# Functional Programming Practicals

#### Shin-Cheng Mu

#### 1 Functions

1. Define a function *even* ::  $Int \rightarrow Bool$  that determines whether the input is an even number. You may use the following functions:

 $mod :: Int \to Int \to Int$ ,  $(==) :: Int \to Int \to Bool$ .

(Types of the functions written above are not in their most general form.)

- 2. Define a function that computes the area of a circle with given radius r (using 22/7 as an approximation to  $\pi$ ). The return type of the function might be *Double*.
- 3. Type in the definition of *smaller* into your working file. Then try the following:
  - (a) In GHCi, type :t smaller to see the type of smaller.
  - (b) Try applying it to some arguments, e.g. smaller 3 4, smaller 3 1.
  - (c) In your working file, define a new function st3 = smaller 3.
  - (d) Find out the type of st3 in GHCi. Try st3 4, st3 1. Explain the results you see.
- 4. Type in the definition of square in your working file.
  - (a) Define a function quad :: Int  $\rightarrow$  Int such that quad x computes  $x^4$ .
  - (b) Type in this definition into your working file. Describe, in words, what this function does.

twice :: 
$$(a \to a) \to (a \to a)$$
  
twice  $f \ x = f \ (f \ x)$ .

- (c) Define quad using twice.
- 5. Replace the previous *twice* with this definition:

$$\begin{array}{ll} twice & :: (a \rightarrow a) \rightarrow (a \rightarrow a) \\ twice \ f \ = f \cdot f \ . \end{array}$$

- (a) Does *quad* still behave the same?
- (b) Explain in words what this operator  $(\cdot)$  does.
- 6. Let the following identifiers have type:

 $\begin{array}{l} f :: Int \to Char \\ g :: Int \to Char \to Int \\ h :: (Char \to Int) \to Int \to Int \\ x :: Int \\ y :: Int \\ c :: Char \end{array}$ 

Which of the following expressions are type correct?

1.  $(g \cdot f) x c$ 2.  $(g x \cdot f) y$ 3.  $(h \cdot g) x y$ 4.  $(h \cdot g x) c$ 5.  $h \cdot g x c$ 

You may type the expressions into Haskell and see whether they type check. To define f, for example, include the following in your working file:

$$\begin{array}{l} f :: Int \to Char \\ f = undefined \end{array}$$

However, it is better if you can explain why the answers are as they are.

# 2 Products and Sums

1. In GHCi, issue the command

let x = ((1, 'a'), True)

This defines a new symbol x, with value  $((1, 'a'), \mathsf{True})$ .

- (a) Find out the type of x by a GHCi command.
- (b) How do you extract the 1 in x? Type an expression ... x into GHCi such that the result is 1.
- (c) Try to extract a' and True from x too.
- 2. Define a function swap ::  $(a, b) \rightarrow (b, a)$  that, as the name and type suggests, swaps the components
  - (a) Define swap using pattern matching: swap  $(x, y) = \dots$
  - (b) Define swap using fst and snd: swap  $x = \dots$
  - (c) Define *swap* using **case**.
- 3. Define a function  $half :: Int \to Either Int Int$  such that
  - if n is even, half n returns Left k with  $2 \times k = n$ ;
  - if n is odd, half n returns Right k with  $2 \times k + 1 = n$ .

You may use the function div. Find out what it does by youself.

- 4. What are the types of the following expressions?
  - (a)  $\lambda x \to (snd \ x, fst \ x)$ .
  - (b)  $\lambda f x \to f x x$ .
  - (c) Define:

 $\begin{array}{ll} \textit{myEither } f \ g \ x = \ \mathbf{case} \ x \ \mathbf{of} \\ & \text{Left } y \rightarrow f \ y \\ & \text{Right } z \rightarrow g \ z \ . \end{array}$ 

What is the type of myEither?<sup>1</sup>

(d)  $\lambda f \ x \ y \to f \ (fst \ y) \ x.$ (e)  $\lambda f \ x \ y \to fst \ (f \ y \ x).$ (f)  $\lambda x \ y \to x.$ (g)  $\lambda f \ g \ x \to f \ x \ (g \ x).$ 

<sup>&</sup>lt;sup>1</sup>There is such a function called *either*, which is sometimes quite convenient.

#### **3** Inductively Defined Functions on Lists

- 1. Define a function  $fstEven :: [Int] \rightarrow Int$  that returns the first even number of the input list.
- 2. Define a function  $hasZero :: [Int] \rightarrow Bool$  that returns True if and only if there is a 0 in the input list.
- 3. Define a function *myLast* that takes a list and returns the last (rightmost) element.
  - (a) Let the type be  $myLast :: [a] \to a$ . Define myLast.
  - (b) What happens in the previous definition of the input list is empty?
  - (c) Define  $myLast :: [a] \rightarrow Maybe a$ , which returns Nothing if the list is empty.
- 4. Define a function *pos* such that *pos* x *xs* looks for x in *xs* and returns its position. For example, *find* '*a*' "abc" yields 0, and *find* '*a*' "bac" yields 1.
  - (a) Let the type be  $pos :: Eq \ a \Rightarrow a \rightarrow [a] \rightarrow Int$ . In your definition, what happens if x is not in the list?
  - (b) Let the type be  $pos :: Eq \ a \Rightarrow a \rightarrow [a] \rightarrow Maybe Int$ , such that  $pos \ x \ xs$  returns Nothing if x is not in the list.
- 5. Define  $myConcat :: [[a]] \rightarrow [a]$  such that, for example myConcat [[1, 2, 3], [], [4], [5, 6]] = [1, 2, 3, 4, 5, 6]. Hint: use (+).
- 6. Define double ::  $[a] \rightarrow [a]$  such that, for example, double [1, 2, 3] = [1, 1, 2, 2, 3, 3].
- 7. Define *interleave* ::  $[a] \rightarrow [a] \rightarrow [a]$  such that, for example, *interleave* [1, 2, 3, 4] [5, 6, 7] = [1, 5, 2, 6, 3, 7, 4].
- 8. Define  $splitLR :: [Either \ a \ b] \to ([a], [b])$  such that, for example:

splitLR [Left 1, Left 3, Right 'a', Left 2, Right 'b'] = ([1,3,2], "ab").

9. Define a function  $fan :: a \to [a] \to [[a]]$  such that  $fan \ x \ xs$  inserts x into the 0th, 1st...nth positions of xs, where n is the length of xs. For example:

 $fan \ 5 \ [1,2,3,4] = [[5,1,2,3,4], [1,5,2,3,4], [1,2,5,3,4], [1,2,3,5,4], [1,2,3,4,5]] \ .$ 

10. Define  $perms :: [a] \to [[a]]$  that returns all permutations of the input list. For example:

perms [1, 2, 3] = [[1, 2, 3], [2, 1, 3], [2, 3, 1], [1, 3, 2], [3, 1, 2], [3, 2, 1]].

11. Try to define functions *inits* and *tails* yourself, and make sure you understand them. Recall that *inits* [1, 2, 3] = [[], [1], [1, 2], [1, 2, 3]], and *tails* [1, 2, 3] = [[1, 2, 3], [2, 3], [3], [1]].

### 4 Inductively Defined Functions on Natural Numbers

- 1. Define  $mul :: \mathbb{N} \to \mathbb{N} \to \mathbb{N}$  such that  $mul \ m \ n = m \times n$ , by induction on natural number, using addition (+).
- 2. Define  $myMin :: \mathbb{N} \to \mathbb{N} \to \mathbb{N}$  that returns the smaller of its two arguments. There is a built-in operator (min) for this, but try defining it inductively on natural numbers.
- 3. Define a function  $elemAt :: \mathbb{N} \to [a] \to a$  such that  $elemAt \ n \ xs$  yields the nth element of  $xs^2$ .
- 4. Define a function *insertAt* ::  $\mathbb{N} \to a \to [a] \to [a]$  such that *insertAt*  $n \ x \ xs$  inserts x into xs such that the *n*th element of the new list is x.

<sup>&</sup>lt;sup>2</sup>This function is denoted (!!) in the standard library.

# 5 User-Defined Inductive Datatypes

1. Consider the type

data ETree  $a = \text{Tip } a \mid \text{Bin } (ETree a) (ETree a)$ .

- (a) How is it different from the type *Tree* in the lecture note?
- (b) Define Define  $minT :: ETree Int \to Int$ , which computes the minimal element in a tree. The operator for binary minimum in Haskell is  $min :: Ord \ a \Rightarrow a \to a \to a$ .
- 2. Define  $minT :: Tree Int \to Int$ , which computes the minimal element in a tree. The operator for binary minimum in Haskell is  $min :: Ord \ a \to a \to a \to a$ . And the largest Int in Haskell is denoted by maxBound.
- 3. Define  $mapT :: (a \to b) \to Tree \ a \to Tree \ b$ , which applies the functional argument to each element in a tree.
- 4. Define *flatten* :: Tree  $a \to [a]$  that traverses a tree and collects all the labels, in-order, in a list. For example,

flatten (Node 4 (Node 2 (Node 1 Null Null) (Node 3 Null Null)) (Node 6 (Node 5 Null Null) (Node 7 Null Null)))

yields [1, 2, 3, 4, 5, 6, 7]. **Hint**: use (++).

- 5. A binary search tree is a tree of type Tree a, with Ord a, defined by:
  - 1. Null is a binary search tree, and
  - 2. Node x t u is a binary search tree if:
    - every label in t is less than x,
    - every label in u is greater than x, and
    - t and u are also binary search trees.

Define (assuming that t is a binary search tree):

- (a)  $memberT :: Ord \ a \Rightarrow a \rightarrow Tree \ a \rightarrow Bool$ , such that  $memberT \ x \ t$  determines whether x occurs in t, and
- (b) insert  $T :: Ord \ a \Rightarrow a \rightarrow Tree \ a \rightarrow Tree \ a$ , such that insert  $T \ x \ t$  inserts x into t and still returns a binary tree, if x does not appear in t, and returns t if x is in t.