

Functional Program Derivation

Exercises for Day 1

Shin-Cheng Mu Shu-Chun Weng (TA)

2007 Formosan Summer School on
Logic, Language, and Computation
July 4th, 2007

This exam sheet is worth 20 points in total.

1.1 The Expand/Reduce Transformation

1. (1 point) What does this function do?

$$\begin{aligned} \mathit{descend} 0 &= [] \\ \mathit{descend} (n + 1) &= (n + 1) : \mathit{descend} n \end{aligned}$$

2. (1 point) Consider the definition $f = \mathit{sum} \cdot \mathit{descend}$

- (a) Describe in words what this function does.
- (b) Calculate $f 0$.
- (c) Simplify $f (n + 1) = \dots f n \dots$
- (d) From (b) and (c), synthesise a recursive definition of f .

3. (1 point) Recall the datatype definition for internally labelled binary trees:

$$\mathit{data} \mathit{iTree} \alpha = \mathit{Null} \mid \mathit{Node} \alpha (\mathit{iTree} \alpha) (\mathit{iTree} \alpha).$$

Consider the function $\mathit{mapiTree}$ defined below:

$$\begin{aligned} \mathit{mapiTree} f \mathit{Null} &= \mathit{Null}, \\ \mathit{mapiTree} f (\mathit{Node} a t u) &= \mathit{Node} (f a) (\mathit{mapiTree} f u) (\mathit{mapiTree} f t). \end{aligned}$$

What does this function do?

4. (1 point) Define a function $\mathit{sumiTree}$ computing the sum of all node values in an iTree .

5. **(2 points)** The function $one\ x = 1$ returns 1, what ever the input is. The function $sizeiTree$, computing the size of a tree, can be specified by:

$$sizeiTree = sumiTree \cdot mapiTree\ one.$$

Derive a definition of $sizeiTree$ which does not construct an intermediate tree.

1.2 Proof by Induction

1. **(2 points)** Prove $(xs \ ++\ ys) \ ++\ zs = xs \ ++\ (ys \ ++\ zs)$. Hint: induction on xs .
2. **(2 points)** The function $concat$ concatenates a list of lists:

$$\begin{aligned} concat\ [] &= [], \\ concat\ (xs : xss) &= xs \ ++\ concat\ xss. \end{aligned}$$

E.g. $concat\ [[1, 2], [3, 4], [5]] = [1, 2, 3, 4, 5]$. Prove that:

$$sum \cdot concat = sum \cdot map\ sum.$$

Hint: you may need one of the properties proved in the lecture.

3. **(2 points)** Prove that $map\ f \cdot map\ g = map\ (f \cdot g)$.
4. **(2 points)** The function $swapTree$ is defined by:

$$\begin{aligned} swapiTree\ Null &= Null, \\ swapiTree\ (Node\ a\ t\ u) &= Node\ a\ (swapiTree\ u)\ (swapiTree\ t). \end{aligned}$$

Prove that $swapiTree\ (swapiTree\ t) = t$ for all t .

1.3 Accumulating Parameters

1. **(3 points)** Recall the standard definition of factorial:

$$\begin{aligned} fact\ 0 &= 1, \\ fact\ (n + 1) &= (n + 1) \times fact\ n. \end{aligned}$$

This program also implicitly uses space linear to n in the call stack.

- (a) Introduce $factit\ n\ m = \dots$ where m is an accumulating parameter.
- (b) Express $fact$ in terms of $factit$.
- (c) Construct a space efficient implementation of $factit$.

2. (**3 points**) Given an *iTree*, the following function *flatten* returns a list of all labels in the tree, in left-to-right order:

$$\begin{aligned} \textit{flatten} \textit{Null} &= [], \\ \textit{flatten} (\textit{Node} \ a \ t \ u) &= \textit{flatten} \ t \ ++ \ [a] \ ++ \ \textit{flatten} \ u. \end{aligned}$$

Unfortunately, *flatten* is slow. Let us try to improve it.

- (a) Introduce *flatcat* $t \ xs = \dots$ where *xs* is an accumulating parameter.
- (b) Express *flatten* in terms of *flatcat*.
- (c) Construct an efficient implementation of *flatten*. You will need some properties of $(++)$ proved in one of the exercises.