

Program Construction and Reasoning Exercises for Day 2

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1 In-Class Exercises

1.1 Folds and Fold-Fusion

1. Given functions $f :: \alpha \rightarrow \beta$ and $g :: \alpha \rightarrow \gamma$, $split\ f\ g :: \alpha \rightarrow (\beta, \gamma)$ is a function defined by:

$$split\ f\ g\ a = (f\ a, g\ a).$$

Recall the definition of *steep* and *sum*. The definition of *steepsum* can be re-written as:

$$steepsum = split\ steep\ sum.$$

Also recall that the identity function *id* on lists is a fold: $id = foldr\ (\cdot)\ []$. Use the fold-fusion theorem to fuse $steepsum \cdot id$ into one fold.

2. Recall the definition of *scanr* from the lecture:

$$scanr\ f\ e = map\ (foldr\ f\ e) \cdot tails$$

and its implementation as a fold:

$$scanr\ f\ e = foldr\ (sc\ f)\ [e] \\ \mathbf{where}\ sc\ f\ x\ (y:ys) = f\ x\ y : y : ys$$

- (a) Expand $scanr\ (+)\ 0\ [1, 2, 3]$ step by step:

$$scanr\ (+)\ 0\ [1, 2, 3] \\ = foldr\ (sc\ (+))\ [0]\ [1, 2, 3] \\ = \dots$$

- (b) Derive the implementation of $scanr\ f\ e$ by fusing $map\ (foldr\ f\ e) \cdot tails$ into one fold.

2 Take-Home Exercise (Due Date: July 10th)

You do not have to do the exercises below if you have completed any of the exercises from Day 1. Exercise 1 is worth 40 points while exercise 2 is worth 50 points.

1. The function *filter p* selects from a list all elements satisfying a predicate *p*. For example, *filter even* [1, 2, 3, 4] = [2, 4].

- (a) Give a recursive definition of *filter*:

$$\begin{aligned} \textit{filter } p [] &= \dots \\ \textit{filter } p (x:xs) &= \dots \end{aligned}$$

- (b) Define *filter p* in terms of *foldr*.

- (c) Prove, by fold-fusion, that

$$\textit{filter } p \cdot \textit{map } f = \textit{map } f \cdot \textit{filter } (p \cdot f).$$

Hint: apply fold-fusion on both sides, and show that they are equal to the same fold.

2. Given two functions *h*₁ and *h*₂, the function *split h*₁ *h*₂ computes the pair of their results:

$$\textit{split } h_1 h_2 xs = (h_1 xs, h_2 xs).$$

In the special case when both *h*₁ and *h*₂ are defined by *foldr*:

$$\begin{aligned} h_1 &= \textit{foldr } f_1 e_1, \\ h_2 &= \textit{foldr } f_2 e_2, \end{aligned}$$

the following “banana-split” rule allows us to express *split h*₁ *h*₂ using one single *foldr*:

$$\begin{aligned} \textit{split } h_1 h_2 &= \textit{foldr } g (e_1, e_2), \\ \mathbf{where } g x (y, z) &= (f_1 x y, f_2 x z). \end{aligned}$$

It optimises two traversal through the list to only one traversal. It is called “banana-split” because folds used to be written using a notation called “banana brackets”.

- (a) The function *split sum length* return the pair of sum and length of the input list. Use the banana-split rule to express *split sum length* by a fold.
- (b) Prove the banana-split rule by fold fusion. Hint: recall that *split h*₁ *h*₂ = *split h*₁ *h*₂ · *id*, and *id* is a fold.