, q_2, i), a, (q_1', q_2', j)) \in \delta$ if
  - $(q_1, a, q_1') \in \delta_1$ and $(q_2, a, q_2') \in \delta_2$,
  - $j = 1$ if $i = 0$,
  - $j = i$ if $i \neq 0$ and $q_i \notin F_i$, and
  - $j = (i + 1)$ mod 2 if $i \neq 0$ and $q_i \in F_i$
- $L(M) = L(M_1) \cap L(M_2)$

Büchi Automata Complementation



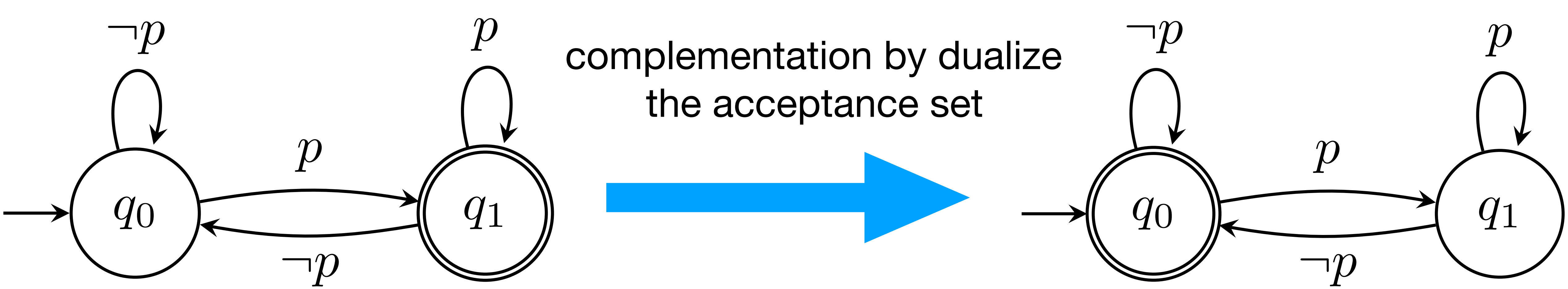
Does the right one exactly accept the complement of the left one?

Büchi Automata Complementation



Does the right one exactly accept the complement of the left one? ✗

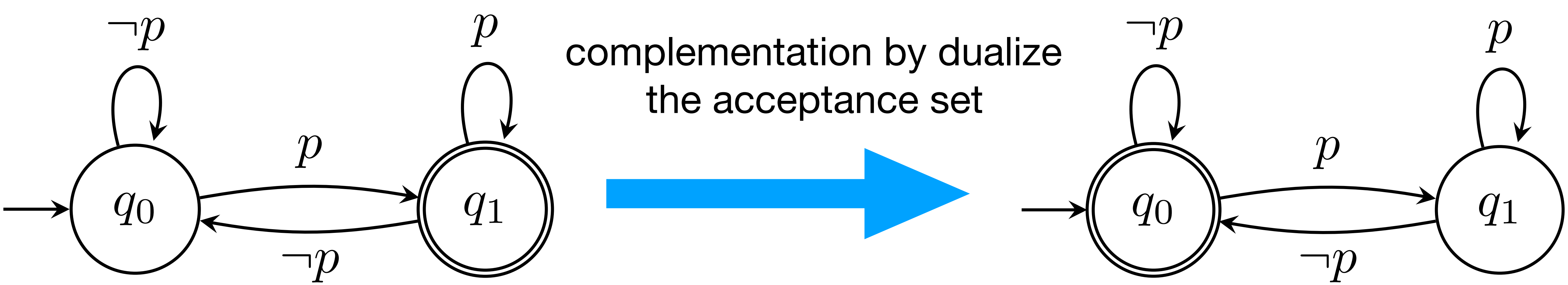
Büchi Automata Complementation



Does the right one exactly accept the complement of the left one? ✗

Complementation of NBW is much harder than that of NFA

Büchi Automata Complementation



Does the right one exactly accept the complement of the left one? ✗

Complementation of NBW is much harder than that of NFA

We may express specifications using logic formulas

LTL Model Checking

- Express the behavior of a system as a Büchi automaton *M* (usually converted from a Kripke structure)
- Express the specification as a formula f in *linear temporal logic* (LTL)
- Translation $\neg f$ to a Büchi automaton $A_{\neg f}$ *with labels on states*
- Check if $L(M) \cap L(A_{\neg f})$ is empty

Linear Temporal Logic
Syntax

- AP is a finite set of atomic propositions
- The alphabet Σ is defined as 2^{AP}
- A linear temporal logic (LTL) formula is defined as follows
  - For every $p \in AP$, p is an LTL formula
  - If f and g are LTL formulas, then so are $\neg f$, $f \wedge g$, $\boldsymbol{X} f$, and f $\boldsymbol{U}$ g
- $\boldsymbol{X}$ and $\boldsymbol{U}$ are (future) temporal operators

Linear Temporal Logic Semantics

- A state is a subset of AP, containing exactly those propositions that evaluate to true in that state

- An LTL formula is interpreted over an infinite sequence of states $\rho = s_0 s_1 \ldots$.

$$(\rho,\ i) \models p \quad \text{iff} \quad p \in s_i$$

$$(\rho,\ i) \models \neg\, f \quad \text{iff} \quad (\rho,\ i) \not\models f$$

$$(\rho,\ i) \models f \wedge g \quad \text{iff} \quad (\rho,\ i) \models f \text{ and } (\rho,\ i) \models g$$

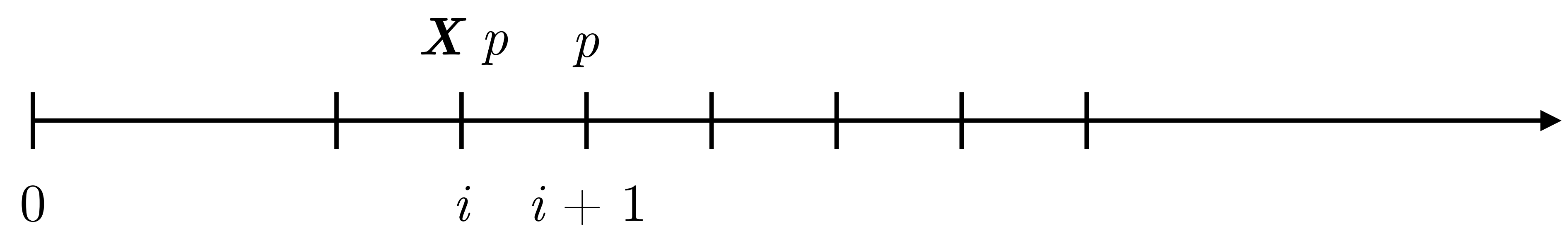
$$(\rho,\ i) \models \boldsymbol{X}\, f \quad \text{iff} \quad (\rho,\ i+1) \models f$$

$$(\rho,\ i) \models f\ \boldsymbol{U}\ g \quad \text{iff} \quad \text{exists } j \geq i \text{ such that } (\rho,\ j) \models g \text{ and}$$

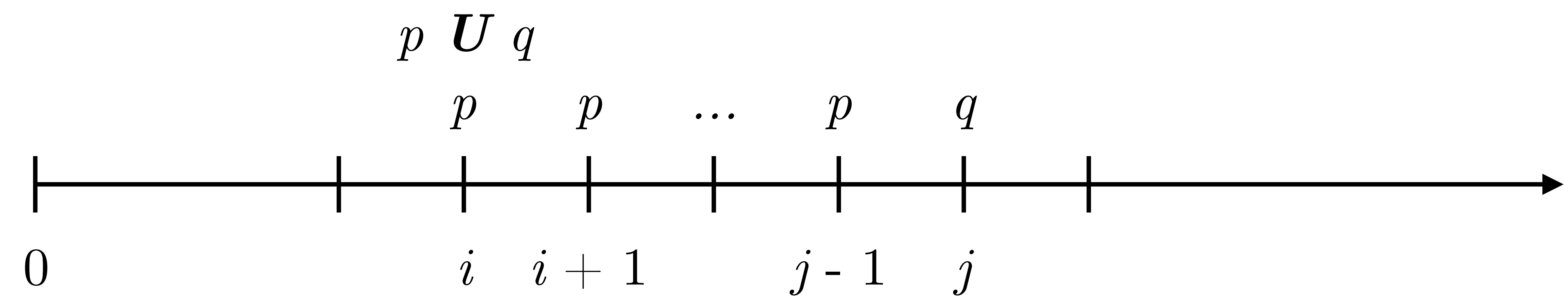
$$\text{for all } i \leq k < j,\ (\rho,\ k) \models f$$

Next and Until

- $(\rho,\ i) \vDash \boldsymbol{X}\, f$ iff $(\rho,\ i+1) \vDash f$

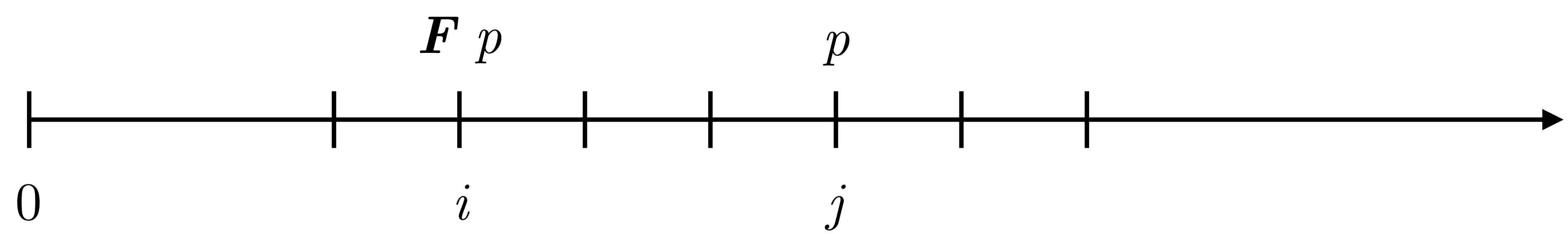


- $(\rho,\ i) \vDash f\ \boldsymbol{U}\ g$ iff exists $j \geq i$ such that $(\rho,\ j) \vDash g$ and for all $i \leq k < j$, $(\rho,\ k) \vDash f$

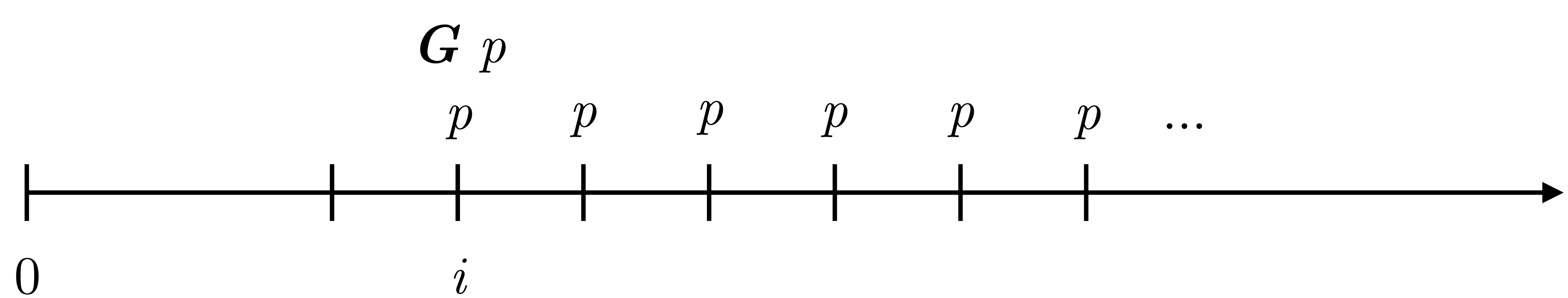


Future and Global

- $(\rho,\ i) \vDash \boldsymbol{F}\ f$ iff $(\rho,\ j) \vDash f$ for some $j \geq i$

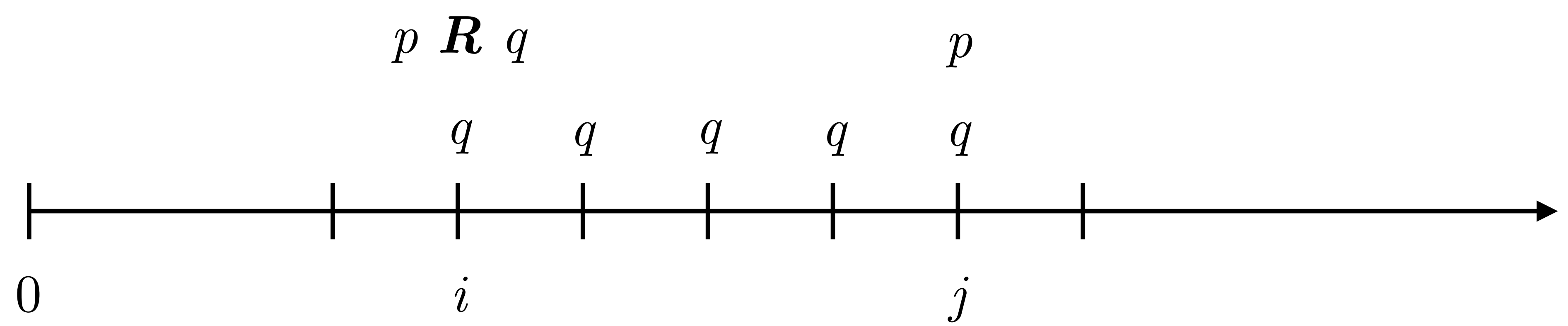


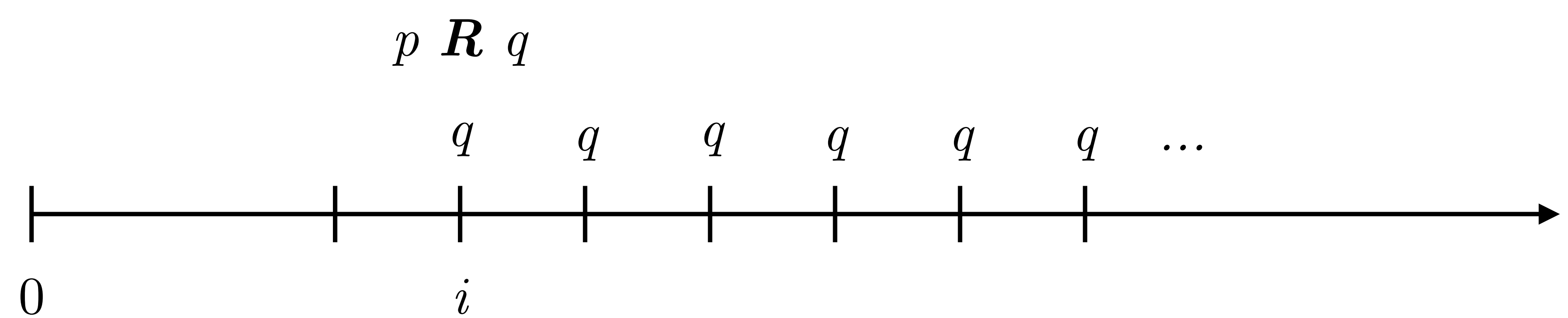
- $(\rho,\ i) \vDash \boldsymbol{G}\ f$ iff $(\boldsymbol{\rho},\ j) \vDash f$ for all $j \geq i$



Release

- $(\rho, i) \vDash f \, \boldsymbol{R} \, g$ iff exists $j \geq i$ such that $(\rho, j) \vDash f$ and for all $i \leq k \leq j$, $(\rho, k) \vDash g$; or for all $j \geq i$, $(\rho, j) \vDash g$





Abbreviations

- $true := p \vee \neg p$
- $false := \neg true$
- $f \vee g := \neg(\neg f \wedge \neg g)$
- $f \rightarrow g := \neg f \vee g$
- $f \leftrightarrow g := (f \rightarrow g) \wedge (g \rightarrow f)$
- $f\ \boldsymbol{R}\ g := \neg(\neg f\ \boldsymbol{U}\ \neg g)$
- $\boldsymbol{F}\ g := true\ \boldsymbol{U}\ g$
- $\boldsymbol{G}\ f := false\ \boldsymbol{R}\ f$

$$\bigcirc = \boldsymbol{X},\ \Diamond = \boldsymbol{F},\ \Box = \boldsymbol{G}$$

Exercise

- Express the following sentences in LTL formulas.
  - “p occurs infinitely often”
  - “whenever a message is sent, eventually an acknowledgement will be received”

Satisfaction, Validity, and Congruence

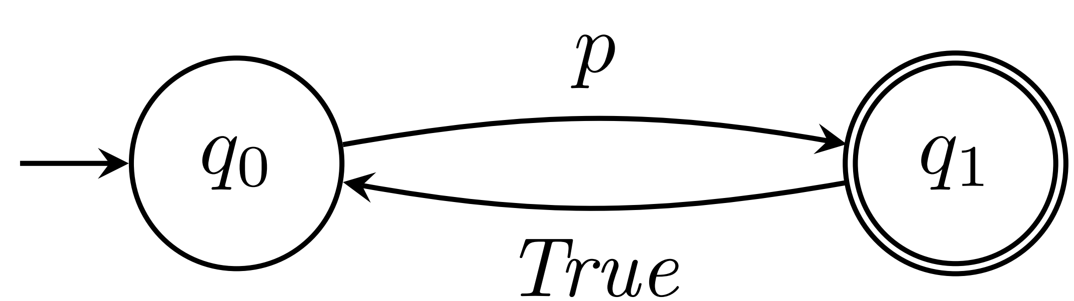
- $\rho \models f$: a state sequence ρ *satisfies* an LTL formula f
  - $\rho \models f$ iff $(\rho, 0) \models f$
- $\models f$: an LTL formula f is *valid*
  - $\models f$ iff $\rho \models f$ for all ρ
- $f \cong g$: two formulas f and g are *congruent*
  - $f \cong g$ iff $\models \boldsymbol{G}\ (f \leftrightarrow g)$

Congruent Formulas

- $\neg\ \boldsymbol{X}\ f \cong \boldsymbol{X}\ \neg\ f$
- $\neg\ \boldsymbol{F}\ g \cong \boldsymbol{G}\ \neg\ g$
- $\neg\ \boldsymbol{G}\ f \cong \boldsymbol{F}\ \neg\ f$
- $\boldsymbol{G}\ \boldsymbol{G}\ f \cong \boldsymbol{G}\ f$
- $\boldsymbol{F}\ \boldsymbol{F}\ g \cong \boldsymbol{F}\ g$
- $\neg\ \neg\ f \cong f$

Expressive Power of LTL

- LTL is strictly less expressive than NBW
- "even p" can be expressed in NBW but not LTL



- NBW is as expressive as QPTL (Quantified Propositional Temporal Logic)
- "even p" in QPTL: $\exists\ t.\ t \wedge \boldsymbol{G}\ (t \leftrightarrow \boldsymbol{X}\ \neg t) \wedge \boldsymbol{G}\ (t \rightarrow\ p)$

From LTL to Labeled NGW

- Translate an LTL formula *f* to a labeled NGW (with labels on states)
  - Take the *negation normal form* (NNF) of f
  - Expand f_{NNF} into basic formulas as the initial states
  - Construct successors of states based on $\boldsymbol{X}$-formulas
  - For each subformula $g\ \boldsymbol{U}\ h$, create an acceptance set such that h will become true eventually

NNF: negation only occurs right before propositions

Basic Formulas

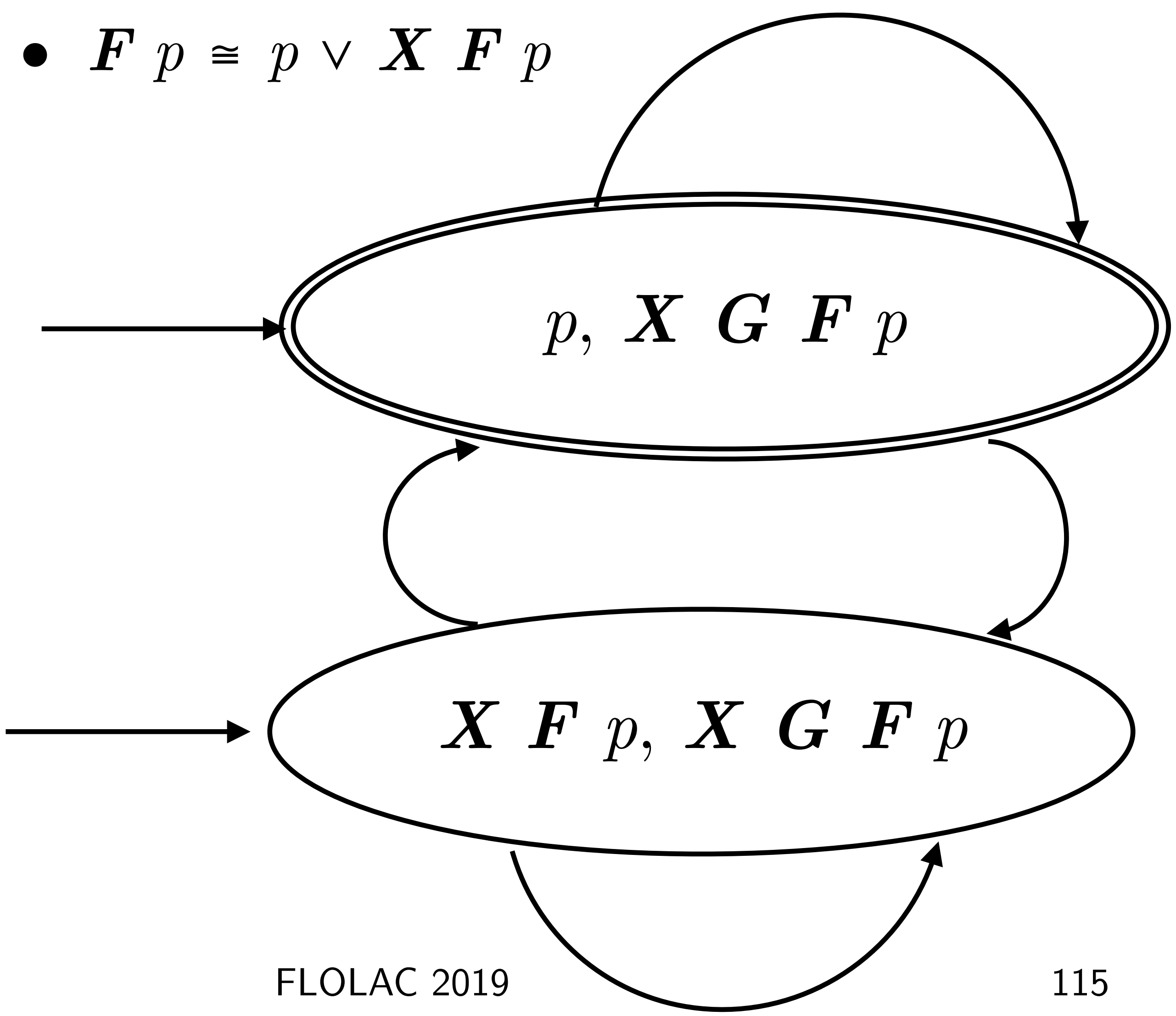
- A *literal* is either a proposition or its negation
- A basic formula is either a literal or an $\boldsymbol{X}$-formula

Expansion Formulas

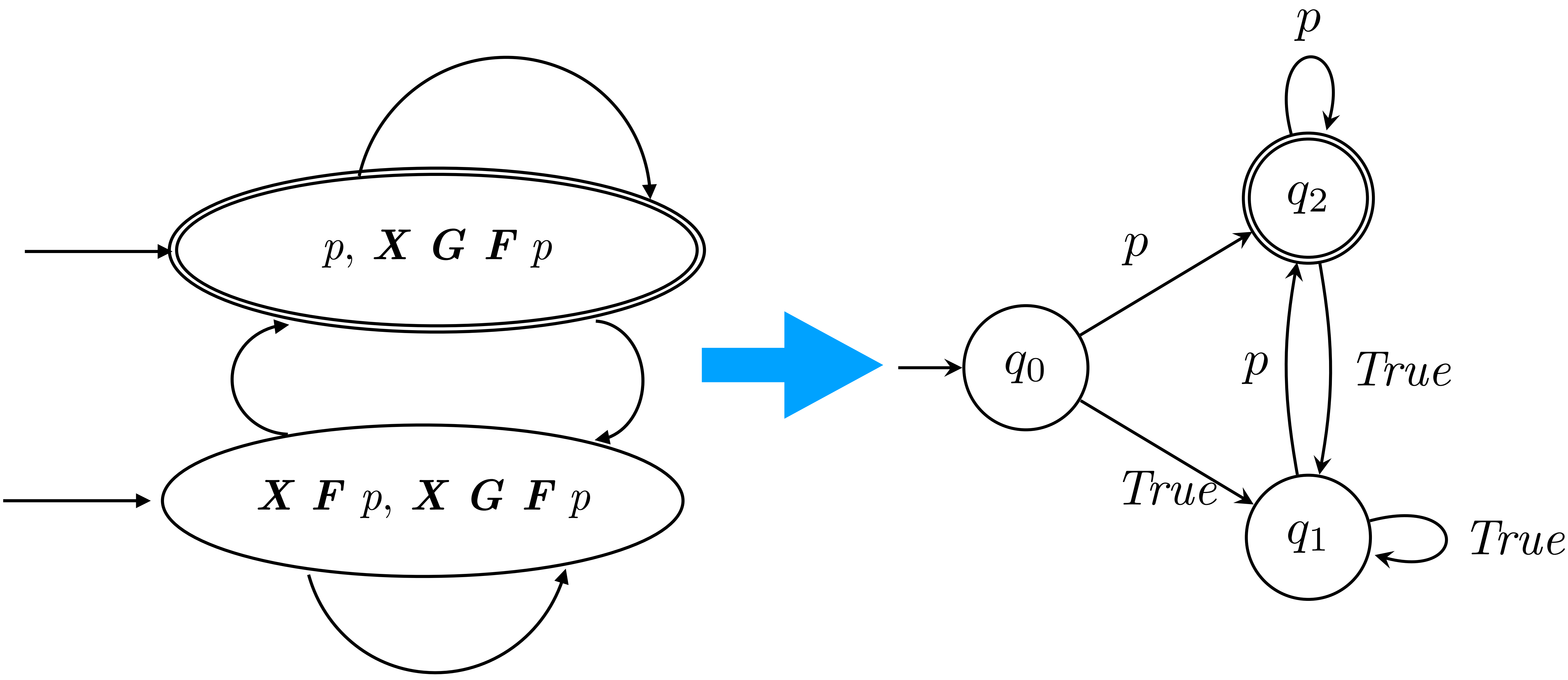
- $\boldsymbol{F}\, g \cong g \vee \boldsymbol{X}\ \boldsymbol{F}\, g$
- $\boldsymbol{G}\, f \cong f \wedge \boldsymbol{X}\ \boldsymbol{G}\, f$
- $f\ \boldsymbol{U}\ g \cong g \vee (f \wedge \boldsymbol{X}\ (f\ \boldsymbol{U}\ g))$
- $f\ \boldsymbol{R}\ g \cong g \wedge (f \vee \boldsymbol{X}\ (f\ \boldsymbol{R}\ g))$

From LTL to Labeled NGW Example

- $f := \boldsymbol{G}\ \boldsymbol{F}\ p$
- $\boldsymbol{G}\ \boldsymbol{F}\ p \cong (p \vee \boldsymbol{X}\ \boldsymbol{F}\ p) \wedge \boldsymbol{X}\ \boldsymbol{G}\ \boldsymbol{F}\ p \cong (p \wedge \boldsymbol{X}\ \boldsymbol{G}\ \boldsymbol{F}\ p) \vee (\boldsymbol{X}\ \boldsymbol{F}\ p \wedge \boldsymbol{X}\ \boldsymbol{G}\ \boldsymbol{F}\ p)$
- $\boldsymbol{F}\ p \cong p \vee \boldsymbol{X}\ \boldsymbol{F}\ p$



From Labeled NGW to NGW



From NGW to NBW

- Apply the same technique in the intersection of NBW
- Use an index i to remember the next acceptance set in $\{F_1,\ F_2,$ ..., $F_n\}$ to be passed
- Once a state in F_i is passed, increase the index i by 1
- If every $F_i \in \{F_1,\ F_2,\ ...,\ F_\mathrm{n}\}$ has been passed at least once, change the index to 0 and set the index to 1 in the successors
- A run is accepting if the index 0 is passed infinitely many times

Tools

- LTL2BA (http://www.lsv.fr/~gastin/ltl2ba/index.php)
- LTL3BA (https://sourceforge.net/projects/ltl3ba/)
- SPIN (http://spinroot.com/spin/whatispin.html)
- NuSMV (http://nusmv.fbk.eu)
- GOAL (http://goal.im.ntu.edu.tw/wiki/doku.php)