Temporal Logics & Model Checking

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Specifications, descriptions, & verification

- specification:
  - The user’s requirement
- description (implementation):
  - The user’s description of the systems
  - No strict line between description and specification.
- verification:
  - Does the description satisfy the specification ?

Formal specification & automated verification

- formal specificaton:
  - specification with rigorous mathematical notations
- automated verification:
  - verification with support from computer tools.

Why formal specifications ?

- to make the engineers/users understand the system to design through rigorous mathematical notations.
- to avoid ambiguity/confusion/misunderstanding in communication/discussion/reading.
- to specify the system precisely.
- to generate mathematical models for automated analysis.

- *But according to Goedel's incompleteness theorem, it is impossible to come up with a complete specification.*

Why automated verification ?

- to somehow be able to verify complexer & larger systems
- to liberate human from the labor-intensive verification tasks
  - to set free the creativity of human
- to avoid the huge cost of fixing early bugs in late cycles.
- to compete with the core verification technology of the future.

Specification & Verification ?

- Specification → Complete & sound.
- Verfication

  → Reducing bugs in a system.

  → Making sure there are very few bugs.

Very difficult!

Competitiveness of high-tech industry!

A way to survive for the students!

A way to survive for Taiwan!

$4 billion development effort

> 50% system integration & validation cost

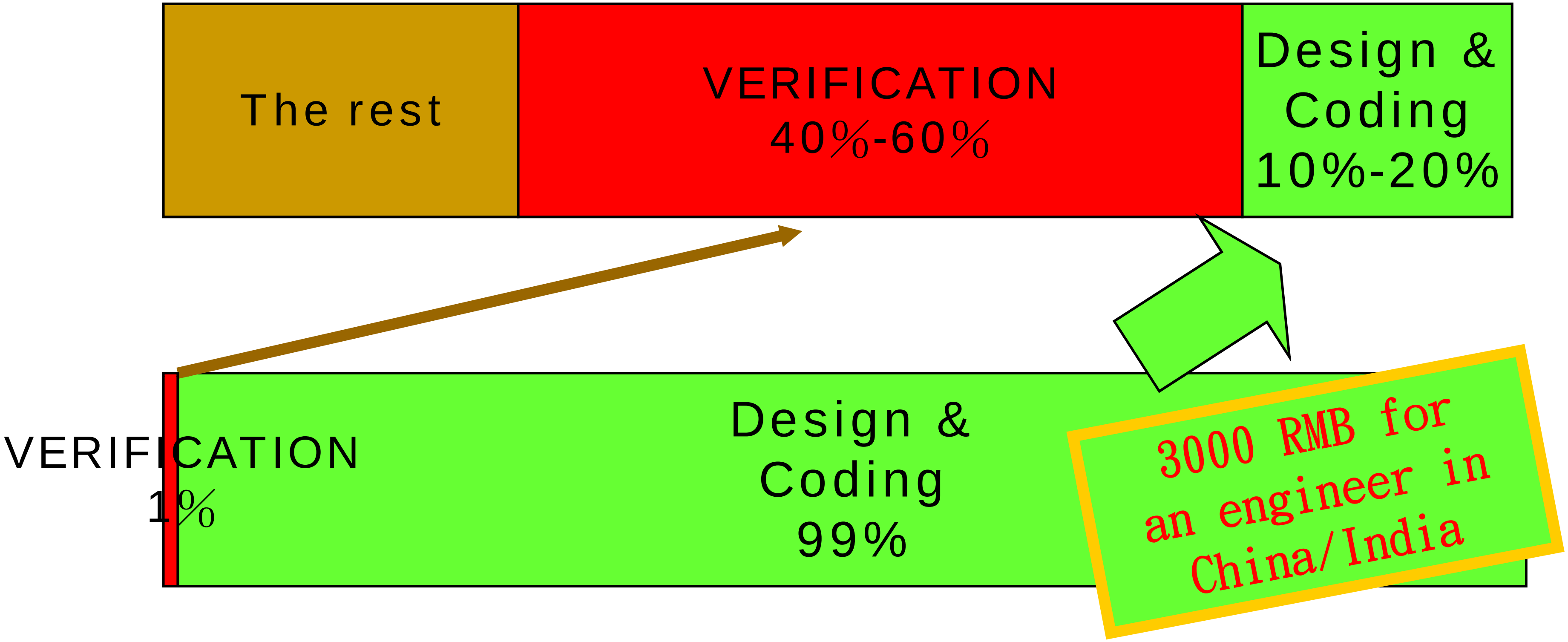
2,500,000+1,500,000 lines of codes (most in Ada)

400 horses
100 microprocessors

Bugs in complex software

- They take effects only with special event sequences.
  - the number of event sequences is factorial and super astronomical!
- It is impossible to check all traces with test/simulation.

Budget appropriation



The rest

VERIFICATION 40%-60%

Design & Coding 10%-20%

VERIFICATION 1%

Design & Coding 99%

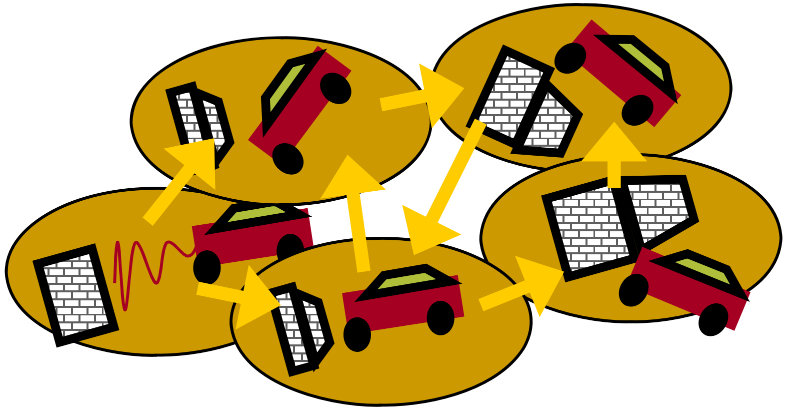
3000 RMB for an engineer in China/India

Training in Taiwan College

Three technologies in verification

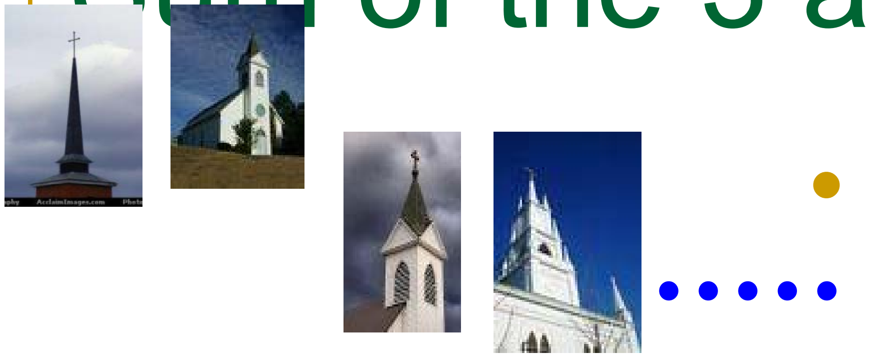
- Testing (real wall for real cars)
  - Expensive
  - Low coverage
  - Late in development cycles

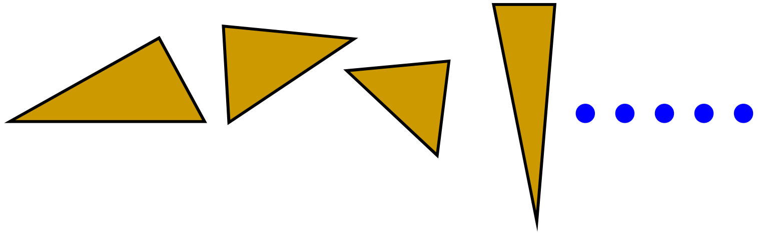
- Simulation(virtual wall for virtual ca
  - Economic
  - Low coverage
  - Don't know what you haven't seen.

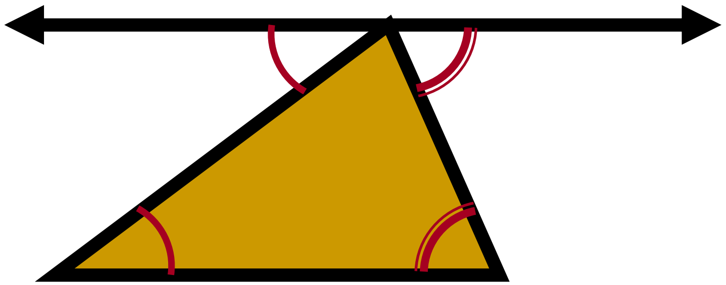


- Formal Verification (virtual car checked)
  - Expensive
  - Functional completeness
    - 100% coverage
  - Automated!
    - With algorithms and proofs.

Sum of the 3 angles = 180 ?

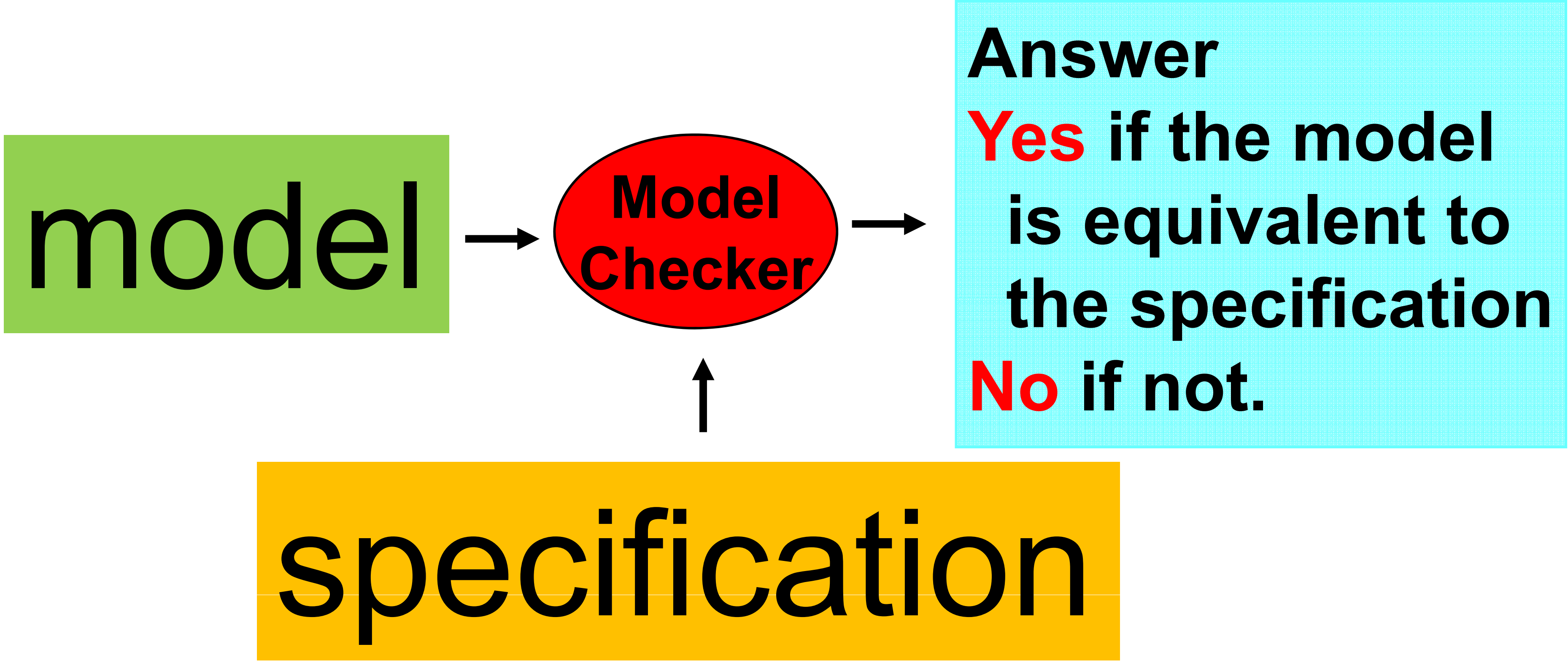






- Testing (check all ∆s you see)
  - Expensive
  - Low coverage
  - Late in development cycles
- Simulation (check all ∆s you draw)
  - Economic
  - Low coverage
  - Don't know what you haven't seen.
- Formal Verification(we prove it.)
  - Expensive
  - Functional completeness
    - 100% coverage
  - Automated!
    - With algorithms and proofs.

Model-checking
- a general framework for verification of sequential systems

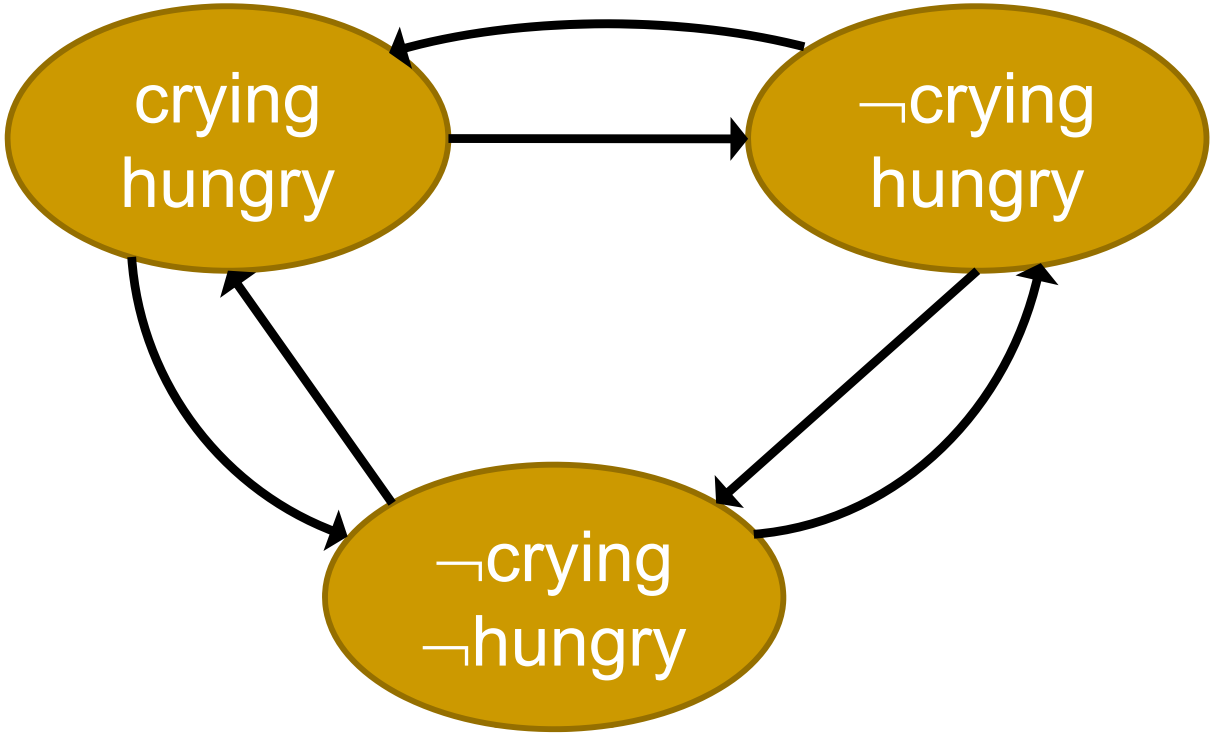


Answer
Yes if the model is equivalent to the specification
No if not.

Models & Specifications - formalism

Whenever a baby cries, it is hungry.

- Logics: □(crying → hungry)
- Graphs:



Models & Specifications
- fairness assumptions

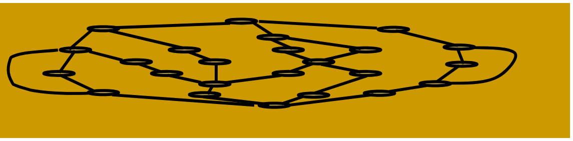
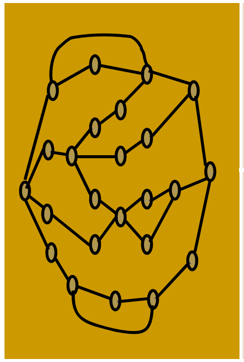
Some properties are almost impossible to verify without assumptions.

Example: $\Box(\text{start} \rightarrow \Diamond \text{ finish})$

To verify that a program halts, we assume

- CPU does not burn out.
- OS gives the program a *fair* share of CPU time.
- All the drivers do not stuck.
- ………

Model-checking - frameworks in our lecture

Spec Model							Logics			
			traces		Trees		Linear		Branching	
			F=∅	F≠∅	F=∅	F≠∅	F=∅	F≠∅	F=∅	F≠∅
	traces	F=∅	✓	✓			✓	✓		
		F≠∅	✓	✓			✓	✓		
	Trees	F=∅			☑	✓			☑	✓
		F≠∅			✓	✓			✓	✓
Logics	Linear	F=∅					☑	☑		
		F≠∅					☑	☑		
	Branching	F=∅							✓	✓
		F≠∅							✓	✓

✓: known; ☑: discussed in the lecture

History of Temporal Logic

- Designed by philosophers to study the way that time is used in natural language arguments
- Reviewed by Prior [PR57, PR67]

- Brought to Computer Science by Pnueli [PN77]
- Has proved to be useful for specification of concurrent systems

Framework

- Temporal Logic is a class of Modal Logic
- Allows qualitatively describing and reasoning about changes of the truth values over time
- Usually implicit time representation
- Provides variety of temporal operators (*sometimes, always*)
- Different views of time (branching vs. linear, discrete vs. continuous, past vs. future, etc.)

Outline

- Linear
  - LPTL (Linear time Propositional Temporal Logics)
- Branching
  - CTL (Computation Tree Logics)
  - CTL* (the full branching temporal logics)

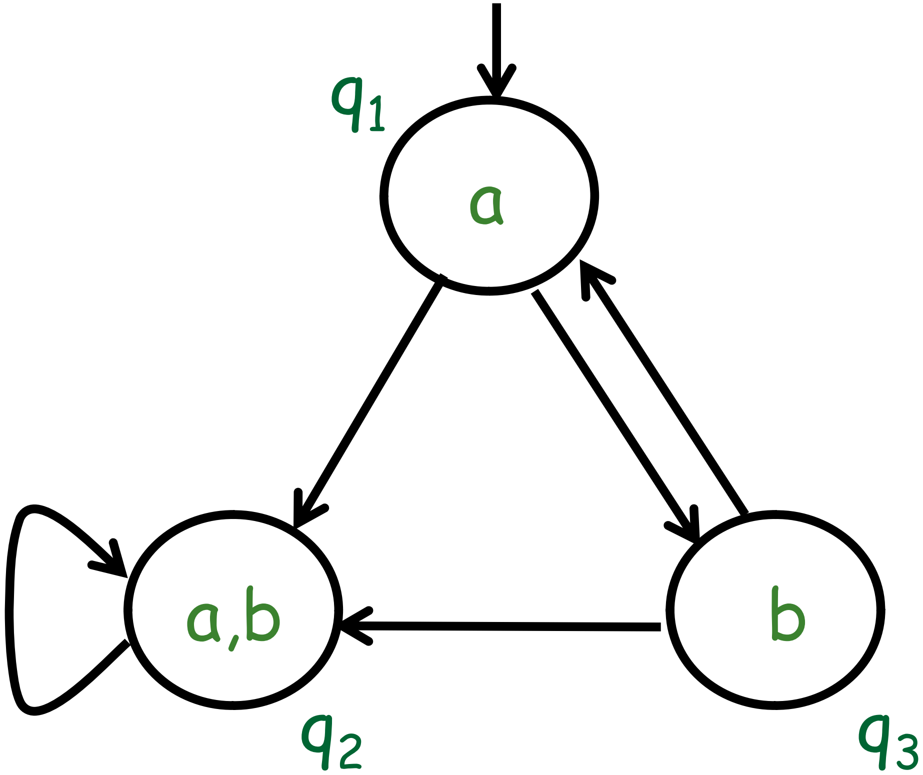
Kripke structure

$A = (S, S_0, R, L)$

- S
  - a set of all states of the system
- $S_0 \subseteq S$
  - a set of initial states
- $R \subseteq S \times S$
  - a transition relation between states
- $L : S \mapsto 2^P$
  - a function that associates each state with set of propositions true in that state

Kripke Model

- Set of states S
  - $\{q_1,q_2,q_3\}$
- Set of initial states S_0
  - $\{q_1\}$
- Set of atomic propositions AP
  - {a,b}



Example of Kripke Structure

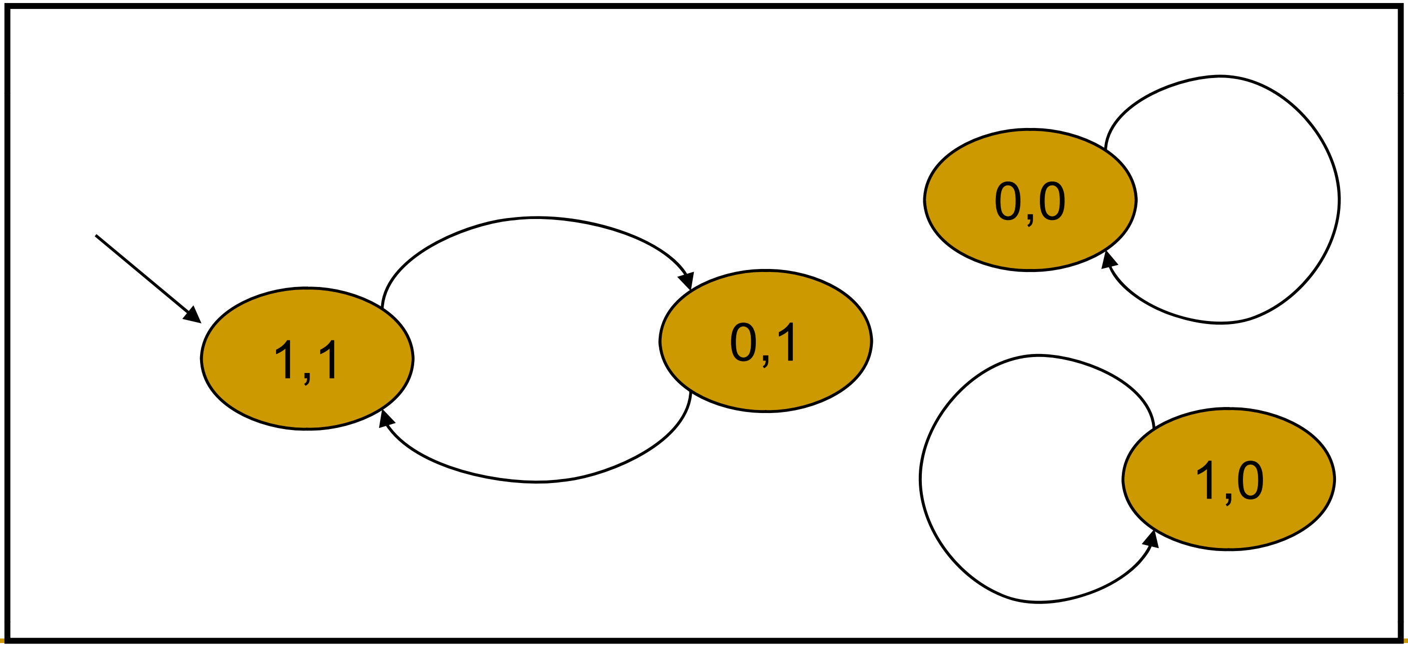
Suppose there is a program

```
initially x=1 and y=1;
while true do
  x:=(x+y) mod 2;
endwhile
```

where x and y range over *D*={0,1}

Example of Kripke Structure

- $S=D\text{x}D$
- $S_0=\{(1,1)\}$
- $R=\{((1,1),(0,1)),((0,1),(1,1)),((1,0),(1,0)),((0,0),(0,0))\}$
- L((1,1))={x=1,y=1},L((0,1))={x=0,y=1}, L((1,0))={x=1,y=0},L((0,0))={x=0,y=0}



BNF, syntax definitions Note!

Be sure how to read BNF !

- used for define syntax of context-free language
- important for the definition of
  - automata predicates and
  - temporal logics
- Used throughout the lectures!
- In exam: violate the syntax rules ➔ no credit.

$$A ::= c \mid x \mid (M) \mid A_1 + A_2 \mid A_1 – A_2$$

$$M ::= c \mid x \mid (A) \mid M_1 * M_2 \mid M_1 / M_2$$

c is an integer

x is a variable name.

BNF, syntax definitions

$A ::= c \mid x \mid (M) \mid A_1{+}A_2 \mid A_1{-}A_2$

$M ::= c \mid x \mid (A) \mid M_1{*}M_2 \mid M_1/M_2$

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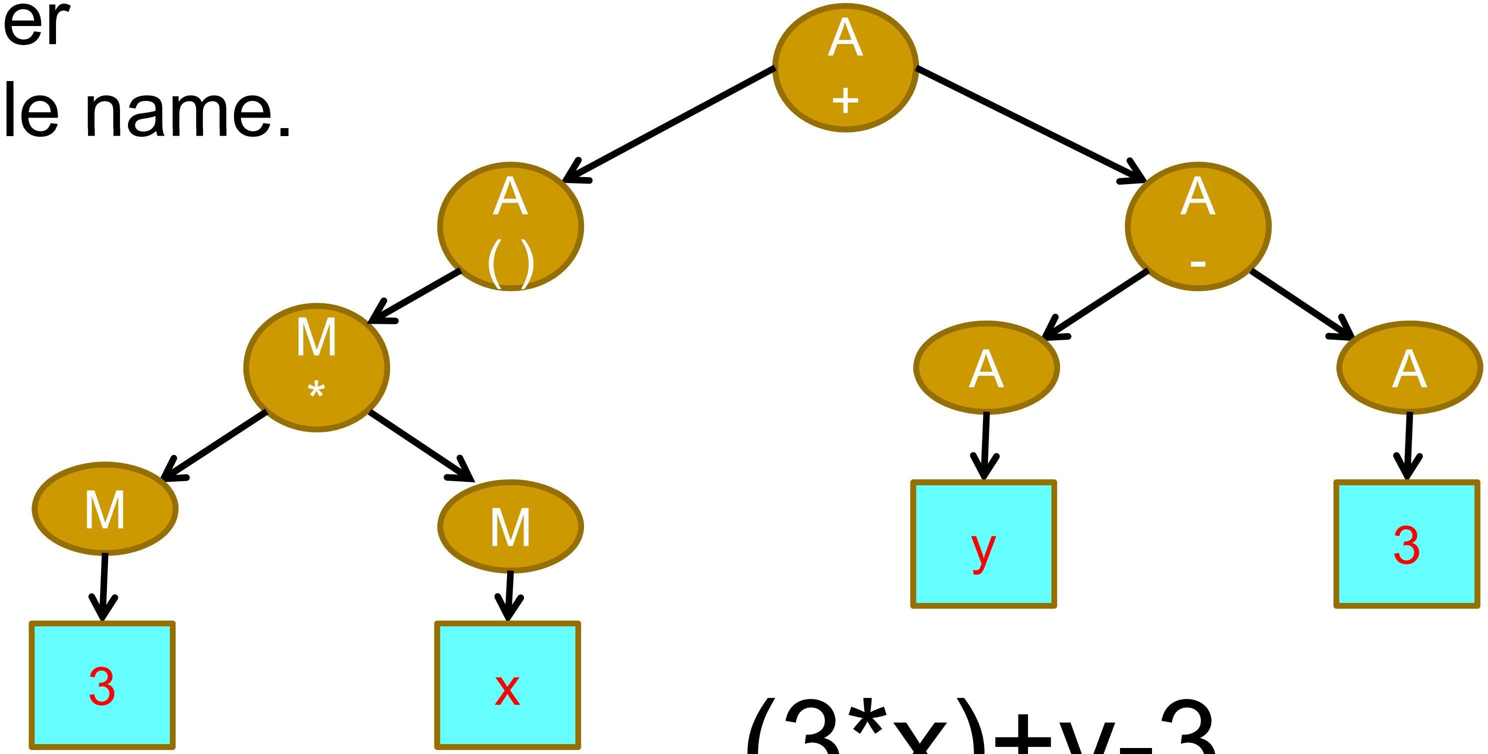


(3*x)+y-3

BNF, syntax definitions - derivation trees (from top down)

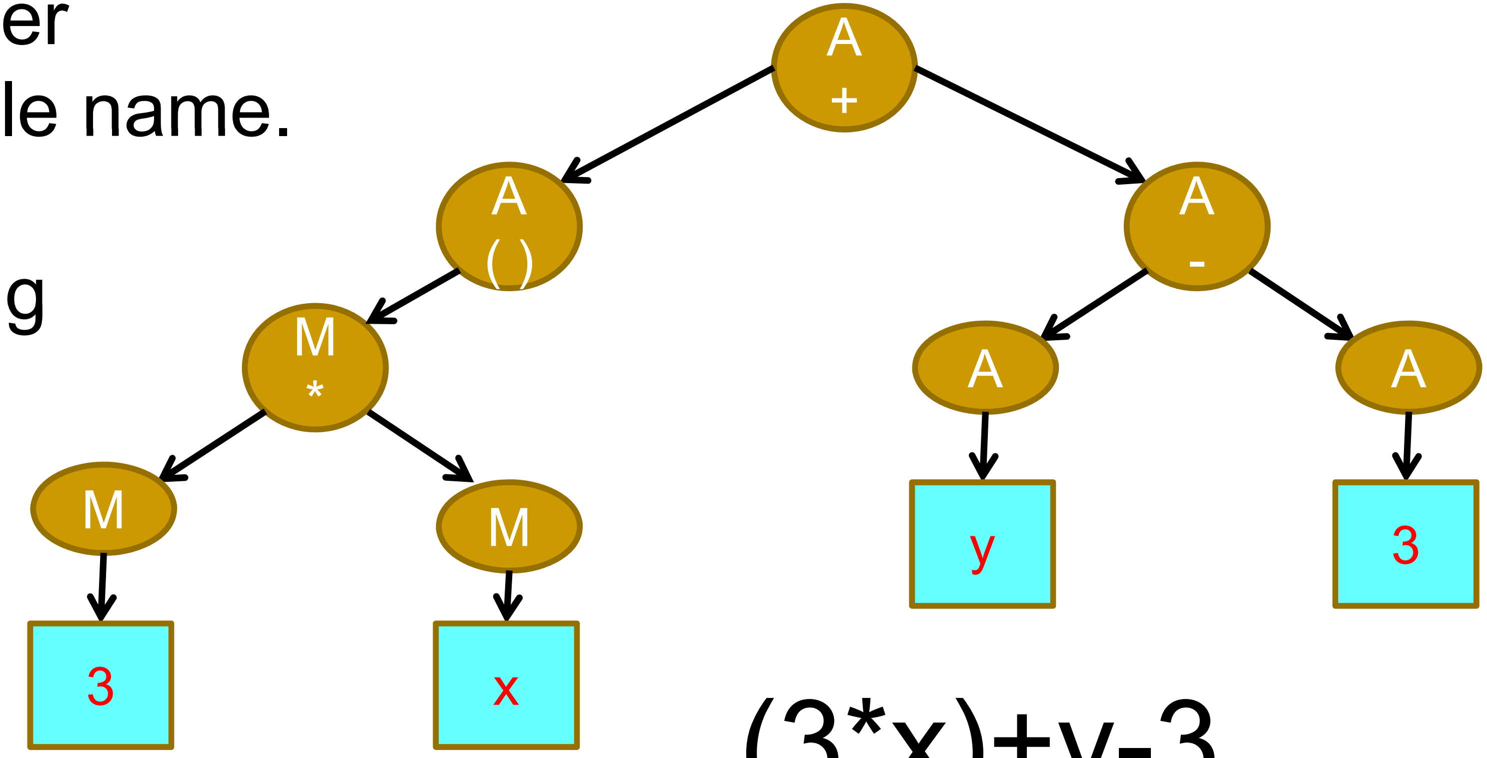
$A ::= c \mid x \mid (M) \mid A_1{+}A_2 \mid A_1{-}A_2$

$M ::= c \mid x \mid (A) \mid M_1{*}M_2 \mid M_1/M_2$

c is an integer

x is a variable name.

used in string generation.



(3*x)+y-3

BNF, syntax definitions - parsing trees (from bottom up)

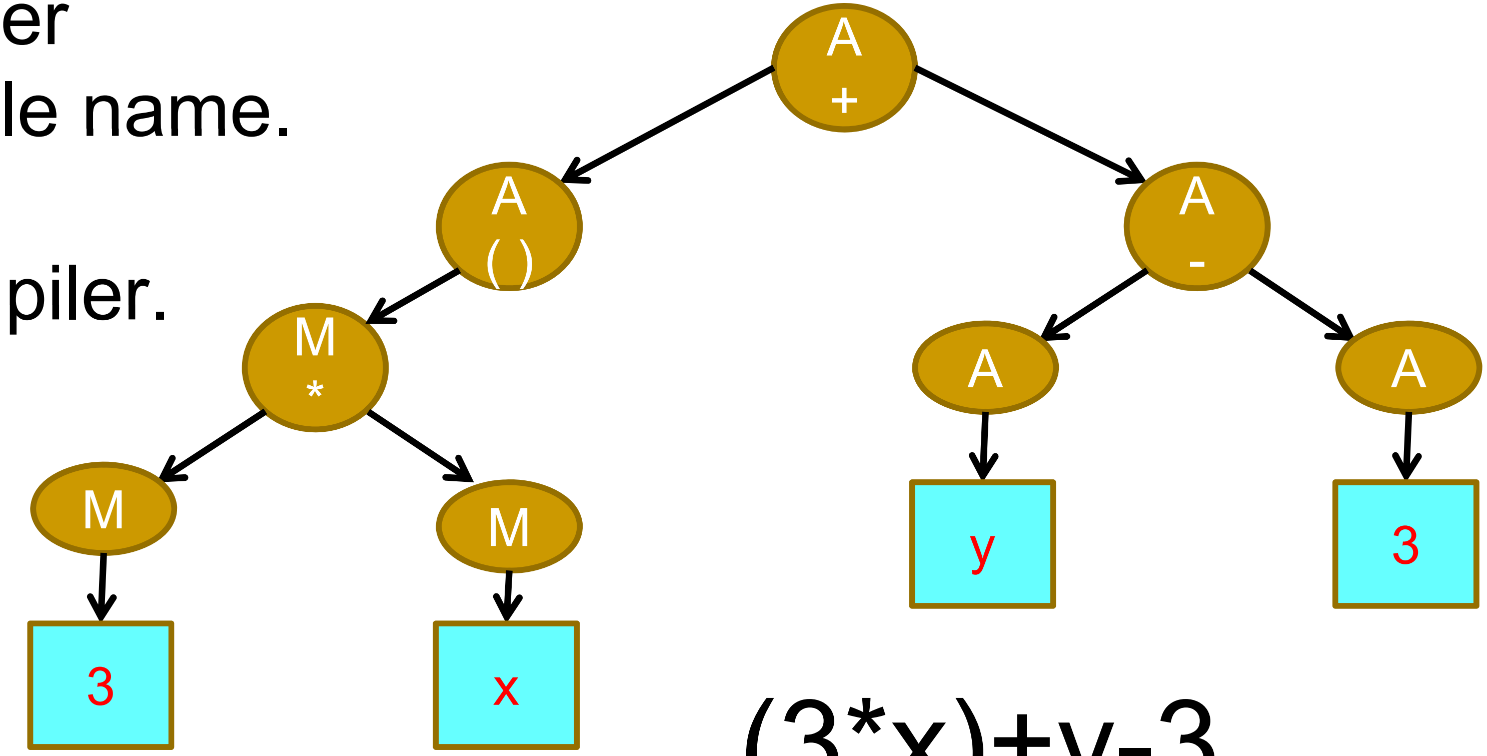
A ::= c | x | (M) | A_1+A_2 | A_1–A_2

M ::= c | x | (A) | M_1*M_2 | M_1/M_2

c is an integer

x is a variable name.

used in compiler.



(3*x)+y-3

Temporal Logics : Catalog

propositional ↔ first-order

global ↔ compositional

branching ↔ linear-time

points ↔ intervals

discrete ↔ continuous

past ↔ future

Temporal Logics

- Linear
  - LPTL (Linear time Propositional Temporal Logics)
    - LTL, PTL, PLTL
- Branching
  - CTL (Computation Tree Logics)
  - CTL* (the full branching temporal logics)

Amir Pnueli
1941

- Professor, Weizmann Institute
- Professor, NYU
- Turing Award, 1996

Presentation of a gift at ATVA /FORTE 2005, Taipei

LPTL (PTL, LTL)
Linear-Time Propositional Temporal Logic

Conventional notation：

- propositions : ***p*, *q*, *r*, …**
- sets : ***A*, *B*, *C*, *D*, …**
- states : ***s***
- state sequences : ***S***
- formulas : ***φ*,*ψ***
- Set of natural number：*N* **= {0, 1, 2, 3, …}**
- Set of real number：*R*

LPTL

Given ***P*** : a set of propositions,

a Linear-time structure : *state sequence*

$S = s_0\, s_1\, s_2\, s_3\, s_4 \ldots\ s_k \ldots\ldots$

s_k is a function of P where ***P*** ***{true,false}***

or $s_k \in 2^P$

example: P={a,b}

{a}{a,b}{a}{a}{b}…

Syntax definitions
Note!

Be sure how to read BNF !

- used for define syntax of context-free language
- important for the definition of
  - automata predicates and
  - temporal logics
- Used throughout the lectures!
- In exam: violate the syntax rules ➔ **no credit.**

A ::= (M) | A1 + A2 | A1 – A2
M ::= (A) | M1 * M2 | M1 / M2

LPTL
- syntax

syntax definition in BNF

$$\psi ::= \mathbf{true} \mid p \mid \neg\psi \mid \psi_1 \vee \psi_2 \mid \bigcirc\psi \mid \psi_1 \mathrm{U} \psi_2$$

abbreviation

$$\mathbf{false} \equiv \neg\, \mathrm{true}$$

$$\psi_1 \wedge \psi_2 \equiv \neg\, ((\neg\psi_1) \vee (\neg\psi_2))$$

$$\psi_1 \rightarrow \psi_2 \equiv (\neg\psi_1) \vee \psi_2$$

$$\Diamond\psi \equiv \mathbf{true}\ \mathrm{U}\psi$$

$$\Box\psi \equiv \neg\Diamond\neg\psi$$

LPTL

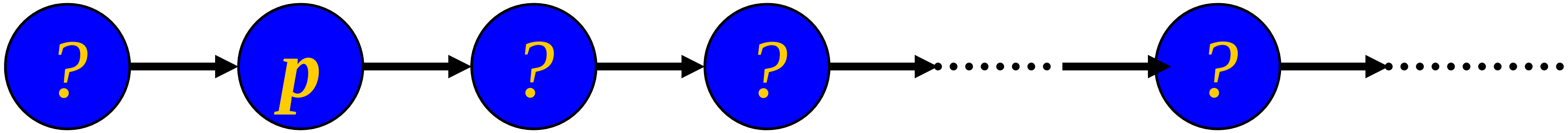
- syntax

Exam.	Symbol in CMU	
$\bigcirc p$	$\mathrm{X}p$	*p is* **true on next state**
$p \mathrm{U} q$	$p \mathrm{U} q$	**From now on, *p* is always true until *q* is true**
$\Diamond p$	$\mathrm{F}p$	**From now on, there will be a state where *p* is eventually (sometimes) true**
$\Box p$	$\mathrm{G}p$	**From now on, *p* is always true**

LPTL

- syntax

$\bigcirc p$ $\quad$ $\mathbf{X}p$ $\quad$ ***p is*** **true on next state**



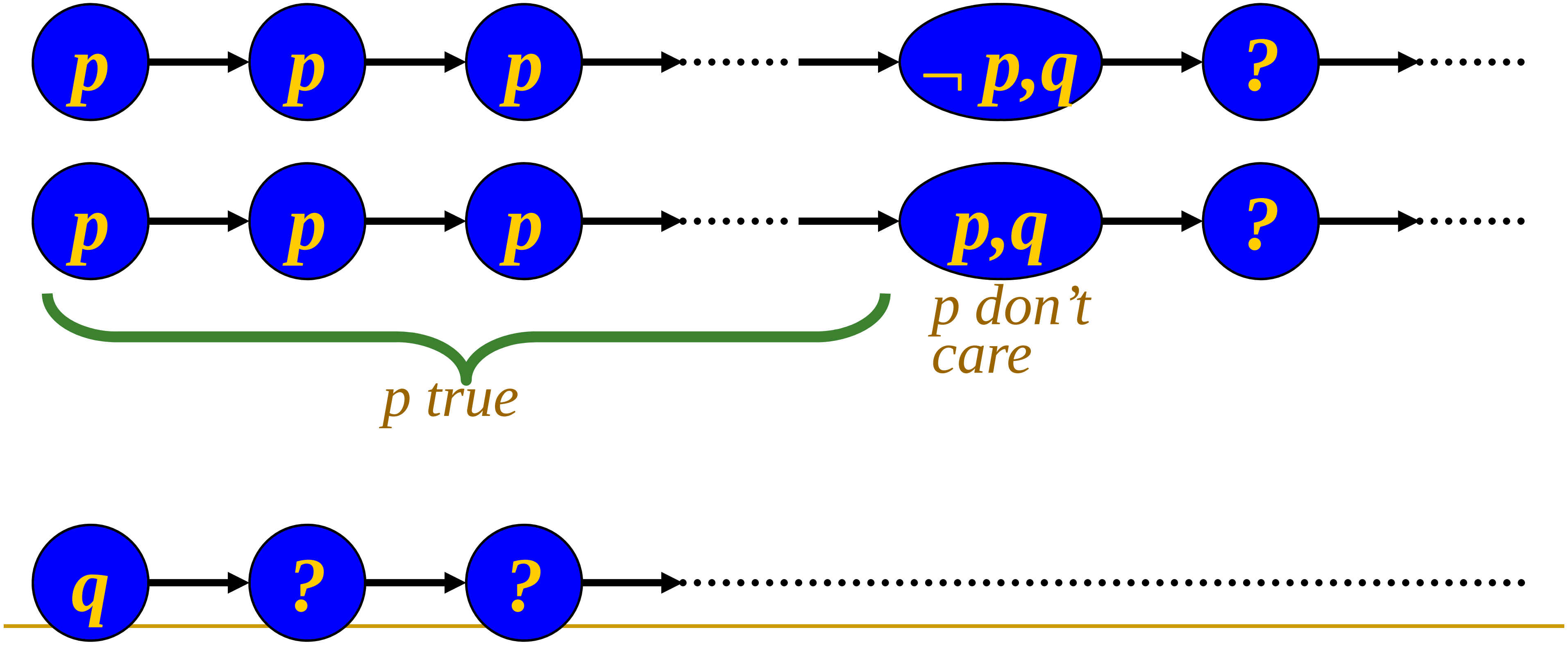
? : don’t care

LPTL
- syntax

p*U*q

p*U*q

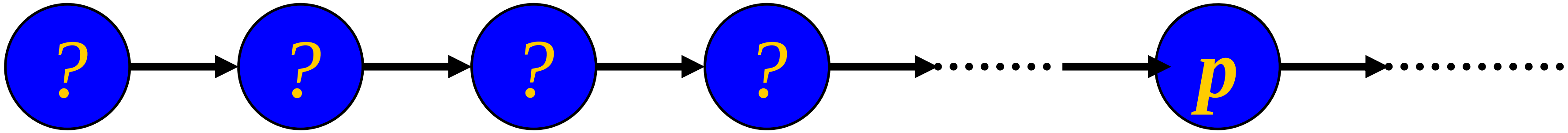
From now on, *p* is always true until *q* is true



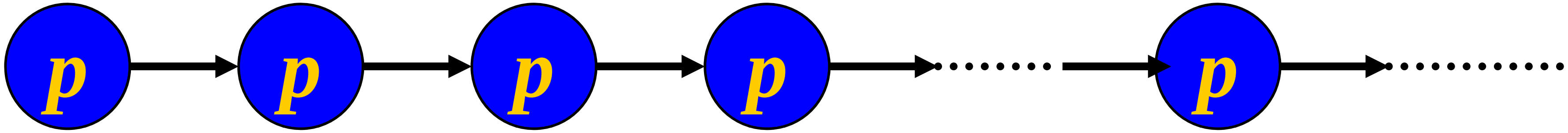
LPTL

- syntax

$\Diamond p$ Fp From now on, there will be a state where p is eventually (sometimes) true



$\Box p$ Gp From now on, p is always true



LPTL

- syntax

Two operator for Fairness

- $\Diamond^{\infty} p \equiv \Box\Diamond p$; *p* will happen infinitely many times
  infinitely often
- $\Box^{\infty} p \equiv \Diamond\Box p$; *p* will be always true after some time in the future
  almost everywhere

LPTL

- semantics

suffix path :

$$S = s_0\, s_1\, s_2\, s_3\, s_4\, s_5 \ldots\ldots\ldots$$

$$S^{(0)} = s_0\, s_1\, s_2\, s_3\, s_4\, s_5 \ldots\ldots\ldots$$

$$S^{(1)} = s_1\, s_2\, s_3\, s_4\, s_5\, s_6 \ldots\ldots\ldots$$

$$S^{(2)} = s_2\, s_3\, s_4\, s_5\, s_6 \ldots\ldots\ldots$$

$$S^{(3)} = s_3\, s_4\, s_5\, s_6 \ldots\ldots\ldots$$

$$S^{(k)} = s_k\, s_{k+1}\, s_{k+2}\, s_{k+3} \ldots\ldots\ldots$$

LPTL

- semantics

Given a state sequence

$$\boldsymbol{S= s_0\, s_1\, s_2\, s_3\, s_4 \ldots\, s_k \ldots\ldots}$$

We define $S \models \psi$ (S satisfies ψ） inductively as：

- $S \models \text{true}$
- $S \models p \;\Leftrightarrow\; s_0(p)=\text{true}$, or equivalently $p \in s_0$
- $S \models \neg\psi \Leftrightarrow S \models \psi$ is false
- $S \models \psi_1 \vee \psi_2 \;\Leftrightarrow S \models \psi_1$ or $S \models \psi_2$
- $S \models \bigcirc\psi \Leftrightarrow S^{(1)} \models \psi$
- $S \models \psi_1 U \psi_2 \;\Leftrightarrow \exists k \geq 0 (S^{(k)} \models \psi_2 \wedge \forall 0 \leq j < k (S^{(j)} \models \psi_1))$

LPTL
- semantics

- If a state sequence *S* satisfies φ (S⊨φ)

  then *S* is a model of φ*.*

- If there is a state sequence *S satφ,*

  then φ is *satisfiable*;

  else φ *is unsatisfiable*.

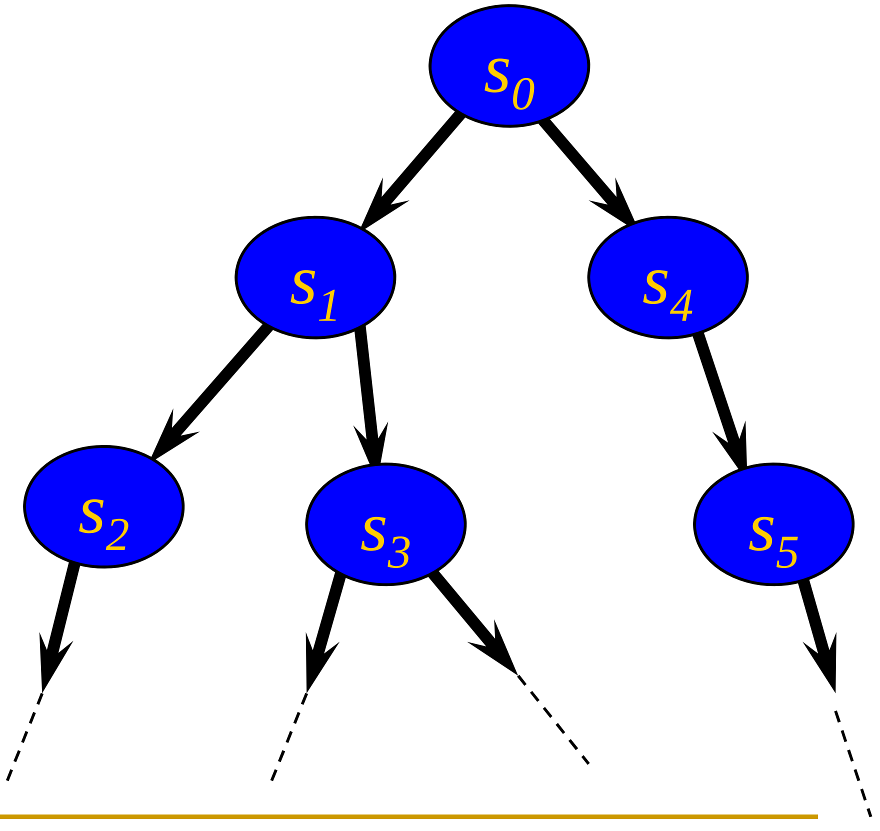
- If for all state sequence *S* ⊨ φ ，

  then φ is *valid*. （⊨φ）
- A formula φ characterizes its set of models.

Branching Temporal Logics

Basic assumption of tree-like structure

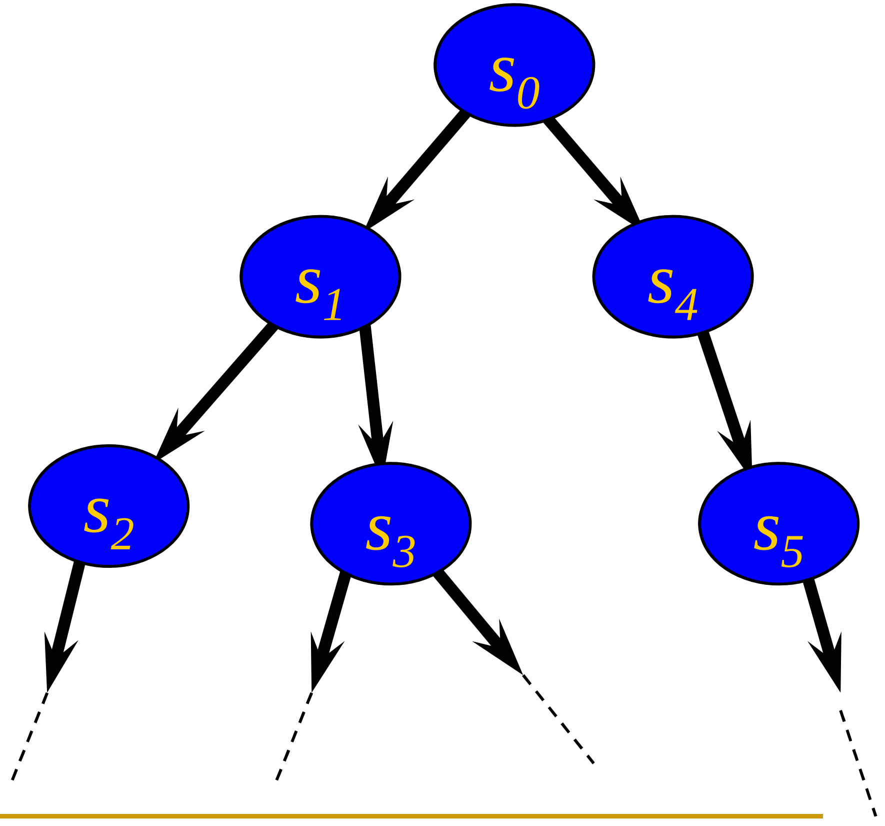
- Every node is a function of P→{true,false}
- Every state may have many successors



Branching Temporal Logics

Basic assumption of tree-like structure

- Every path is isomorphic as ***N***
  - Correspond to a state sequence



Path : $s_0\ s_1\ s_3 \ldots\ldots$

$s_0\ s_1\ s_2 \ldots\ldots$

$s_1\ s_3 \ldots\ldots$

$s_4\ s_5 \ldots\ldots$

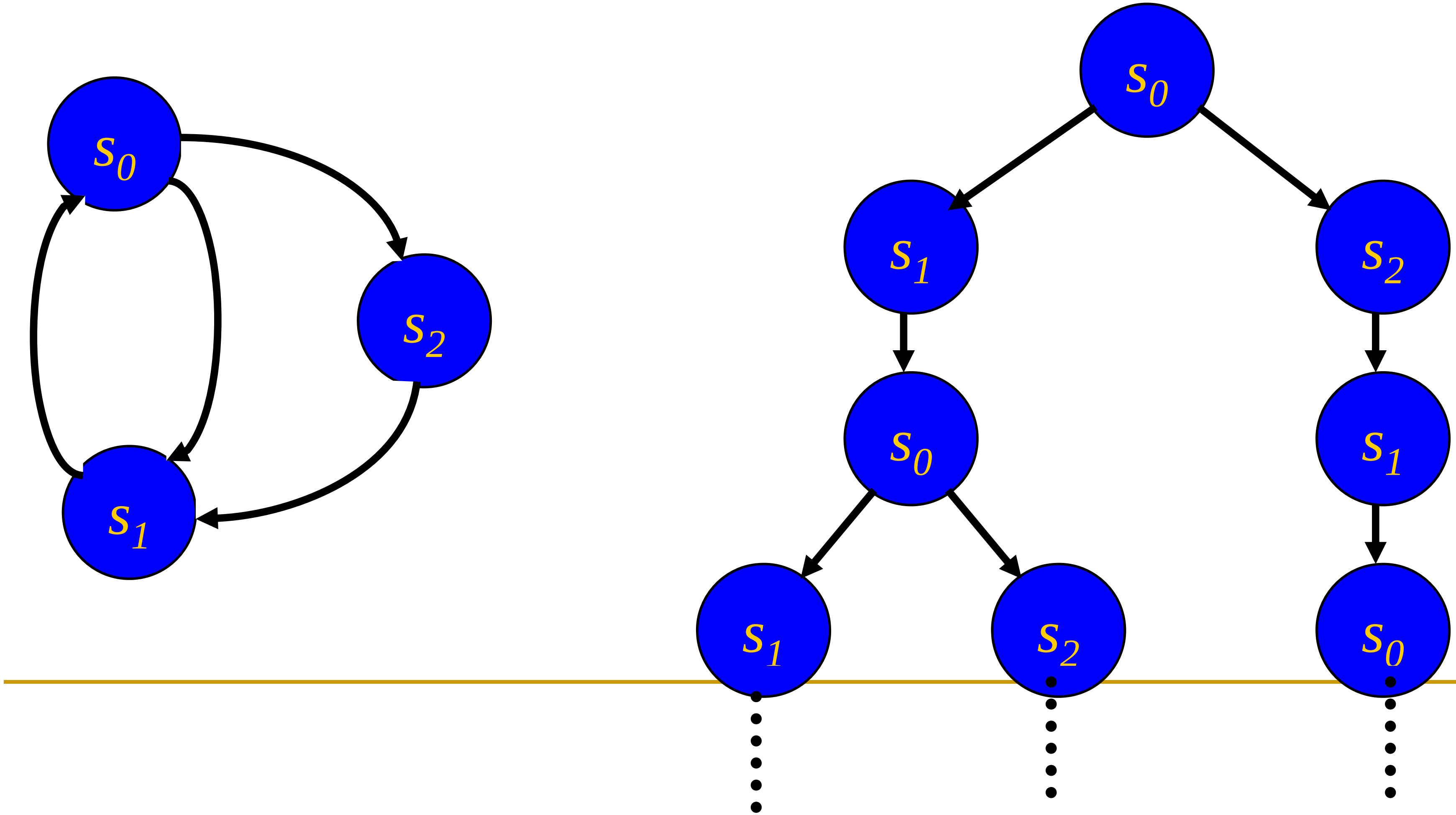
Branching Temporal Logic

It can accommodate infinite and dense state successors

- In CTL and CTL*, it can’t tell
  - Finite and infinite
    - Is there infinite transitions？
  - Dense and discrete
    - Is there countable（ω）transitions？

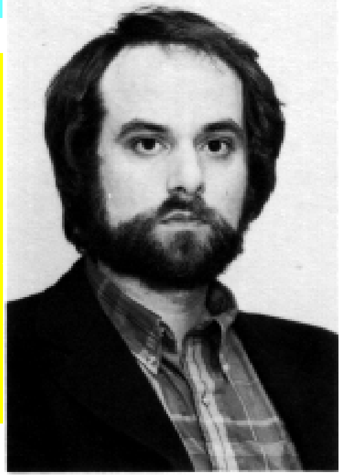
Branching Temporal Logic

Get by flattening a finite state machine



CTL(Computation Tree Logic)

Edmund M. Clarke
Professor, CS & ECE
Carnegie Mellon University



E. Allen Emerson
Professor, CS
The University of Texas at Austin

Chin-Laung Lei
Professor, EE
National Taiwan University

CTL(Computation Tree Logic) - syntax

$\varphi ::= \text{true} \mid p \mid \neg\varphi \mid \varphi_1 \vee \varphi_2 \mid \exists\bigcirc\varphi \mid \forall\bigcirc\varphi$

$\quad \mid \exists\varphi_1 \mathrm{U} \varphi_2 \mid \forall\varphi_1 \mathrm{U} \varphi_2$

abbreviation：

false	≡	¬ true
$\varphi_1 \wedge \varphi_2$	≡	$\neg((\neg\varphi_1) \vee (\neg\varphi_2))$
$\varphi_1 \rightarrow \varphi_2$	≡	$(\neg\varphi_1) \vee \varphi_2$
$\exists\Diamond\varphi$	≡	$\exists \text{true } \mathrm{U}\varphi$
$\forall\Box\varphi$	≡	$\neg\exists\Diamond\neg\varphi$
$\forall\Diamond\varphi$	≡	$\forall \text{true } \mathrm{U}\varphi$
$\exists\Box\varphi$	≡	$\neg\forall\Diamond\neg\varphi$

CTL
- semantics

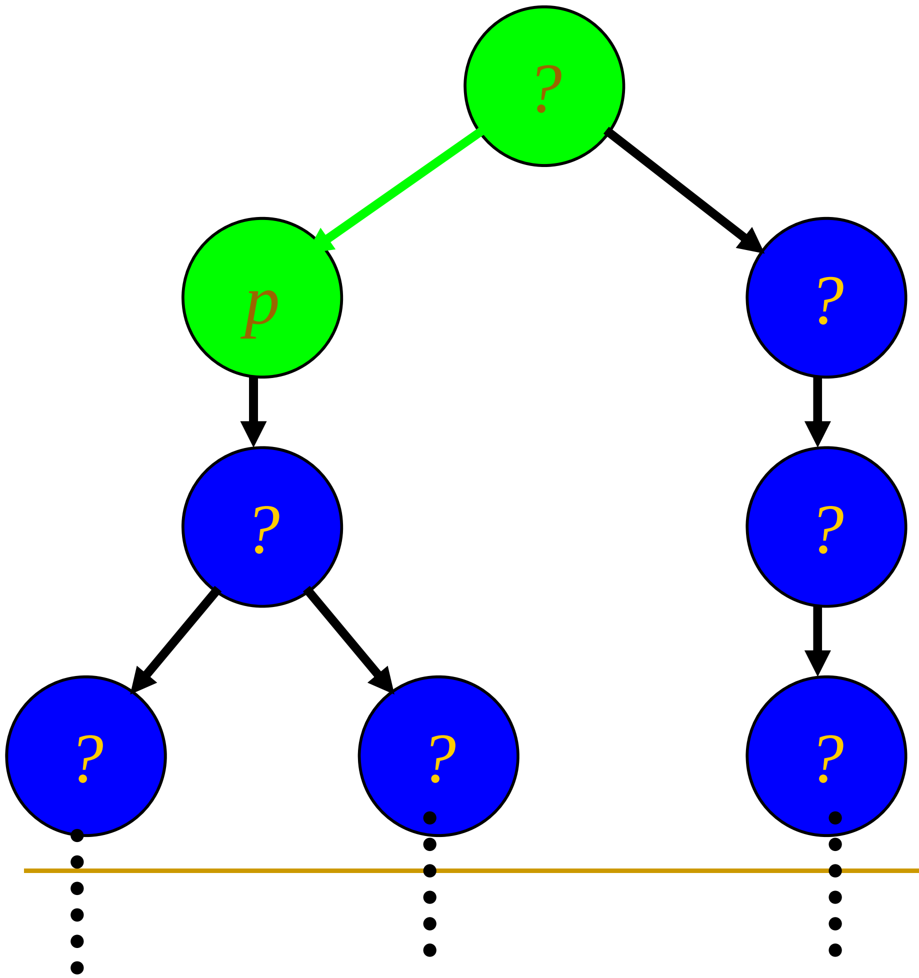
example	symbol in CMU	
$\exists \bigcirc p$	EXp	there exists a path where p is true on next state
$\exists p \mathcal{U}\ q$	$p$EU$q$	from now on, there is a path where p is always true until q is true
$\forall \bigcirc p$	AXp	for all path where p is true on next state
$\forall p \mathcal{U}\ q$	$p$AU$q$	from now on, for all path where p is always true until q is true

CTL

- semantics

∃○*p* **EX*p***

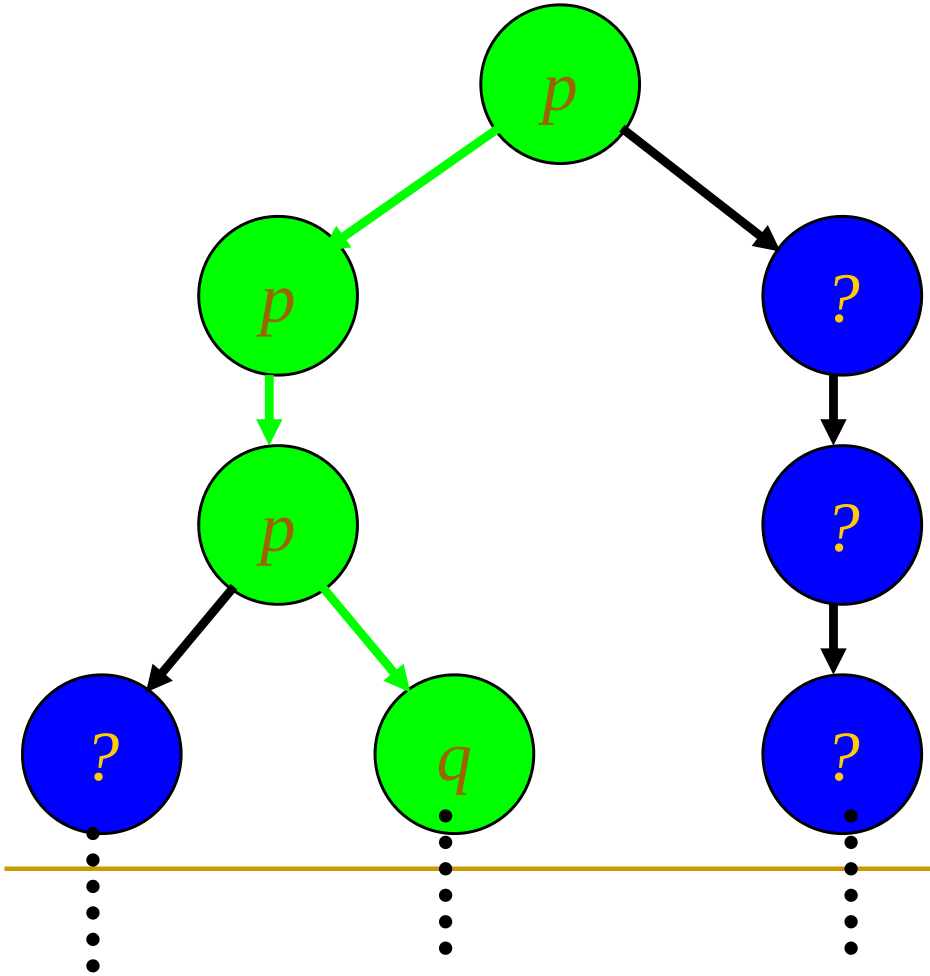
there exists a path where *p* is true on next state



CTL
- semantics

∃*p*U *q* ***p*EU*q***

from now on, there is a path where *p* is always true until *q* is true

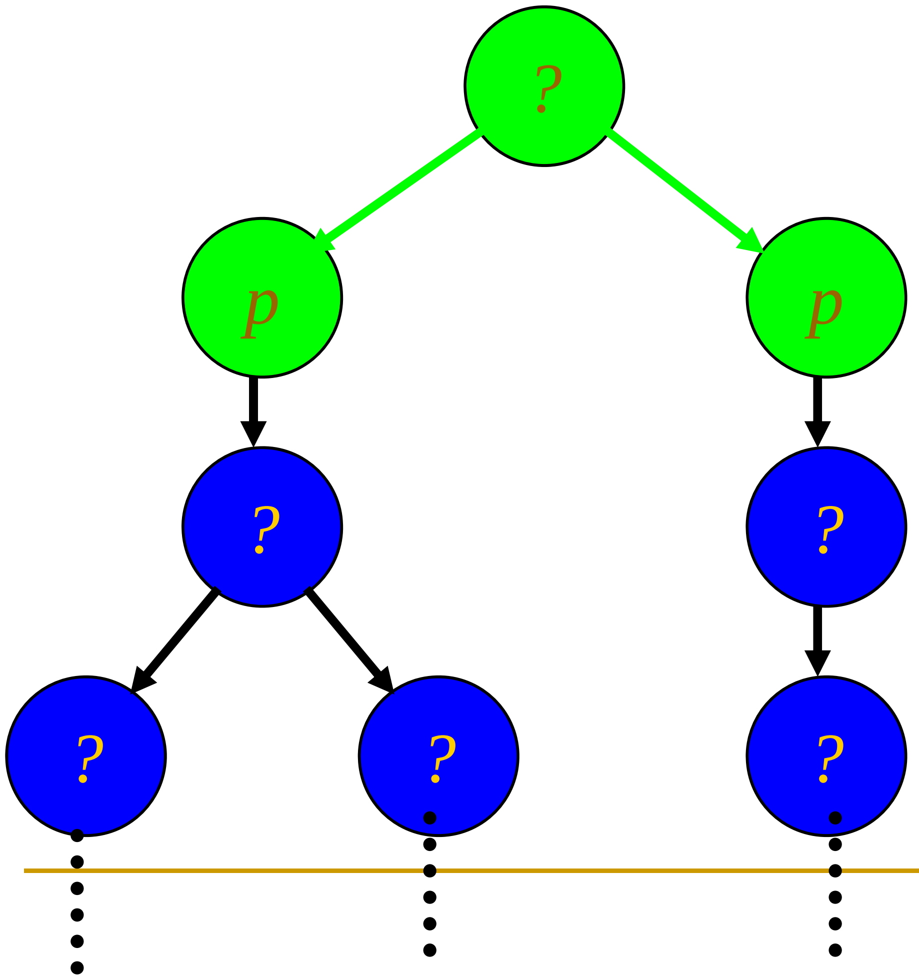


CTL

- semantics

$\forall \bigcirc p$ AX*p*

for all path where *p* is true on next state

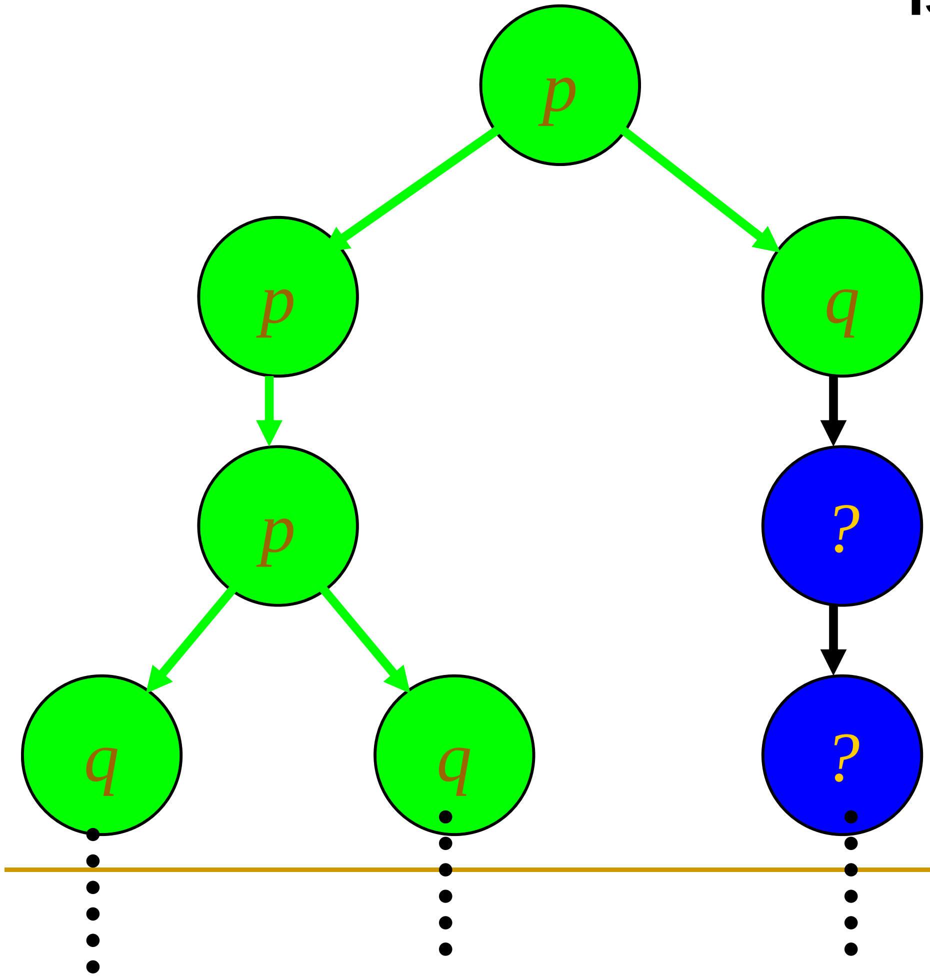


CTL

- semantics

$\forall p \cup\ q$ $\qquad$ ***p*AU*q***

from now on, for all path where *p* is always true until *q* is true



CTL

- semantic

Assume there are

- a tree stucture ***M***,
- one state ***s*** in ***M*,** and
- a CTL fomula **φ**

M*,*s⊨***φ*** means ***s*** in ***M*** satisfy **φ**

CTL

- semantics

s-path ：a path in *M* which stats from ***s***

s_0 -path:

$s_0\, s_1\, s_2\, s_3\, s_5$ ………

$s_0\, s_1\, s_6\, s_7\, s_8$ ………

s_1 -path:

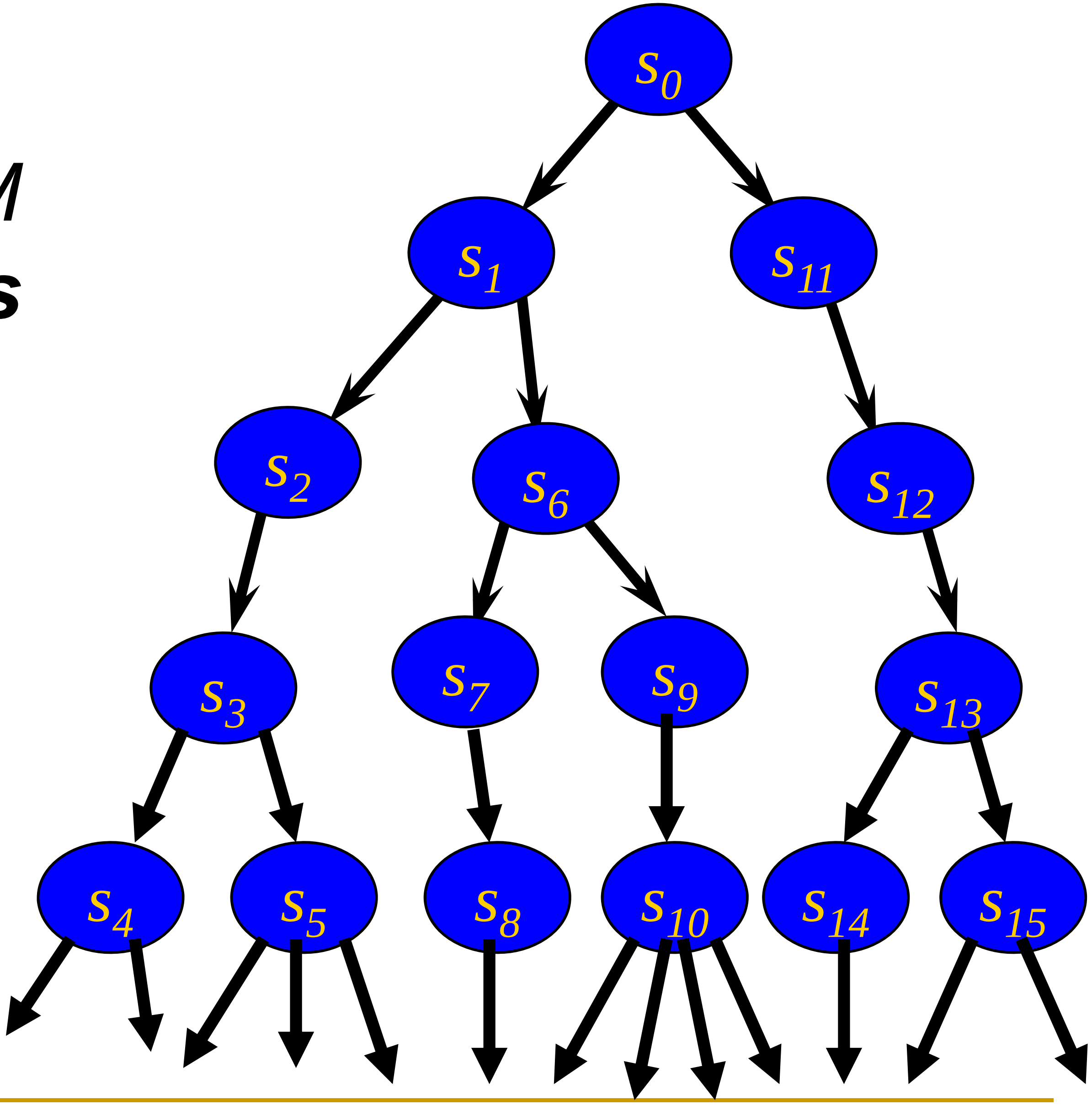
$s_1\, s_2\, s_3\, s_5$ …………

s_2 -path:

$s_2\, s_3\, s_5$ …………

s_{13} -path:

$s_{13}\, s_{15}$ …………

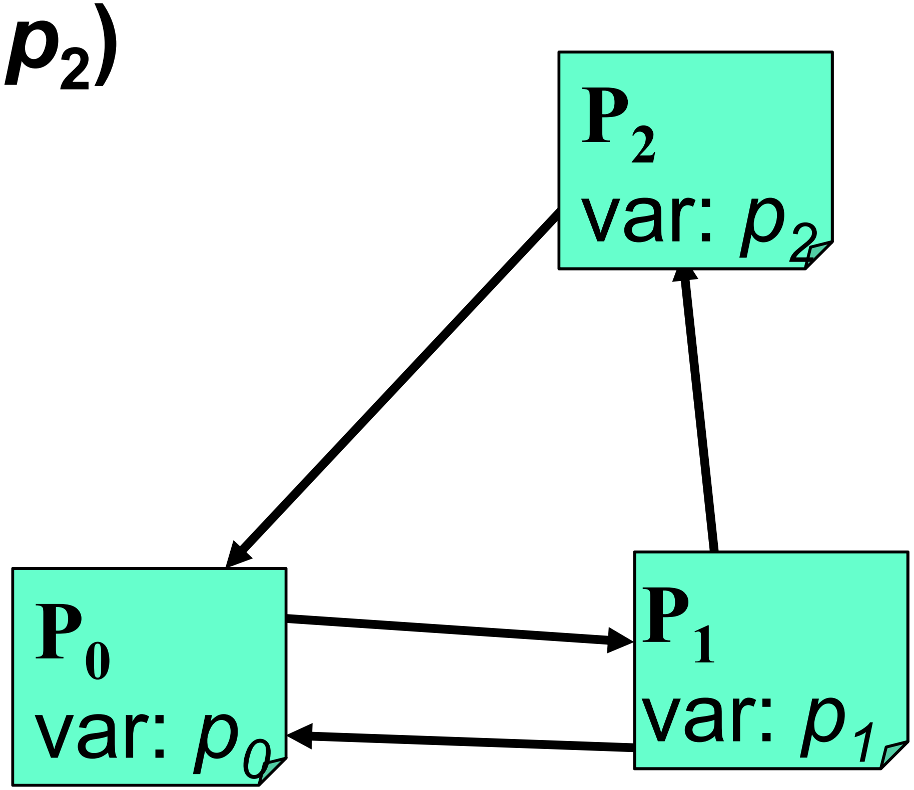


CTL
- semantics

- $M,s \models \text{true}$
- $M,s \models p \;\Leftrightarrow\; p \in s$
- $M,s \models \neg\varphi \Leftrightarrow$ it is false that $M,s \models \varphi$
- $M,s \models \varphi_1 \vee \varphi_2 \;\Leftrightarrow M,s \models \varphi_1$ or $M,s \models \varphi_2$
- $M,s \models \exists\bigcirc\varphi \Leftrightarrow \exists$ s-path $= s_0\ s_1\ \ldots\ldots(M,s_1 \models \varphi)$
- $M,s \models \forall\bigcirc\varphi \Leftrightarrow \forall$ s-path $= s_0\ s_1\ \ldots\ldots(M,s_1 \models \varphi)$
- $M,s \models \exists\varphi_1 U\varphi_2 \;\Leftrightarrow \exists$ s-path $= s_0\ s_1\ \ldots\ldots, \exists k\geq 0$
  $(M,s_k \models \varphi_2 \wedge \forall 0\leq j<k(M,s_j \models \varphi_1)$
- $M,s \models \forall\varphi_1 U\varphi_2 \;\Leftrightarrow \forall$s-path $= s_0\ s_1\ \ldots\ldots, \exists k\geq 0$
  $(M,s_k \models \varphi_2 \wedge \forall 0\leq j<k(M,s_j \models \varphi_1)$

CTL
- examples (I)

P_0:(p_0:=0 | p_0 := $p_0 \vee p_1 \vee p_2$)
P_1:(p_1:=0 | p_1 := $p_0 \vee p_1$)
P_2:(p_2:=0 | p_2 := $p_1 \vee p_2$)



If P_0 is true, it is possible that P_2 can be true after the next two cycles.

$\forall\Box(p_0 \rightarrow \exists\bigcirc\ \exists\bigcirc\ p_2)$

CTL
- examples (II)

1. If there are dark clouds, it will rain.

   $\forall\Box(\text{dark-clouds} \rightarrow \forall\Diamond \text{rain})$

2. if a buttefly flaps its wings, the New York stock could plunder.

   $\forall\Box(\text{buttefly-flap-wings} \rightarrow \exists\Diamond \text{NY-stock-plunder})$

3. if I win the lottery, I will be happy forever.

   $\forall\Box(\text{win-lottery} \rightarrow \forall\Box\ \text{happy})$

4. In an execution state, if an interrupt occurs in the next cycle, the interrupt handler will execute at the 2nd next cycle.

   $\forall\Box(\text{exec} \rightarrow \forall\bigcirc(\text{intrpt} \rightarrow \forall\bigcirc(\text{intrpt-handler})))$

CTL
- examples (III)

In an execution state, if an interrupt occurs in the next cycle, the interrupt handler will execute at the 2nd next cycle.

$$\forall\Box(\text{exec}\rightarrow\forall\bigcirc(\text{intrpt}\rightarrow\forall\bigcirc(\text{intrpt-handler})))$$

Some possible mistakes:

$\forall\Box(\text{exec}\rightarrow((\forall\bigcirc\ \text{intrpt})\rightarrow\forall\bigcirc\text{intrpt-handler}))$

$\forall\Box(\text{exec}\rightarrow((\forall\bigcirc\ \text{intrpt})\rightarrow\forall\bigcirc\forall\bigcirc\ \text{intrpt-handler}))$

CTL
- examples (IIIa)

Please draw a Kripke structure that tells

$$\forall\bigcirc(\text{intrpt}\rightarrow\forall\bigcirc(\text{intrpt-handler}))$$

from

$$(\forall\bigcirc\ \text{intrpt})\rightarrow\forall\bigcirc\text{intrpt-handler}$$

and

$$(\forall\bigcirc\ \text{intrpt})\ \rightarrow\forall\bigcirc\forall\bigcirc\ \text{intrpt-handler}$$

CTL
- important classes

- $\forall\Box\eta$: safety properties
  - η is always true in all computations from now.
- $\exists\Diamond\eta$: reachability properties
  - η is eventually true in some computation from now.
  - $\forall\Box\eta \equiv \neg\exists\Diamond\neg\eta$
- $\forall\Diamond\eta$: inevitabilities
  - η is eventually true in all computations from now.
- $\exists\Box\eta$
  - $\forall\Diamond\eta \equiv \neg\exists\Box\neg\eta$

CTL*

- syntax

- CTL* fomula (state-fomula)

  $\varphi ::= true \mid p \mid \neg\varphi_1 \mid \varphi_1 \vee \varphi_2 \mid \exists\psi \mid \forall\psi$

- path-fomula

  $\psi ::= \varphi \mid \neg\psi_1 \mid \psi_1 \vee \psi_2 \mid \bigcirc\psi_1 \mid \psi_1 U \psi_2$

CTL* is the set of all state-fomulas!

CTL*
- examples (1/4)

In a fair concurrent environment, jobs will eventually finish.

$\forall(((\Box\Diamond \textit{execute}_1) \wedge (\Box\Diamond \textit{execute}_2)) \rightarrow \Diamond \textit{finish})$

or

$\forall(((\Diamond^{\infty} \textit{execute}_1) \wedge (\Diamond^{\infty} \textit{execute}_2)) \rightarrow \Diamond \textit{finish})$

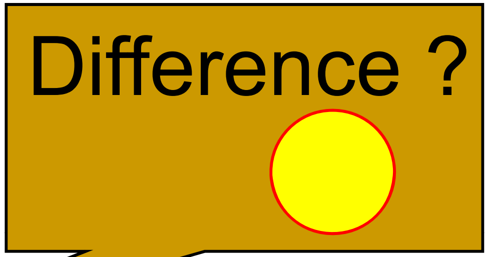
CTL* - examples (2/4)

No matter what, infinitely many comets will hit earth.

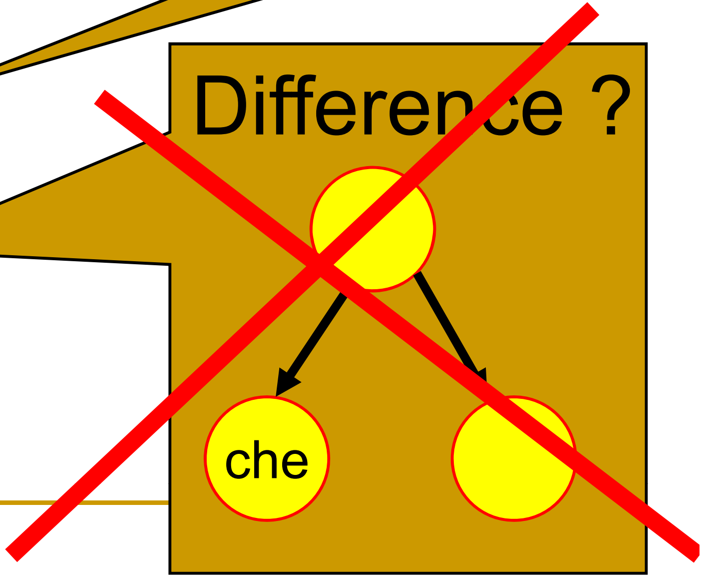
$\forall\Box\bigcirc\Diamond$comet-hit-earth

Why not CTL?

- $\forall\Box\ \forall\bigcirc\forall\Diamond$ comet-hit-earth
- $\forall\Box\forall\bigcirc\exists\Diamond$ comet-hit-earth



Difference ?



Difference ?

Exercise, please construct a model that tells the last from the first

CTL*
- examples (2/4)

No matter what, infinitely many comets will hit earth.

$\forall\Box\Diamond$comet-hit-earth

Or

$\forall\Diamond^{\infty}$ comet-hit-earth

Why not CTL?

- $\forall\Box\ \forall\ \Diamond$ comet-hit-earth
- $\forall\Box\ \exists\ \Diamond$ comet-hit-earth

What is the difference ?
weak
next !

true

CTL*
- Workout

The same according to lemma

- **(1)** ∀□◇comet-hit-earth
- **(2)** ∀□ ∀ ◇ comet-hit-earth
- **(3)** ∀□ ∃ ◇ comet-hit-earth

Please draw Kripke structures that tell

- (1) from (2) and (3)
- (2) from (1) and (3)
- (3) from (1) and (2)

CTL*
- examples (3/4)

If you never have a lover, I will marry you.

$\forall((\Box \text{you-have-no-lover}) \rightarrow \Diamond \text{marry-you})$

Why not CTL ?

- $(\forall\Box\ \text{you-have-no-lover}) \rightarrow \forall\ \Diamond\ \text{marry-you}$
- $(\forall\Box\ \text{you-have-no-lover}) \rightarrow \exists\ \Diamond\ \text{marry-you}$
- $(\exists\Box\ \text{you-have-no-lover}) \rightarrow \forall\ \Diamond\ \text{marry-you}$

CTL*
- Workout

- (1)$\forall((\Box \text{you-have-no-lover}) \rightarrow \Diamond \text{marry-you})$
- (2) $(\forall\Box\ \text{you-have-no-lover}) \rightarrow \forall\ \Diamond\ \text{marry-you}$
- (3) $(\forall\Box\ \text{you-have-no-lover}) \rightarrow \exists\ \Diamond\ \text{marry-you}$
- (4) $(\exists\Box\ \text{you-have-no-lover}) \rightarrow \forall\ \Diamond\ \text{marry-you}$

Please draw trees that tell

- ***(1) from (2)***
- ***(2) from (3)***
- ***(3) from (4)***
- ***(4) from (1)***

CTL*
- examples (4/4)

If I buy lottory tickets infinitely many times, eventually I will win the lottery.

$$\forall((\Box\Diamond \text{buy-lottery}) \rightarrow \Diamond \text{win-lottery})$$

or

$$\forall\ ((\Diamond^{\infty} \text{buy-lottery}) \rightarrow \Diamond\ \text{win-lottery})$$

CTL*

- semantics

suffix path :

$S = s_0\, s_1\, s_2\, s_3\, s_5 \ldots\ldots\ldots$

$S^{(0)} = s_0\, s_1\, s_2\, s_3\, s_5 \ldots\ldots\ldots$

$S^{(1)} = s_1\, s_2\, s_3\, s_5 \ldots\ldots\ldots$

$S^{(2)} = s_2\, s_3\, s_5 \ldots\ldots\ldots$

$S^{(3)} = s_3\, s_5 \ldots\ldots\ldots$

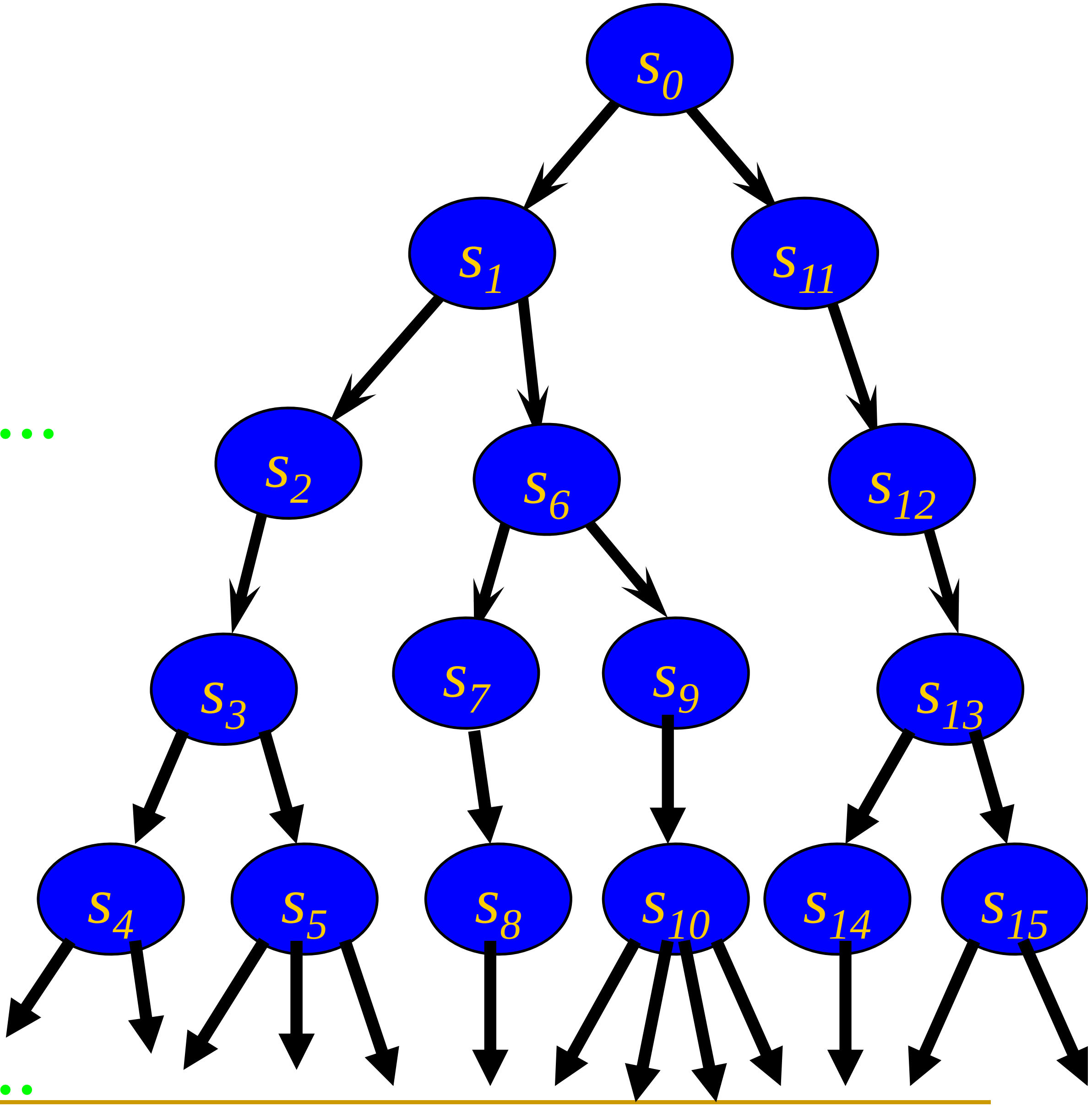
$S^{(4)} = s_5 \ldots\ldots\ldots$

$S = s_0\, s_1\, s_6\, s_7\, s_8 \ldots\ldots\ldots$

$S^{(2)} = s_6\, s_7 s_8 \ldots\ldots\ldots$

$S = s_0\, s_{11}\, s_{12}\, s_{13}\, s_{15} \ldots\ldots\ldots$

$S^{(3)} = s_{13}\, s_{15} \ldots\ldots\ldots\ldots$



CTL*

- semantics

state-fomula

φ::= true | p | $\neg\varphi_1$ | $\varphi_1 \vee \varphi_2$ | $\exists\psi$ | $\forall\psi$

- M,s $\models$ true
- M,s $\models$ p $\Leftrightarrow$ p $\in$s
- M,s $\models$ $\neg\varphi$ $\Leftrightarrow$ M,s $\models$ φ 是false
- M,s $\models$ $\varphi_1 \vee \varphi_2$ $\Leftrightarrow$ M,s $\models$ φ_1 or M,s $\models$ φ_2
- M,s $\models$ $\exists\psi$ $\Leftrightarrow$ $\exists$ s-path = S (S $\models$ ψ)
- M,s $\models$ $\forall\psi$ $\Leftrightarrow$ $\forall$ s-path = S (S $\models$ ψ)

CTL*
- semantics

path-fomula

$\psi ::= \varphi \mid \neg\psi_1 \mid \psi_1 \vee \psi_2 \mid \bigcirc\psi \mid \psi_1 \mathrm{U} \psi_2$

- *If* $S= s_0\, s_1\, s_2\, s_3\, s_4 \ldots\ldots\ldots, S \vDash \varphi \Leftrightarrow M,s_0 \vDash \varphi$
- $S \vDash \neg\psi_1 \Leftrightarrow S \vDash \psi_1$ 是false
- $S \vDash \psi_1 \vee \psi_2 \Leftrightarrow S \vDash \psi_1$ or $S \vDash \psi_1$
- $S \vDash \bigcirc\psi \Leftrightarrow S^{(1)} \vDash \psi$
- $S \vDash \psi_1 \mathrm{U} \psi_2 \Leftrightarrow \exists k \geq 0\ (S^{(k)} \vDash \psi_2 \wedge \forall 0 \leq j < k (S^{(j)} \vDash \psi_1))$

Expressiveness

Given a language $\mathcal{L}$,

- what model sets *L* can express ?
- what model sets L cannot ?

model set: a set of behaviors

A formula = a set of models (behaviors)

- for any $\varphi \in \mathcal{L}$, $[\varphi] \overset{\text{def}}{=} \{M \mid M \vDash \varphi\}$

A language = a set of formulas.

Expressiveness: Given a model set *F*,

F is expressible in $\mathcal{L}$ iff $\exists \varphi \in \mathcal{L}([\varphi]=F)$

Expressiveness

Comparison in expressiveness:

Given two languages L_1 and L_2

Definition: L_1 is ***more expressive than*** L_2 (L_2<L_1)

iff $\forall \varphi \in \mathcal{L}_2$ **([φ] is expressible in $\mathcal{L}_1$)**

Definition: L_1 and L_2 ***are expressively equivalent***

($L_1 \equiv L_2$) iff (L_2<L_1)$\wedge$(L_1<L_2)

Definition: **L_1 、 L_2 are** ***expressively incomparable*** iff

$\neg$((L_2<L_1)$\vee$(L_1<L_2))

Expressiveness
- branching-time logics

What to compare with ?

- finite-state automata on infinite trees.
- 2nd-order logics with monadic prdicate and many successors (SnS)
- 2nd-order logics with monadic and partial-order

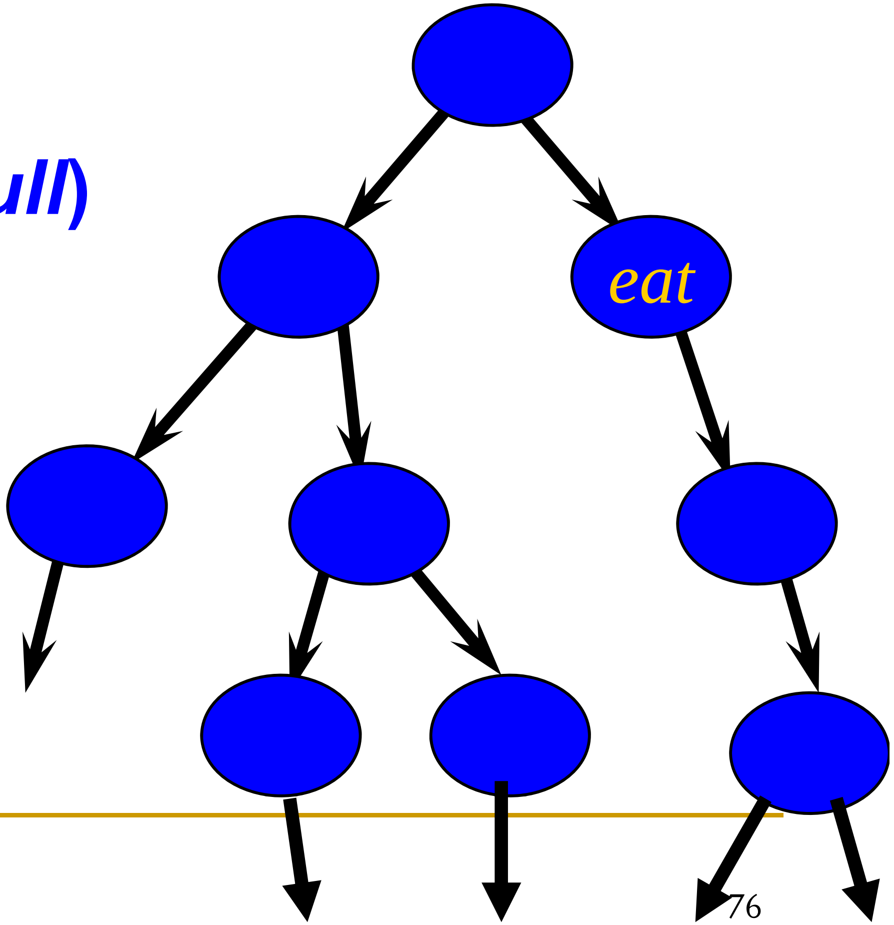
Very little known at the moment,

the fine difference in semantics of branching-structures

Expressiveness - CTL*, example (I)

A tree the distinguishes the following two formulas.

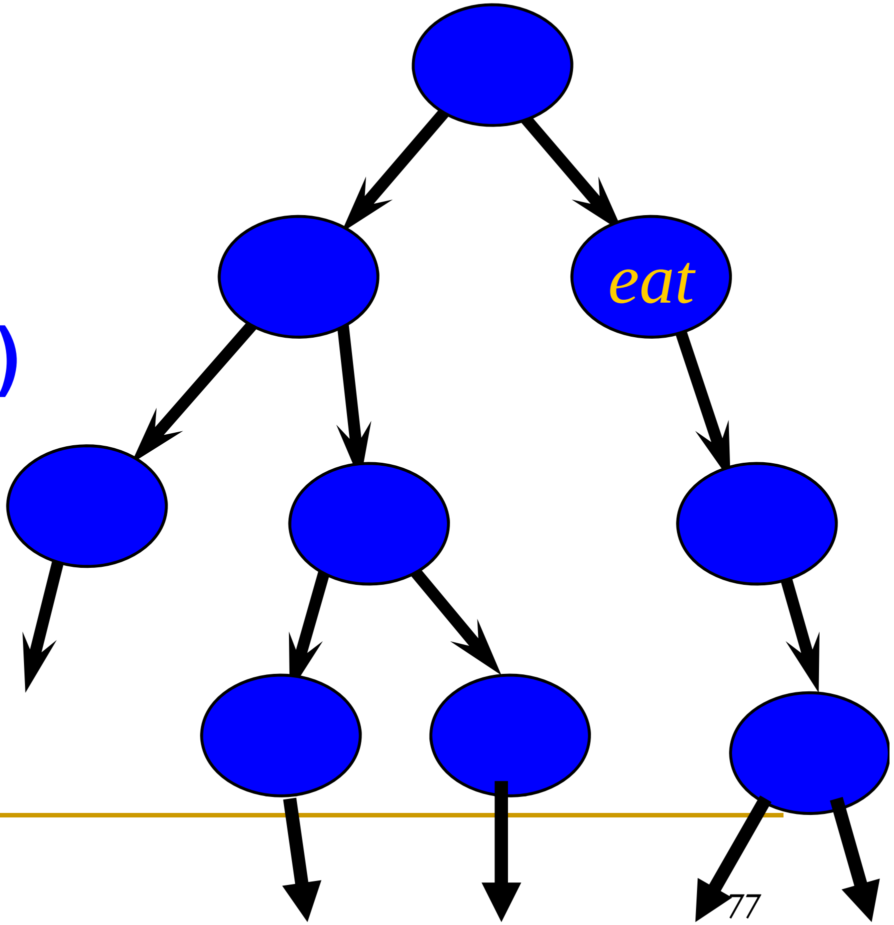
- $\forall((\Diamond eat) \rightarrow \Diamond full)$
  - **Negation:** $\exists((\Diamond eat) \wedge \Box\neg full)$
- $(\forall\Diamond eat) \rightarrow (\forall\Diamond full)$



Expressiveness
- CTL*, example (II)

A tree that distinguishes the following two formulas.

- $\forall((\Box \mathit{eat}) \rightarrow \Diamond \mathit{full})$
- $\forall\Box\ (\mathit{eat} \rightarrow \forall\Diamond \mathit{full})$
  - **Negation:** $\exists\Diamond(\mathit{eat} \wedge \exists\Diamond\neg \mathit{full})$



Expressiveness
- CTL*

With the abundant semantics in CTL*, we can compare the subclasses of CTL*.

With restrictions on the modal operations after $\exists$, $\forall$, we have many CTL$*$ subclasses.

Example:

$B(\neg,\vee,\bigcirc,\boldsymbol{U})$: only $\neg,\vee,\bigcirc,\boldsymbol{U}$ *after* $\exists, \forall$

$B(\neg,\vee,\bigcirc,\Diamond^{\infty})$: only $\neg,\vee,\bigcirc,\Diamond^{\infty}$ after $\exists, \forall$

$B(\bigcirc,\Diamond)$: only $\bigcirc,\Diamond$ after $\exists, \forall$

Expressiveness
- CTL*

CTL* subclass expressiveness heirarchy

$\mathrm{CTL}^* \quad > \quad B(\neg, \vee, \bigcirc, \Diamond, \mathcal{U}, \overset{\infty}{\Diamond})$

$\quad > \quad B(\bigcirc, \Diamond, \mathcal{U}, \overset{\infty}{\Diamond})$

$\quad > \quad B(\neg, \vee, \bigcirc, \Diamond, \mathcal{U})$

$\quad = \quad B(\bigcirc, \Diamond, \mathcal{U})$

$\quad > \quad B(\neg, \vee, \bigcirc, \Diamond)$

$\quad > \quad B(\bigcirc, \Diamond)$

$\quad > \quad B(\Diamond)$

Expressiveness
- CTL*

Some theorems :

- $B(\neg, \vee, \bigcirc, \Diamond, \boldsymbol{U}) \equiv B(\bigcirc, \Diamond, \boldsymbol{U})$

- $\exists \Diamond^{\infty} \boldsymbol{p}$ **is inexpressible in** $B(\bigcirc, \Diamond, U)$ **.**

Expressiveness
- CTL*

Comparing PLTL with CTL*

assumption, all $\varphi \in$ PLTL are interpreted as $\forall\varphi$

Intuition: PLTL is used to specify all runs of a system.

$$\begin{tikzcd}
 & \mathrm{B}(\neg,\vee,\bigcirc,\Diamond,\mathcal{U},\Diamond^{\infty}) \arrow[dr, ">"] & \\
\mathrm{CTL}^* \arrow[ur, ">"] \arrow[dr, ">"'] & & \mathrm{PLTL(F)} \\
 & \mathrm{PLTL} \arrow[ur, ">"'] &
\end{tikzcd}$$

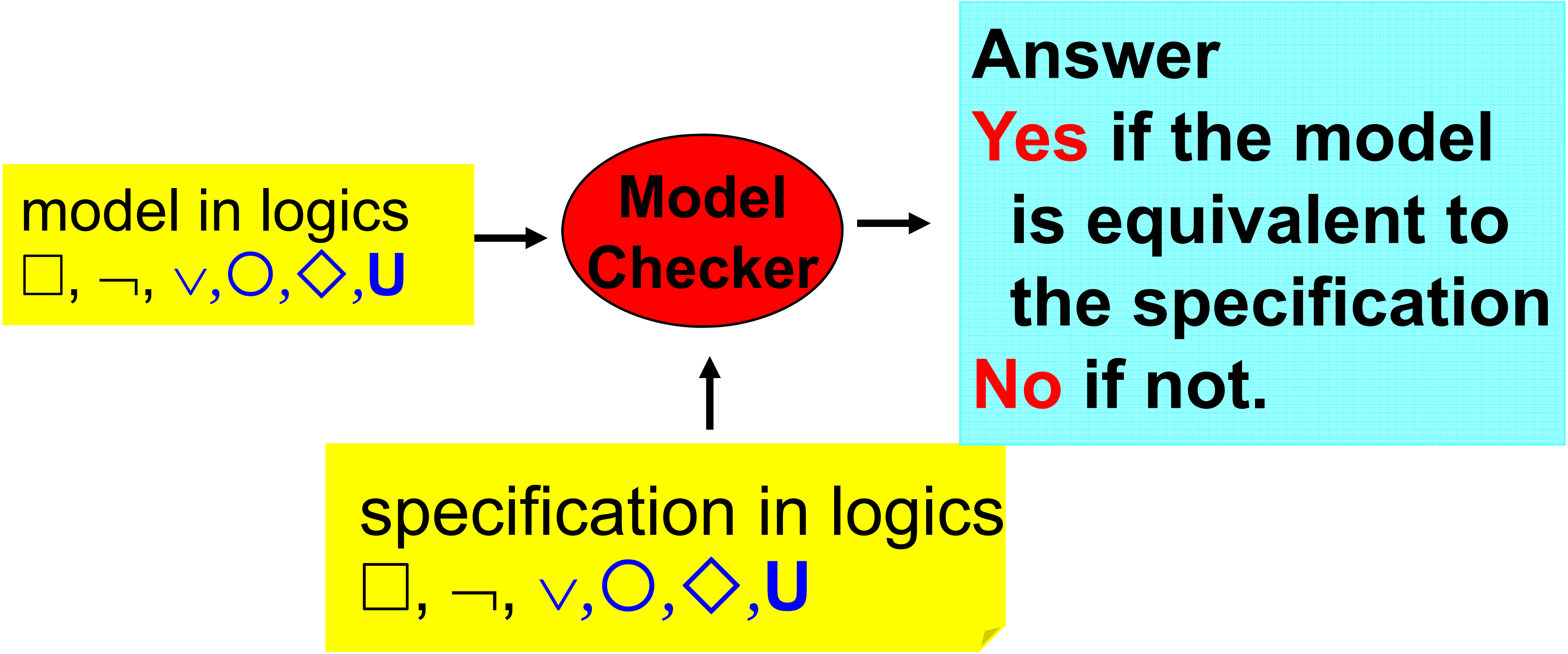
Verification

model (system) formula

specification formula

- LPTL, validity checking $\psi \models \phi$
  - instead, check the satisfiability of $\psi \wedge \neg\phi$
  - construct a tabelau for $\psi \wedge \neg\phi$
- model-checking $M \models \phi$
  - LPTL: M: a Büchi automata, ϕ: an LPTL formula
  - CTL: M: a finite-state automata, ϕ: a CTL formula
- simulation & bisimulation checking M ⊨ M’

Satisfiability-checking framework



model in logics
□, ¬, ∨,○,◇,**U**

Model Checker

specification in logics
□, ¬, ∨,○,◇,**U**

Answer

Yes if the model is equivalent to the specification

No if not.

LPTL
- tableau for satisfiability checking

Tableau for φ

- a finite Kripke structure that fully describes the behaviors of φ
- exponential number of states
- An algorithm can explore a fulfilling path in the tableau to answer the satisfiability.
  - nondeterministic
  - without construction of the tableau
  - PSPACE.

LPTL
- tableau for satisfiability checking

Tableau construction

a preprocessing step: push all negations to the literals.

- $\neg(\psi_1 \wedge \psi_2) \equiv (\neg \psi_1) \vee (\neg \psi_2)$
- $\neg(\psi_1 \vee \psi_2) \equiv (\neg \psi_1) \wedge (\neg \psi_2)$
- $\neg \bigcirc \psi \equiv \bigcirc \neg \psi$
- $\neg \neg \psi \equiv \psi$
- $\neg(\psi_1 \mathbf{U} \psi_2) \equiv (\Box \neg \psi_2) \vee ((\neg \psi_2) \mathbf{U} ((\neg \psi_1) \wedge (\neg \psi_2)))$
- $\neg \Box \psi \equiv \Diamond \neg \psi$

LPTL
- tableau for satisfiability checking

Tableau construction

CL(φ) (closure) is the smallest set of formulas containing φ with the following consistency requirement.

- $\neg p \in CL(\varphi)$ iff $p \in CL(\varphi)$
- If $\psi_1 \vee \psi_2$, $\psi_1 \wedge \psi_2 \in CL(\varphi)$, then $\psi_1, \psi_2 \in CL(\varphi)$
- If $\bigcirc\ \psi \in CL(\varphi)$, then $\psi \in CL(\varphi)$
- If $\psi_1 \mathbf{U} \psi_2 \in CL(\varphi)$, then ψ_1 , ψ_2 , $\bigcirc\ (\psi_1\ \mathbf{U}\ \psi_2\ ) \in CL(\varphi)$
- If $\square\ \psi \in CL(\varphi)$, then ψ, $\bigcirc\ \square\ \psi \in CL(\varphi)$

LPTL

- tableau for satisfiability checking

Tableau (V, E), *node consistency condition*:

A tableau node $v \in V$ is a set $v \subseteq CL(f)$ such that

- $p \in v$ iff $\neg p \notin v$
- If $\psi_1 \vee \psi_2 \in v$, then $\psi_1 \in v$ or $\psi_2 \in v$
- If $\psi_1 \wedge \psi_2 \in v$, then $\psi_1 \in v$ and $\psi_2 \in v$
- if $\Box \psi \in v$, then $\psi \in v$ and $\bigcirc \Box \psi \in v$
- if $\Diamond \psi \in v$, then $\psi \in v$ or $\bigcirc \Diamond \psi \in v$
- if $\psi_1 \mathbf{U} \psi_2 \in v$, then $\psi_2 \in v$ or ($\psi_1 \in v$ and $\bigcirc (\psi_1 \mathbf{U} \psi_2) \in v$)

LPTL
- tableau for satisfiability checking

Tableau (V, E), *arc consisitency condition*:

Given an arc (v,v') $\in$E, if $\bigcirc\ \psi \in$ v, then $\psi \in$ v'

- A node v in (V,E) is initial for φ if $\varphi\ \in$v.

LPTL

- tableau for satisfiability checking

$CL(pUq) = \{pUq, \bigcirc pUq, p, \neg p, q, \neg q\}$

Example: (p U q)

tableau (V,E)

V: $\{p, q, pUq, \bigcirc pUq\}$ $\{p, q, \bigcirc pUq\}$ $\{p, q\}$

$\{p, q, pUq\}$

$\{p, \neg q, pUq, \bigcirc pUq\}$ $\{p, \neg q, \bigcirc pUq\}$ $\{p, \neg q\}$

$\{\neg p, q, pUq, \bigcirc pUq\}$ $\{\neg p, q, pUq\}$ $\{\neg p, q\}$

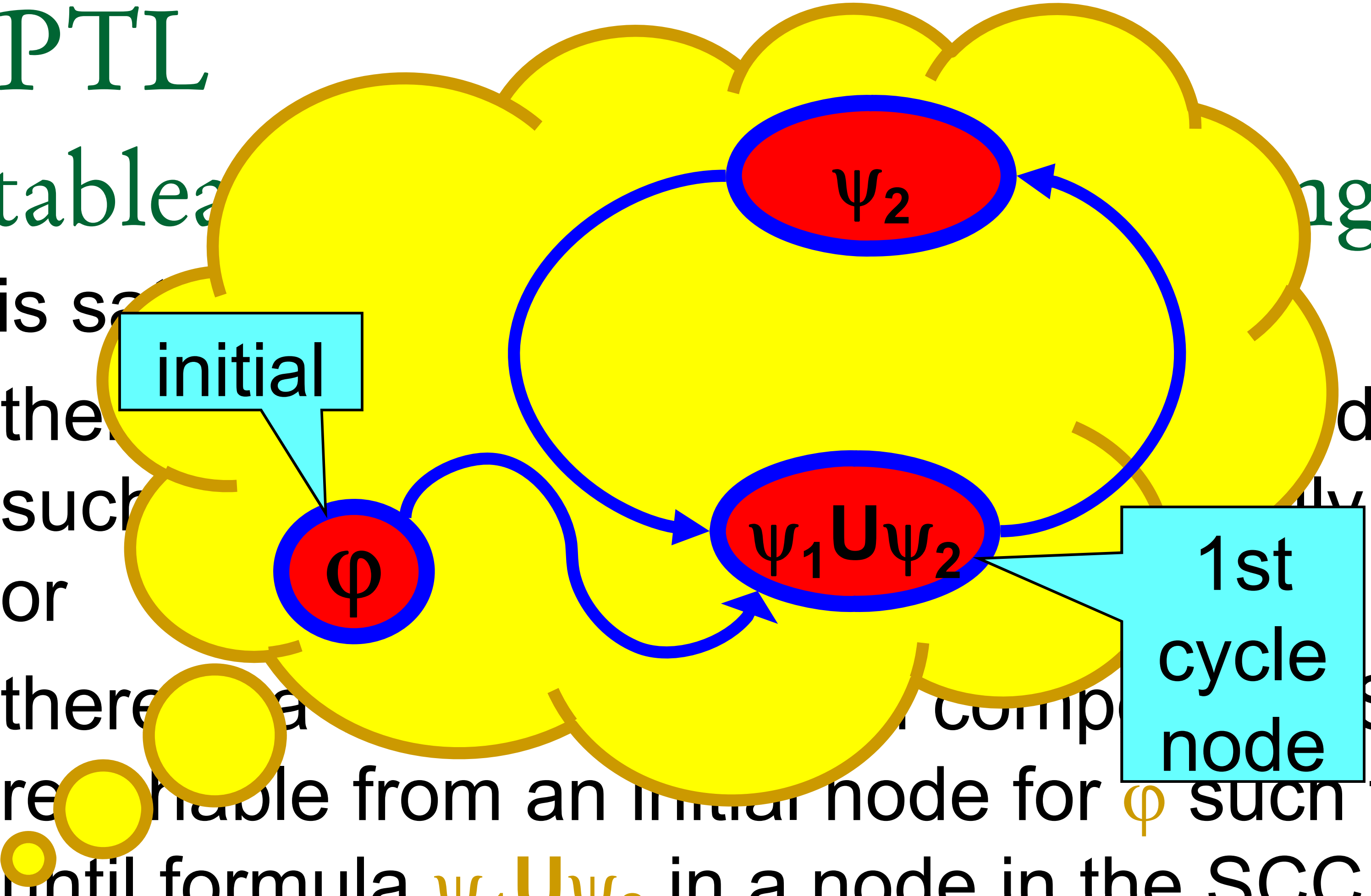
$\{\neg p, q, \bigcirc pUq\}$

$\{\neg p, \neg q, \bigcirc pUq\}$ $\{\neg p, \neg q\}$

E: ?

LPTL

- tablea ng



φ is sa

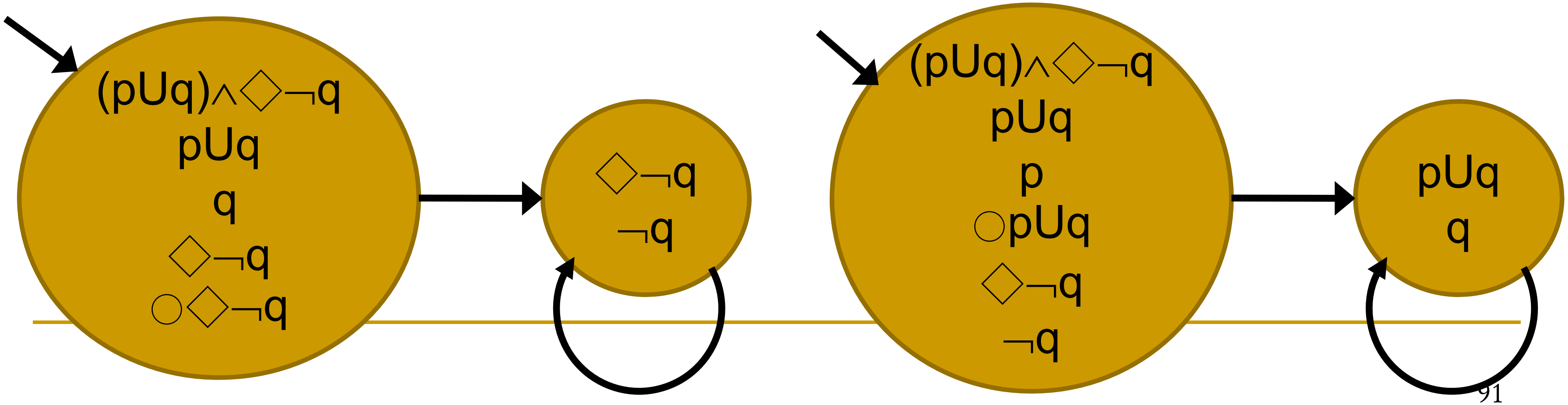
- the de for φ such lly satisfied; or
- there a compo SCC) re hable from an node for φ such that for all until formula $\psi_1 \mathbf{U} \psi_2$ in a node in the SCC, there is also a node in the SCC containing ψ_2 ; or
- there is a cycle reachable from an initial node for φ such that the for all until formulas $\psi_1 \mathbf{U} \psi_2$ in **the first cycle node**, there is also a node in the cycle containing ψ_2 .

LPTL
- tableau for satisfiability checking

Please use tableau method to show that
$pUq \models \Box q$ is false.

1) Convert to negation: $(pUq) \wedge \Diamond\neg q$

2) $CL((pUq) \wedge \Diamond\neg q)$
$= \{(pUq) \wedge \Diamond\neg q,\ pUq,\ \bigcirc pUq,\ p,\ q,\ \Diamond\neg q,\ \bigcirc\Diamond\neg q\ \}$



LPTL
- tableau for satisfiability checking

Please use tableau method to show that $pUq \models \Diamond q$ is true.

1) Convert to negation: $(pUq) \wedge \Box \neg q$

2) $CL((pUq) \wedge \Box \neg q)$

$= \{(pUq) \wedge \Box \neg q,\ pUq,\ \bigcirc pUq,\ p,\ q,\ \Box \neg q,\ \bigcirc \Box \neg q\ \}$

Pf: In each path that is a model of $(pUq) \wedge \Box \neg q$, q must always be satisfied. Thus, pUq is never fulfilled in the model.

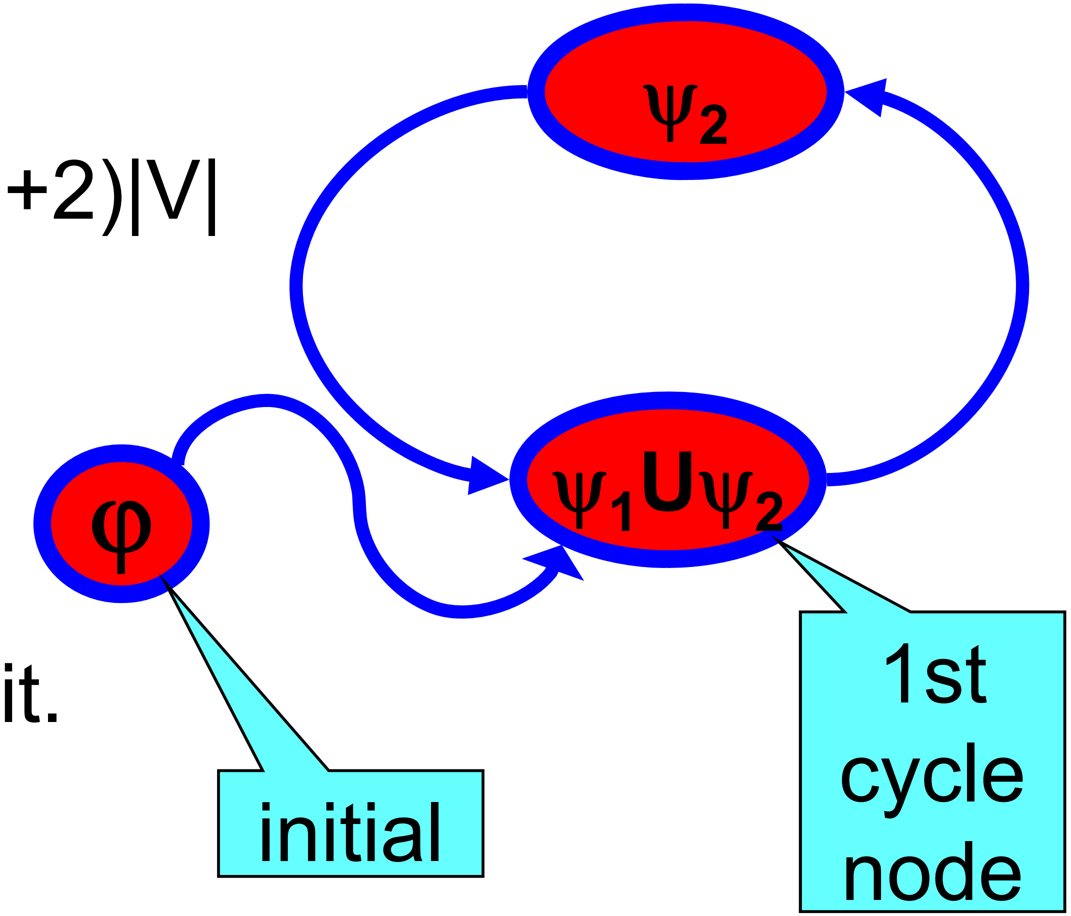
QED

LPTL
- tableau for satisfiability checking

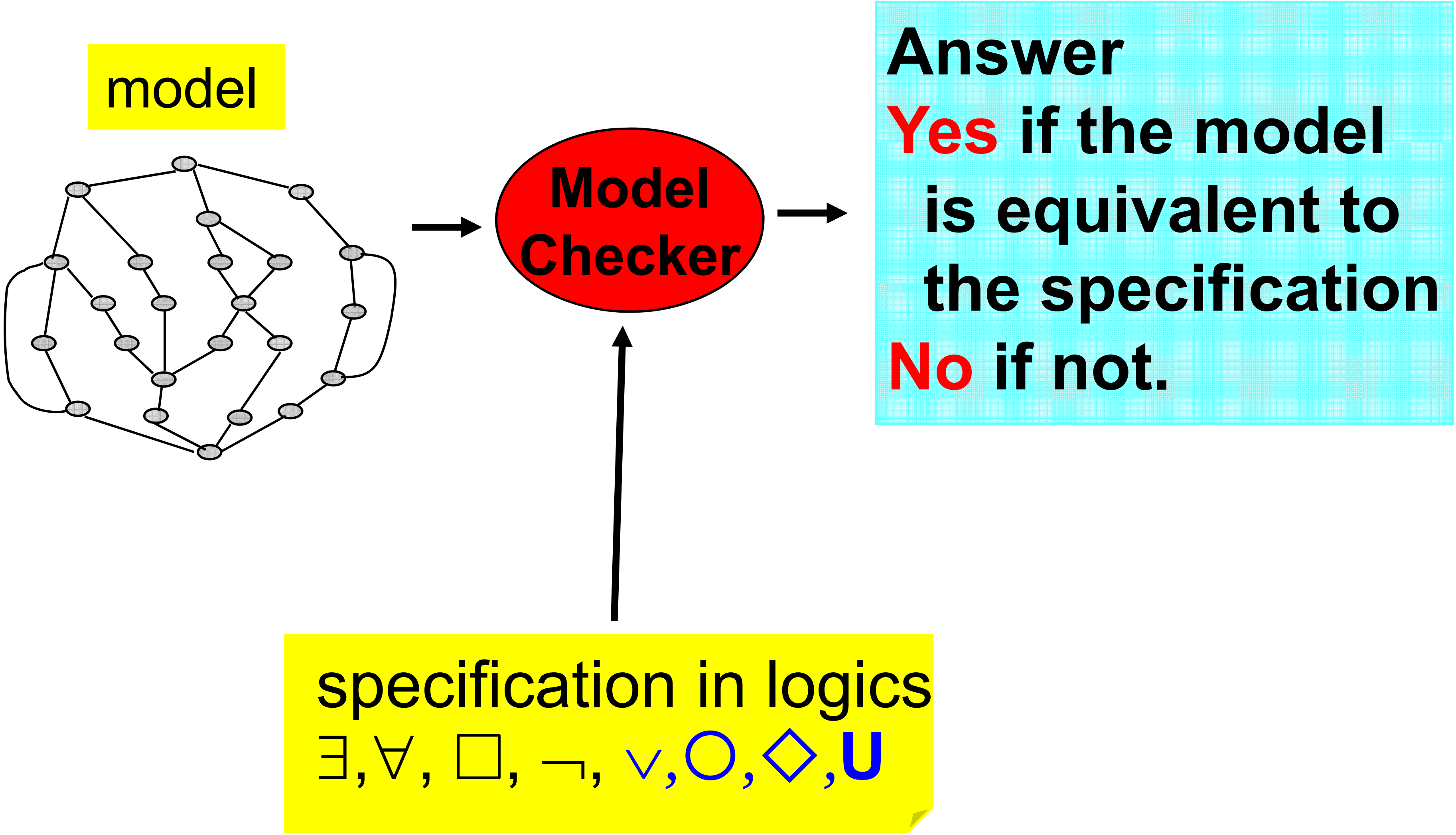
φ is satisfiable iff in (V,E),

there exists ...

- path+cycle≤ (|CL(φ)|+2)|V|
- |CL(φ)| flags to check the until-formulas from the first cycle node.
- nondeterministic PSPACE can solve it.
- PSPACE-complete.



CTL model-checking framework



model

Model Checker

Answer

Yes if the model is equivalent to the specification

No if not.

specification in logics

∃,∀, □, ¬, ∨,○,◇,**U**

CTL
- model-checking

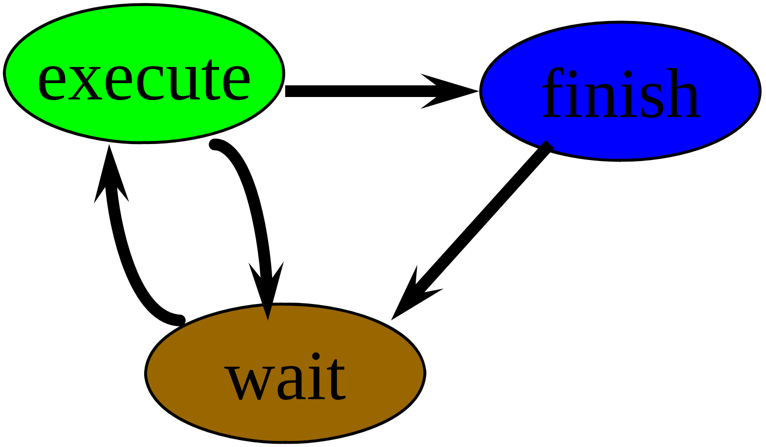
Given a finite Kripke structure M and a CTL formula φ, is M a model of φ ?

- usually, M is a finite-state automata.
- PTIME algorithm.
- When M is generated from a program with variables, its size is easily exponential.

CTL
- model-checking algorithm

techniques

- state-space exploration
  - state-spaces represented as finite Kripke structure
    - directed graph
    - nodes: states or possible worlds
    - arcs: state transitions
- regular behaviors



- Usually the state count is astronomical.

Kripke structure
- *Least fixpoint in modal logics*

Dark-night murder, strategy I:

A suspect will be in the 2nd round iff

- *He/she lied to the police in the 1st round; or*
- *He/she is loyal to someone in the 2nd round*

What is the minimal solution to *2nd[]* ?

$$2nd[i] \equiv Liar[i] \vee \exists j{\neq}i(2nd[j] \wedge Loyal\text{-}to[i,j])$$

Kripke structure

- *Least fixpoint in modal logics*

In a dark night, there was a cruel murder.

- n suspects, numbered 0 through n-1.
- *Liar[i]* iff suspect i has lied to the police in the 1st round investigation.
- *Loyal-to[i,j]* iff suspect i is loyal to suspect j in the same criminal gang.
- *2nd[i]* iff suspect i to be in 2nd round investigation.

What is the minimal solution to *2nd[]* ?

Kripke structure

- *Greatest fixpoint in modal logics*

In a dark night, there was a cruel murder.

- n suspects, numbered 0 through n-1.
- *¬Liar[i]* iff the police cannot prove suspect i has lied to the police in the 1st round investigation.
- *Loyal-to[i,j]* iff suspect i is loyal to j are in the same criminal gang.
- *2nd[i]* iff suspect i to be in 2nd round investigation.

What is the maximal solution to ¬ *2nd[]* ?

Kripke structure
- *Greatest fixpoint in modal logics*

Dark-night murder, strategy II

A suspect will not be in the 2nd round iff

- *We cannot prove he/she has lied to the police; and*
- He/she is loyal to someone not in the 2nd round.

What is the maximal solution to $\neg\,\mathit{2nd}[]$?

$$\neg\,\mathit{2nd}[i] \equiv \neg \mathit{Liar}[i] \wedge \exists j \neq i (\neg \mathit{2nd}[j] \wedge \mathit{Loyal\text{-}to}[i,j])$$

In comparison:

$\neg\,\mathit{2nd}[i] \equiv \neg \mathit{Liar}[i] \wedge \forall j \neq i (\neg \mathit{2nd}[j] \wedge \mathit{Loyal\text{-}to}[i,j])$

~~$\neg\,\mathit{2nd}[i] \equiv \neg \mathit{Liar}[i] \wedge \forall j \neq i (\neg \mathit{2nd}[j] \rightarrow \mathit{Loyal\text{-}to}[i,j])$~~

$\neg\,\mathit{2nd}[i] \equiv \neg \mathit{Liar}[i] \wedge \forall j \neq i (\mathit{Loyal\text{-}to}[i,j] \rightarrow \neg \mathit{2nd}[j]\,)$

Safety analysis

Given M and p (safety predicate), do all states reachable from initial states in M satisfy p ?

- In model-checking:

  Is M a model of $\forall\Box p$?

- Or in risk analysis: Is there a state reachable from initial states in M satisfy p ?

  $$\forall\Box p \equiv \neg\exists\Diamond\neg p \equiv \neg\exists \mathrm{true}\ \mathrm{U}\neg p$$

Reachability analysis: $\exists\Diamond\eta$

Is there a state s reachable from another state s'?

- Encode risk analysis
- Encode the complement of safety analysis
- Most used in real applications

Kripke structure
- safety analysis

Reachability algorithm in graph theory

Given

- a Kripke structure $A = (S, S_0, R, L)$
- a safety predicate **η**,

find a path from a state in S_0 to a state in $[\neg\eta]$.

Solutions in graph theory

- Shortest distance algorithms
- spanning tree algorithms

Kripke structure - safety analysis

```
/* Given A = $(S, S_0, R, L)$*/
safety_analysis($\eta$) /* using least fixpoint algorithm */ {
    for all s, if $\neg\eta \in L(s)$, $L(s)=L(s)\cup\{\exists\Diamond\neg\eta\}$;
    repeat {
        for all s, if $\exists(s,s')(\exists\Diamond\neg\eta \in L(s'))$,
            $L(s)=L(s)\cup\{\exists\Diamond\neg\eta\}$;
    } until no more changes to L(s) for any s.
    if there is an $s_0 \in S_0$ with $\exists\Diamond\neg\eta \in L(s_0)$,
        return 'unsafe,'
        else return 'safe.'
}
```

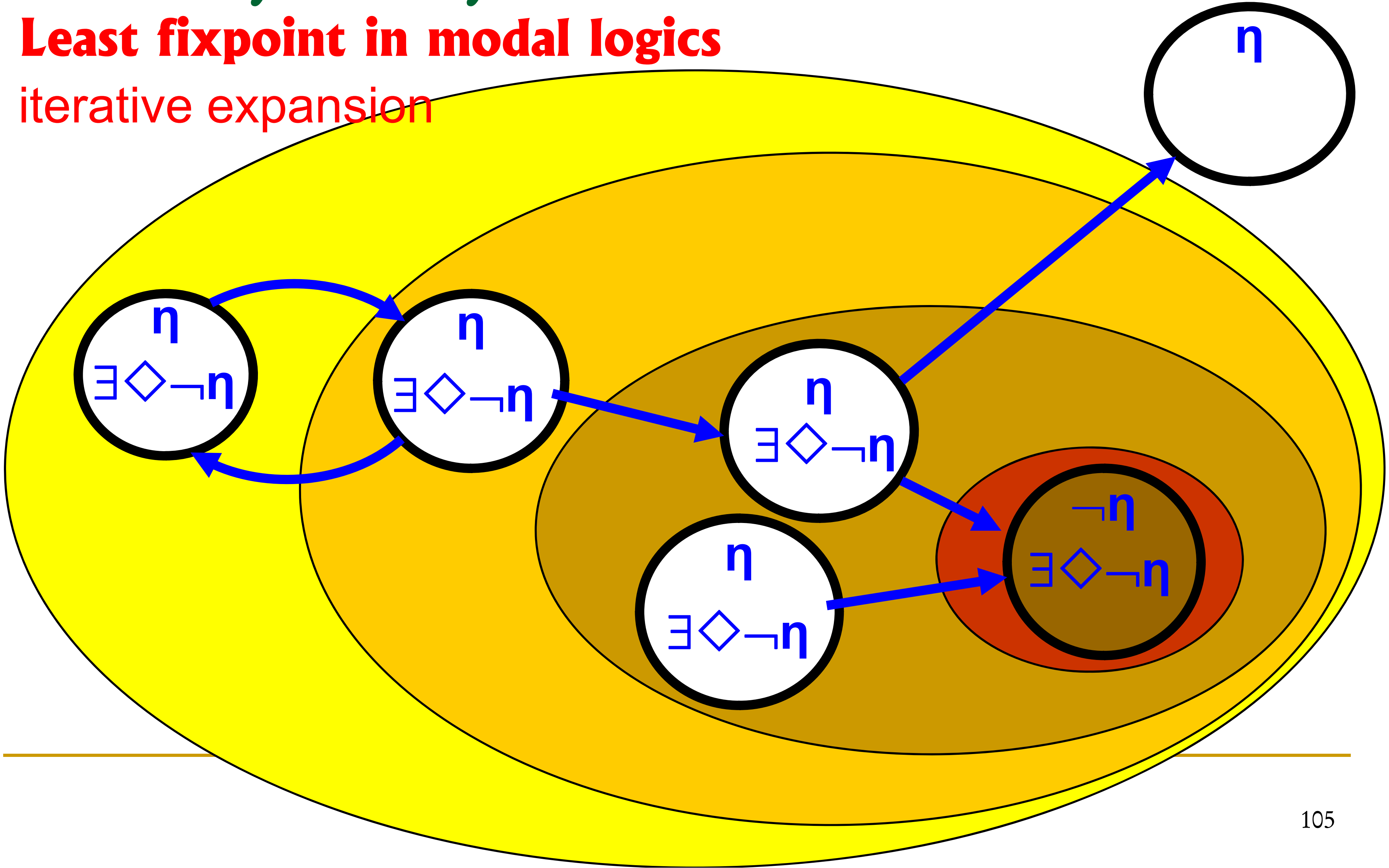
A notation for the possibility of $\neg\eta$

The procedure terminates since S is finite in the Kripke structure.

Kripke structure
- safety analysis

Least fixpoint in modal logics

iterative expansion



Kripke structure
- liveness analysis : $\forall\Diamond\ \eta$

Given

- a Kripke structure $A = (S, S_0, R, L)$
- a liveness predicate η,

can η be true eventually ?

Example:

Can the computer be started successfully ?

Will the alarm sound in case of fire ?

Kripke structure
- liveness analysis

Strongly connected component algorithm in graph theory

Given

- a Kripke structure $A = (S, S_0, R, L)$
- a liveness predicate η,

find a cycle such that

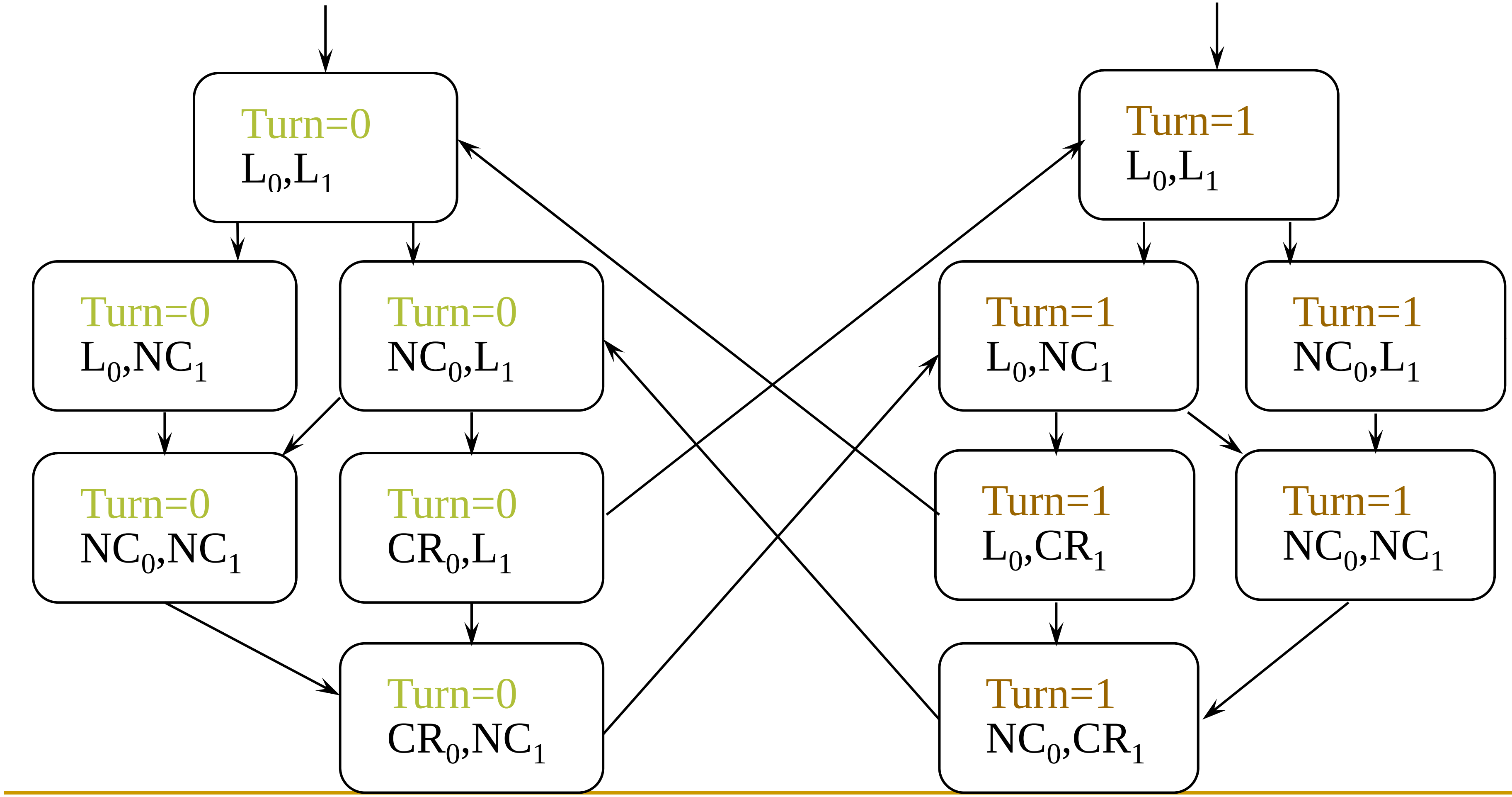
- all states in the cycle are $\neg\eta$
- there is a $\neg\eta$ path from a state in S_0 to the cycle.

Solutions in graph theory

- strongly connected components (SCC)

Kripke structure
- liveness analysis

$\Box\,(Turn=0 \rightarrow \Diamond\, Turn=1)$



Turn=0
L_0,L_1

Turn=0
L_0,NC_1

Turn=0
NC_0,L_1

Turn=0
NC_0,NC_1

Turn=0
CR_0,L_1

Turn=0
CR_0,NC_1

Turn=1
L_0,L_1

Turn=1
L_0,NC_1

Turn=1
NC_0,L_1

Turn=1
L_0,CR_1

Turn=1
NC_0,NC_1

Turn=1
NC_0,CR_1

Kripke structure
- liveness analysis

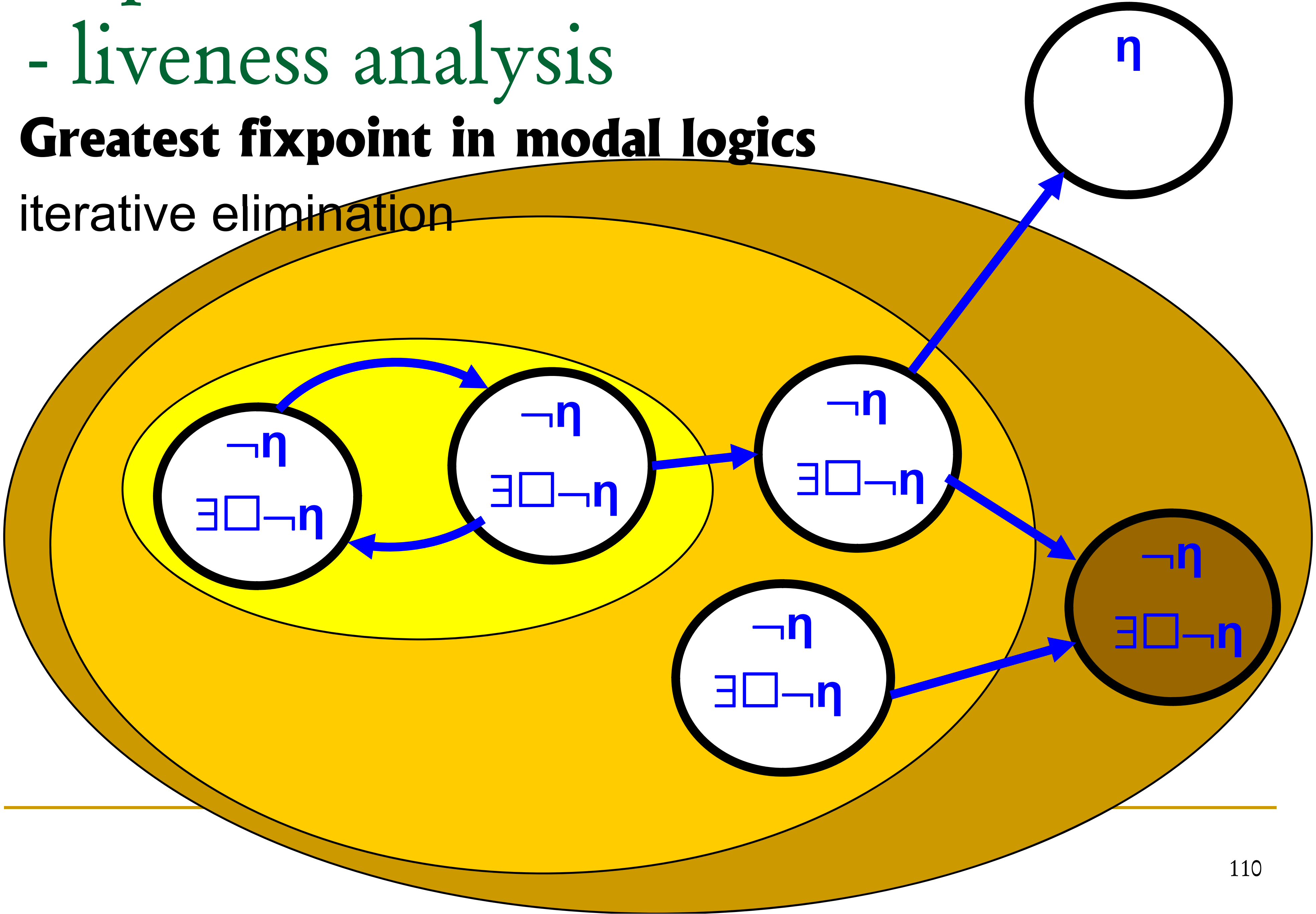
```
liveness($\eta$) /* using greatest fixpoint algorithm */ {
    for all s, if $\neg\eta \in L(s)$, $L(s)=L(s)\cup\{\exists\Box\neg\eta\}$;
    repeat {
        for all s, if $\exists\Box\neg\eta \in L(s)$ and $\forall(s,s')(\exists\Box\neg\eta \notin L(s))$,
            $L(s)=L(s) - \{\exists\Box\neg\eta\}$;
    } until no more changes to L(s) for any s.
    if there is an $s_0 \in S_0$ with $\exists\Box\neg\eta \in L(s_0)$,
        return 'liveness not true,'
        else return 'liveness true.'
}
```

The procedure terminates since S is finite in the Kripke structure.

Kripke structure
- liveness analysis

Greatest fixpoint in modal logics

iterative elimination



CTL model-checking

The NORMAL form needed in CTL model-checking:

1. only modal operators

$$\exists\bigcirc\varphi,\ \exists\ \psi_1\,\mathbf{U}\psi_2\,,\ \exists\Box\varphi$$

2. No modal operators

$$\forall\bigcirc\varphi,\ \forall\ \psi_1\,\mathbf{U}\psi_2\,,\ \forall\Box\varphi,\ \forall\Diamond\varphi,\ \exists\Diamond\varphi$$

3. No double negation: $\neg\neg\varphi$

4. No implication: $\psi_1\Rightarrow\psi_2$

CTL
- model-checking algorithm (1/6)

Given M and φ,

1. Convert φ to NORMAL form.

2. list the elements in Cl (φ) according to their sizes

$$\varphi_0\ \varphi_1\ \varphi_2 \cdots \varphi_n$$

for all $0 \le i < j \le n$，φ_j is not a subformula of φ_i

2. for i=0 to n,

label (φ_i)

See next page!

3. for all initial states s_0 of M, if $\varphi \notin L(s_0)$, return \`No!’

4. return \`Yes!’

CTL
- model-checking algorithm (2/6)

```
label($\varphi$ ) {
case p, return;
case $\neg\varphi$, for all s, if $\varphi \notin L(s)$, $L(s) = L(s) \cup \{\neg\varphi\}$
case$\varphi \vee \psi$, for all s, if$\varphi \in L(s)$ or$\psi \in L(s)$,
    $L(s)=L(s) \cup \{\varphi \vee \psi\}$
case $\exists \bigcirc \varphi$, for all s, if $\exists (s,s')$ with $\varphi \in L(s')$,
    $L(s)=L(s) \cup \{\exists \bigcirc \varphi\}$
case $\exists\ \psi_1 \mathbf{U} \psi_2$ , lfp($\psi_1$ , $\psi_2$ );
case $\exists \Box \varphi$, gfp($\varphi$);
}
```

CTL
- model-checking algorithm (3/6)

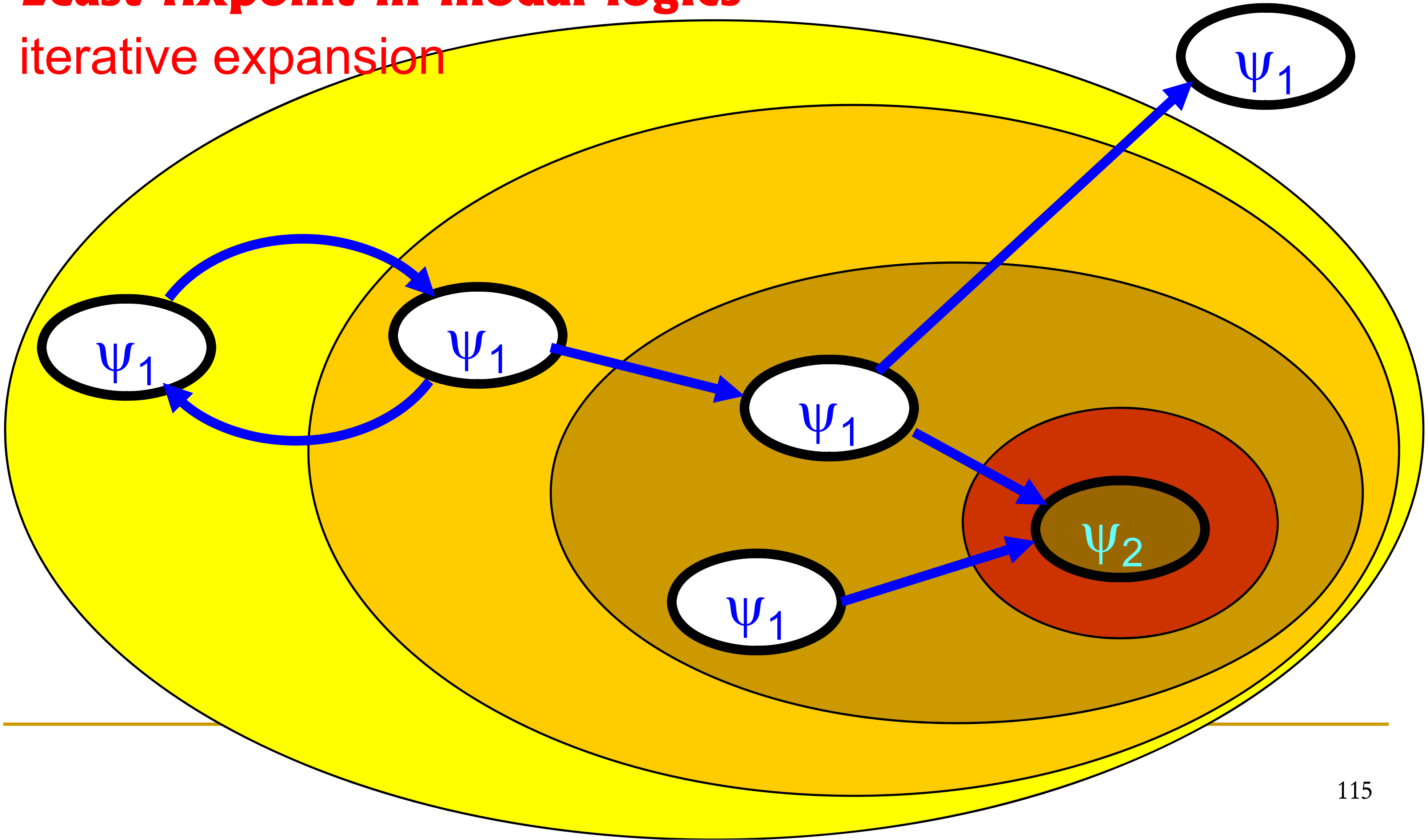
```
lfp($\psi_1$ , $\psi_2$ ) /* least fixpoint algorithm */ {
    for all s, if $\psi_2 \in L(s)$, $L(s)=L(s)\cup\{\exists\psi_1\mathbf{U}\psi_2\}$;
    repeat {
        for all s, if $\psi_1 \in L(s)$ and $\exists(s,s')(\exists\psi_1\mathbf{U}\psi_2 \in L(s'))$,
          $L(s)=L(s)\cup\{\exists\psi_1\mathbf{U}\psi_2\}$;
    } until no more changes to L(s) for any s.
}
```

The procedure terminates since S is finite in the Kripke structure.

CTL - model-checking algorithm (4/6)

Least fixpoint in modal logics

iterative expansion



CTL
- model-checking algorithm (5/6)

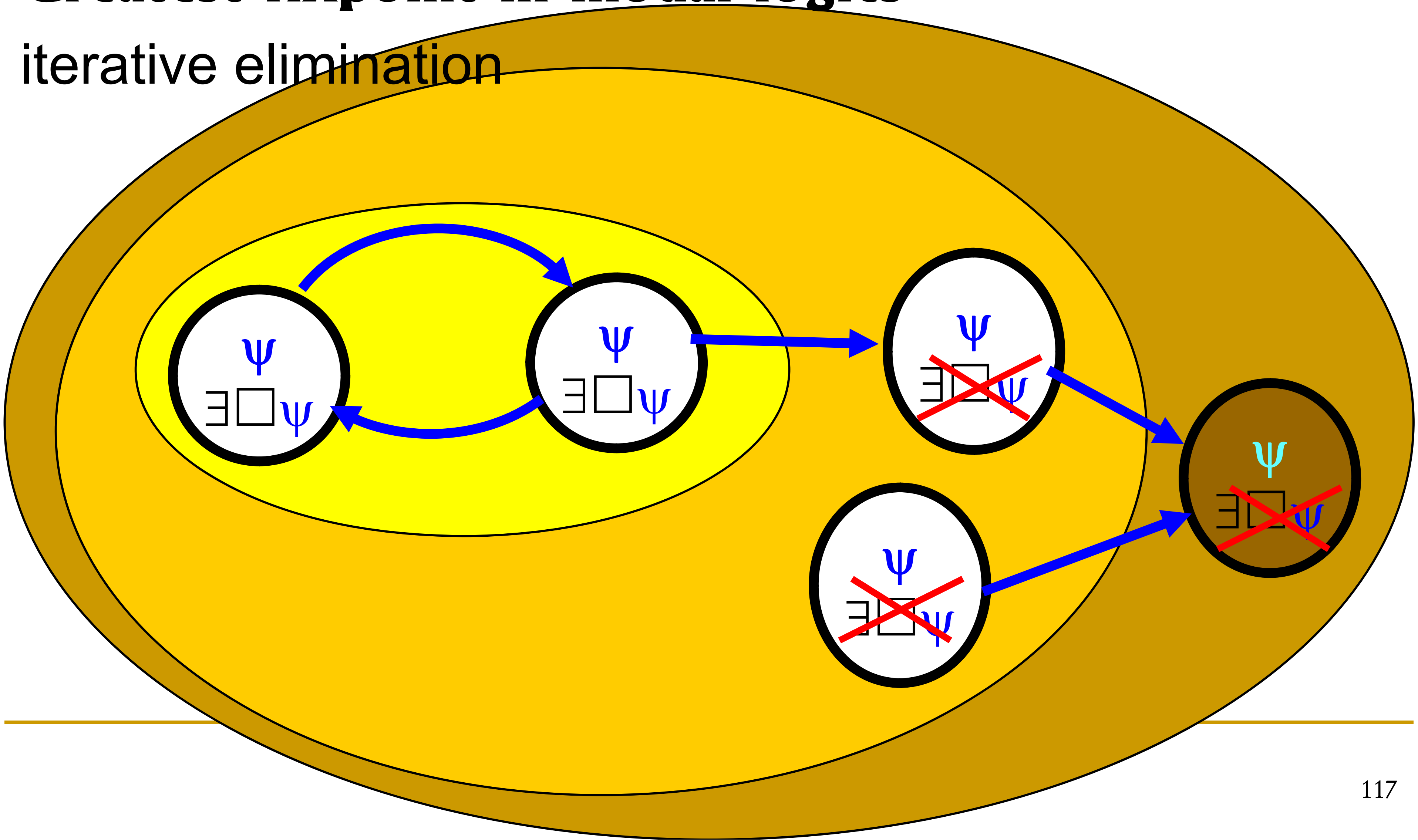
```
gfp($\psi$) /* greatest fixpoint algorithm */ {
    for all s, if $\psi \in L(s)$, $L(s)=L(s)\cup\{\exists\Box\psi\}$;
    repeat {
        for all s, if $\exists\Box\psi\in L(s)$ and $\forall(s,s')(\exists\Box\psi\notin L(s'))$,
            $L(s)=L(s) - \{\exists\Box\psi\}$;
    } until no more changes to L(s) for any s.
}
```

The procedure terminates since S is finite in the Kripke structure.

CTL
- model-checking algorithm (6/6)

Greatest fixpoint in modal logics

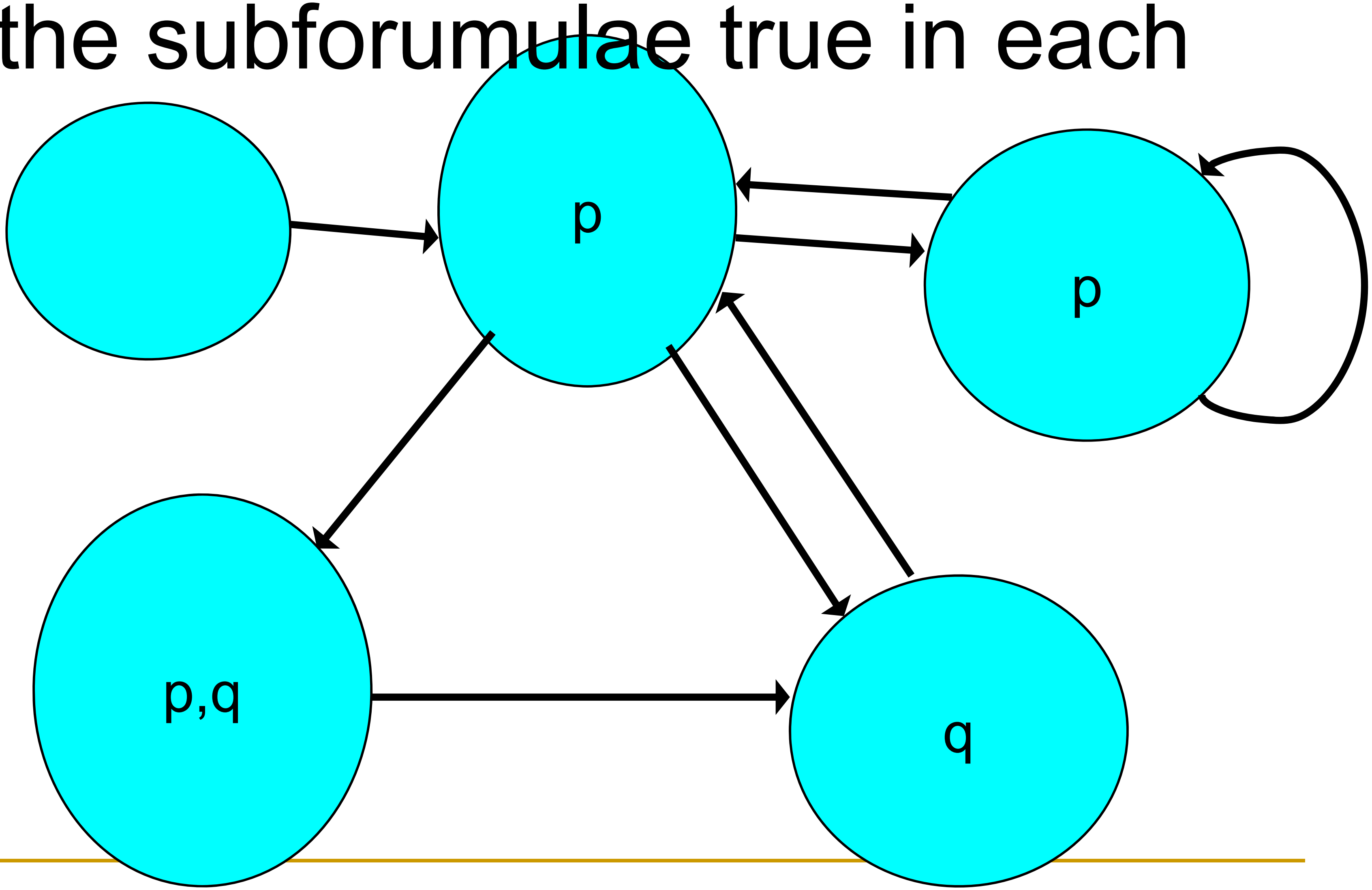
iterative elimination



$(\exists \bigcirc \exists p \, U q) \wedge \exists \Box p$

Labeling funciton:

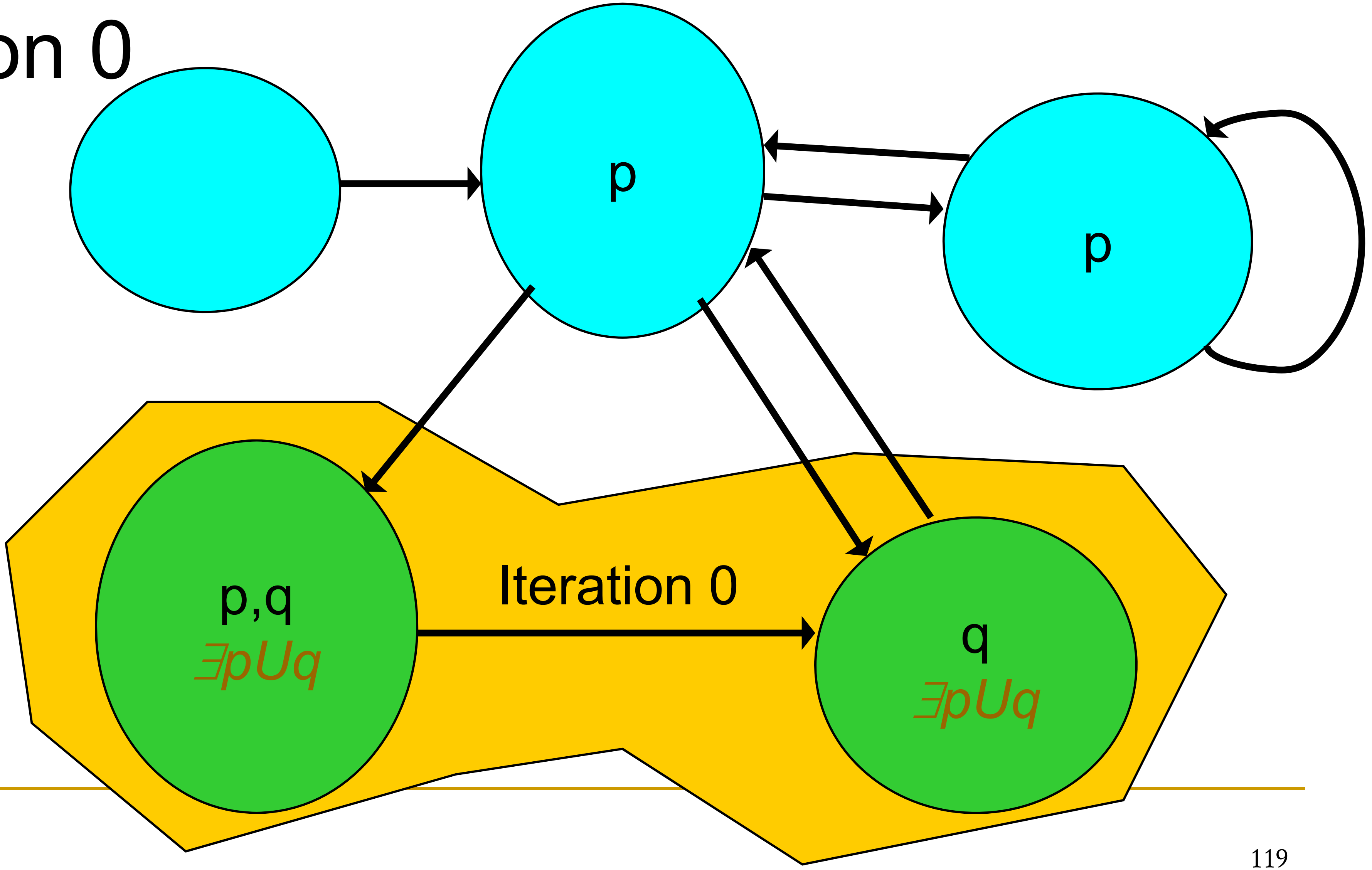
label the subforumulae true in each state.



$(\exists\bigcirc\exists p\,\mathcal{U}q) \wedge \exists\Box p$

Evaluating ∃pUq using least fixpoint

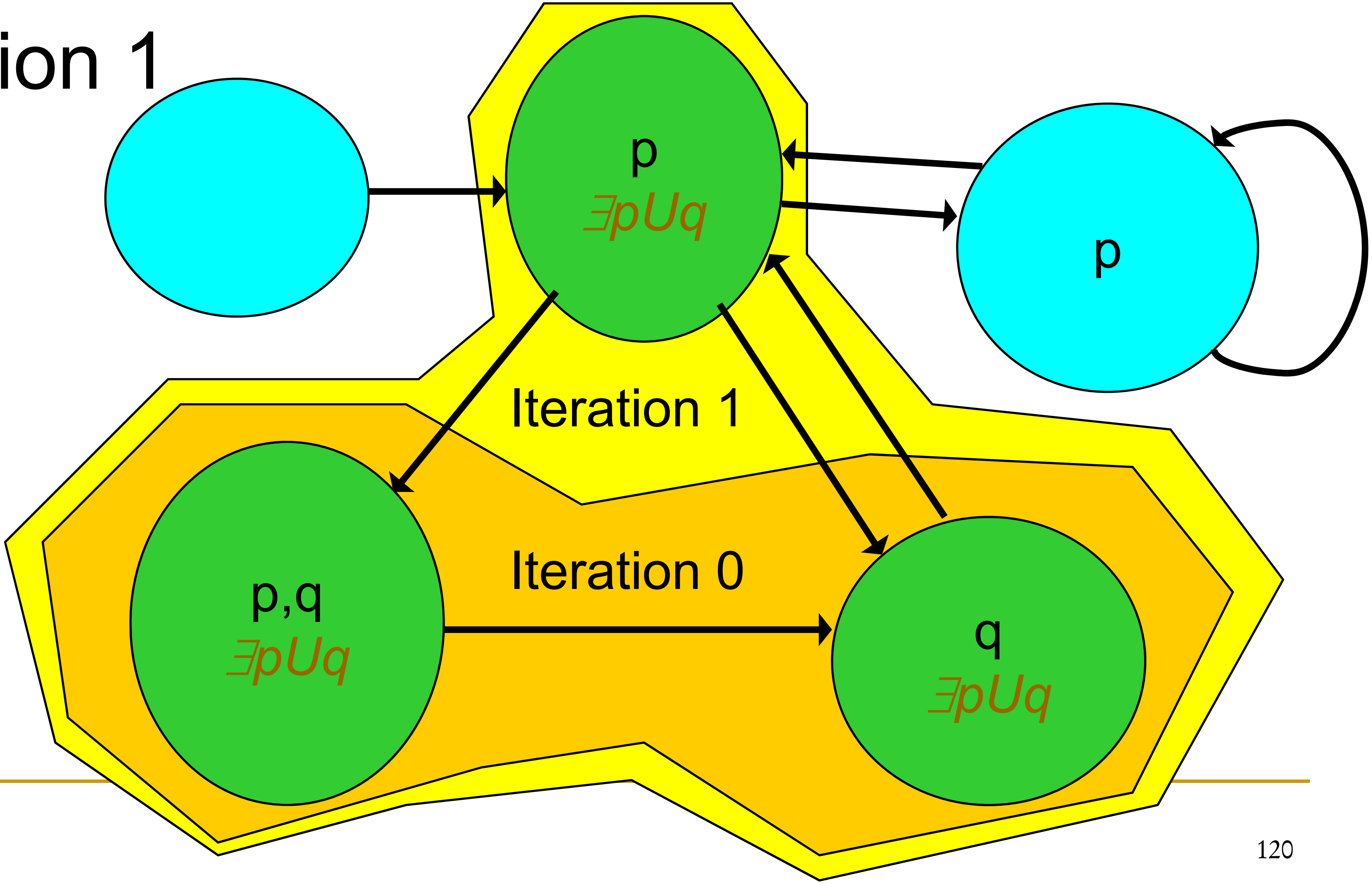
Iteration 0



$(\exists \bigcirc \exists p\, U q) \wedge \exists \Box p$

Evaluating ∃pUq using least fixpoint

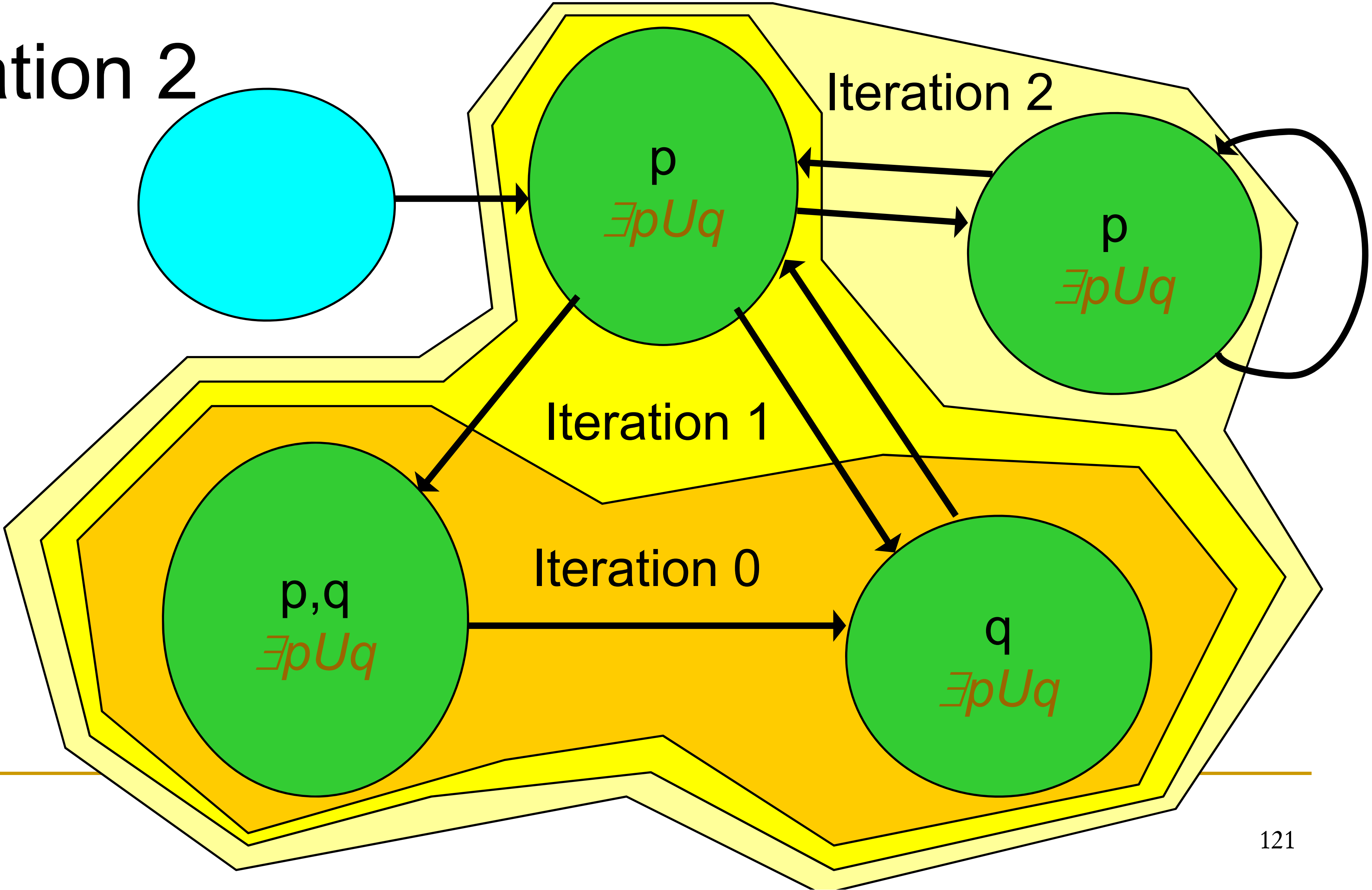
Iteration 1



$(\exists \bigcirc \exists p\, U q) \wedge \exists \Box p$

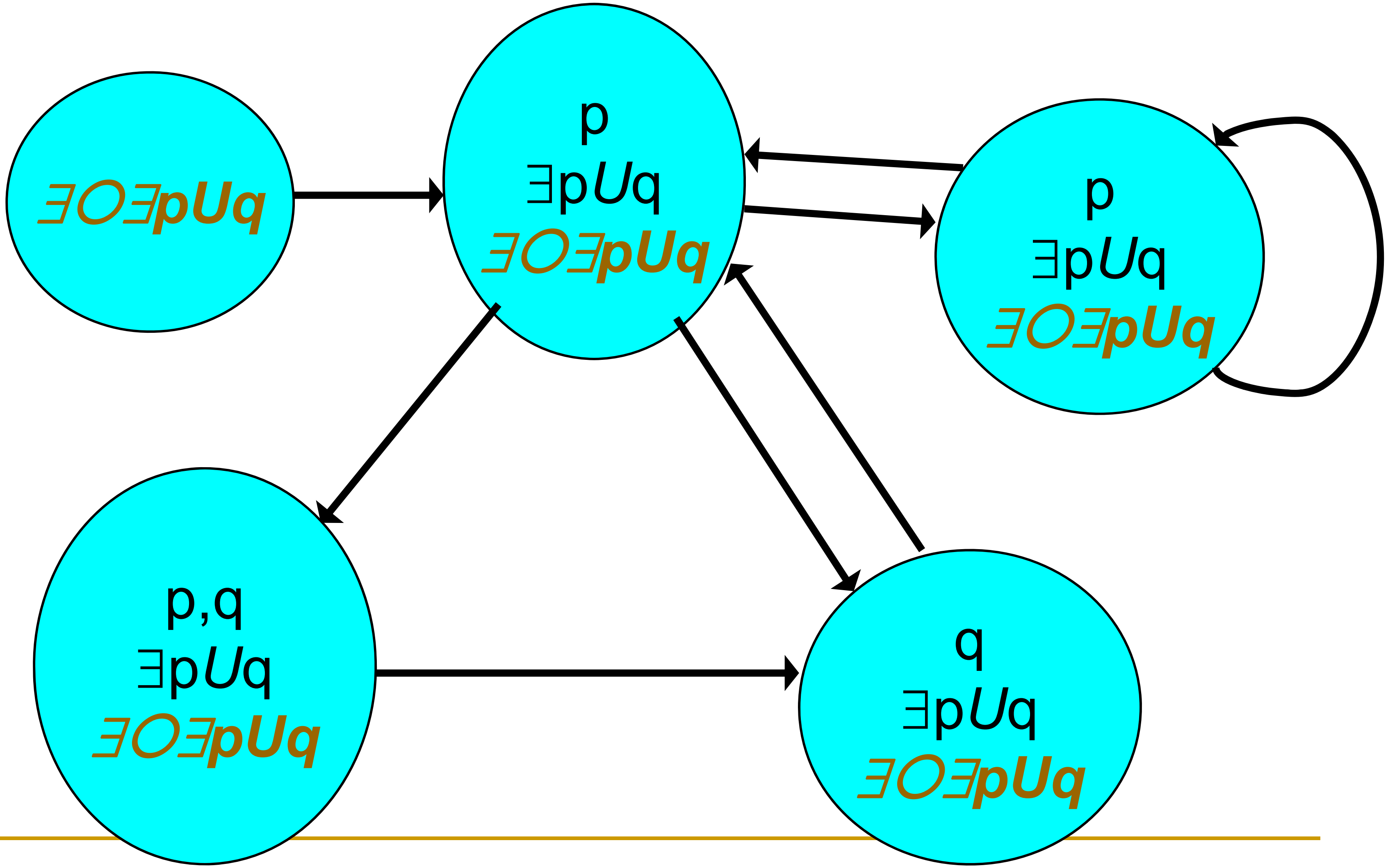
Evaluating ∃pUq using least fixpoint

Iteration 2



(∃○∃p*U*q) ∧ ∃□p

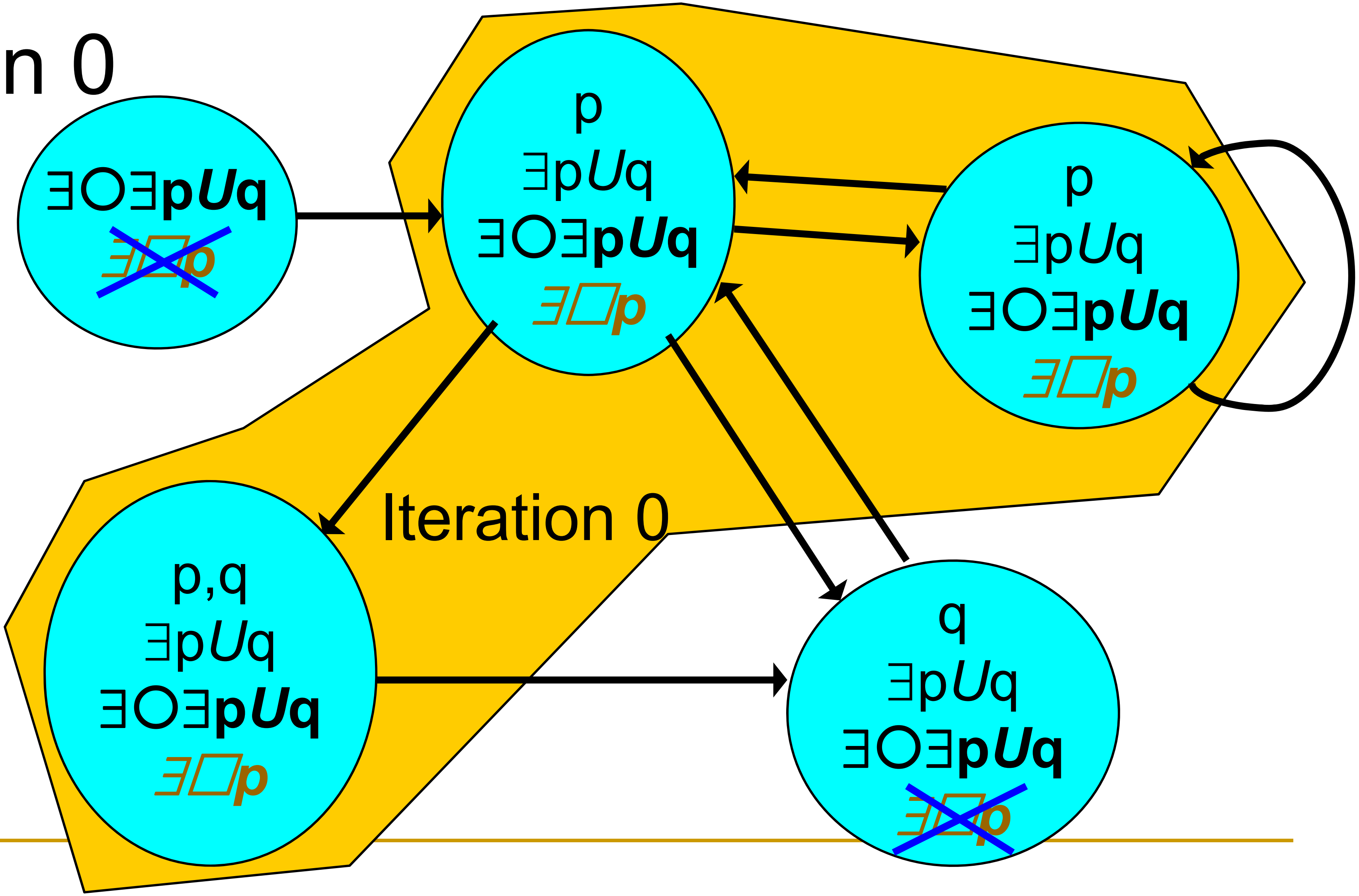
Evaluating ∃○∃pUq



$(\exists\bigcirc\exists p\,\mathcal{U}q) \wedge \exists\Box p$

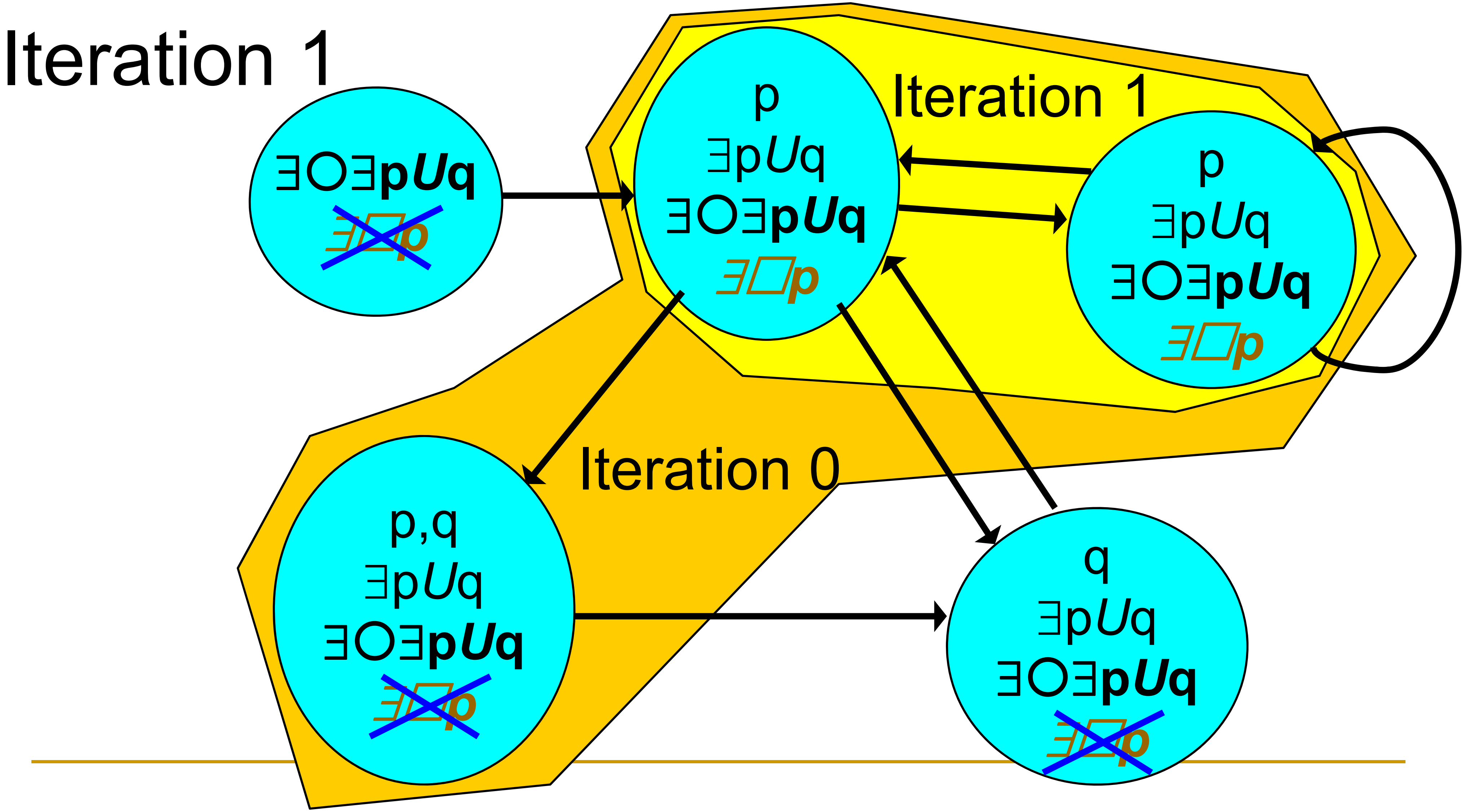
Evaluating $\exists\Box p$ using greatest fixpoint

Iteration 0



(∃○∃p*U*q) ∧ ∃□p

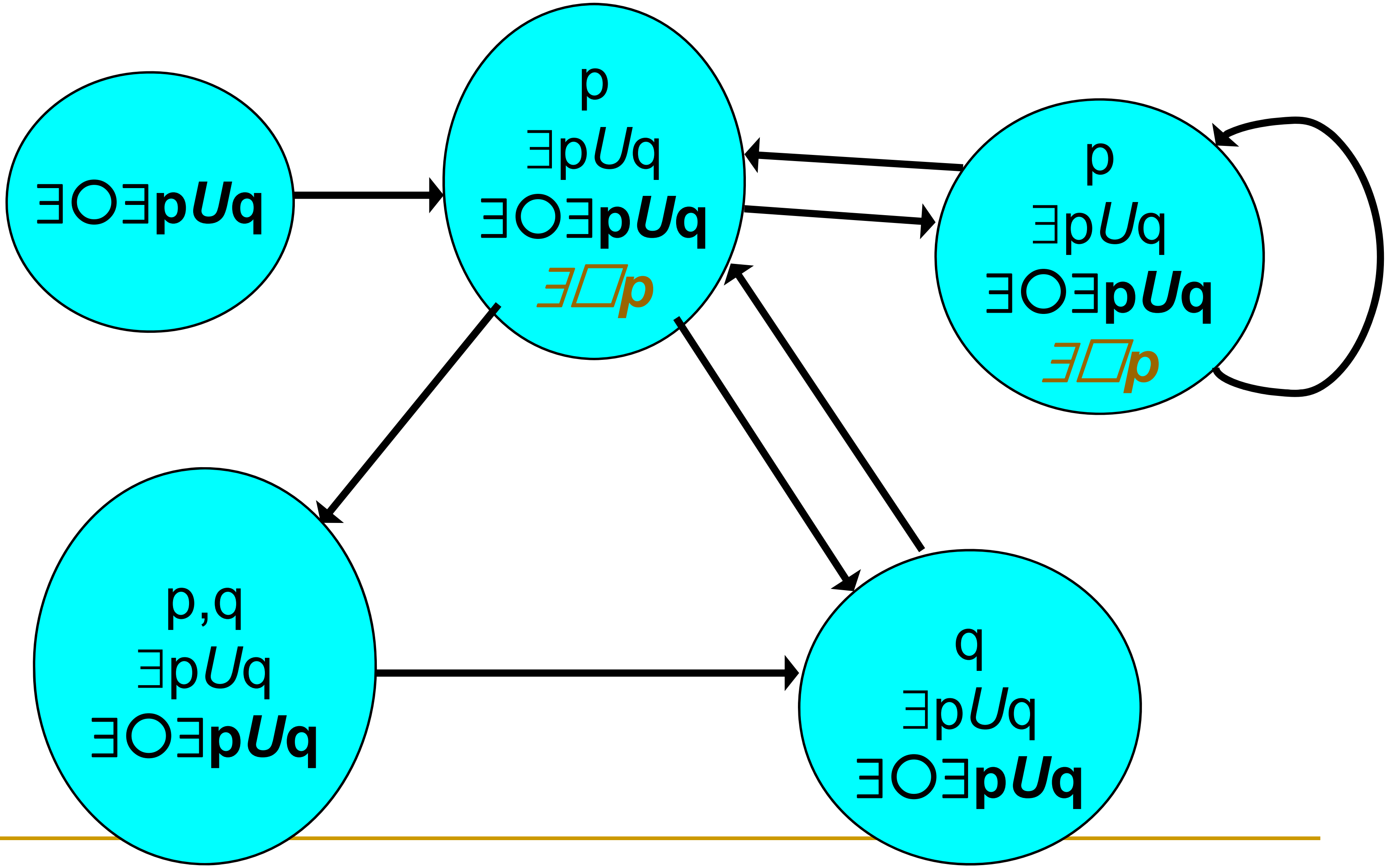
Evaluating ∃□p using greatest fixpoint



$(\exists\bigcirc\exists p\,\mathit{U}q) \wedge \exists\Box p$

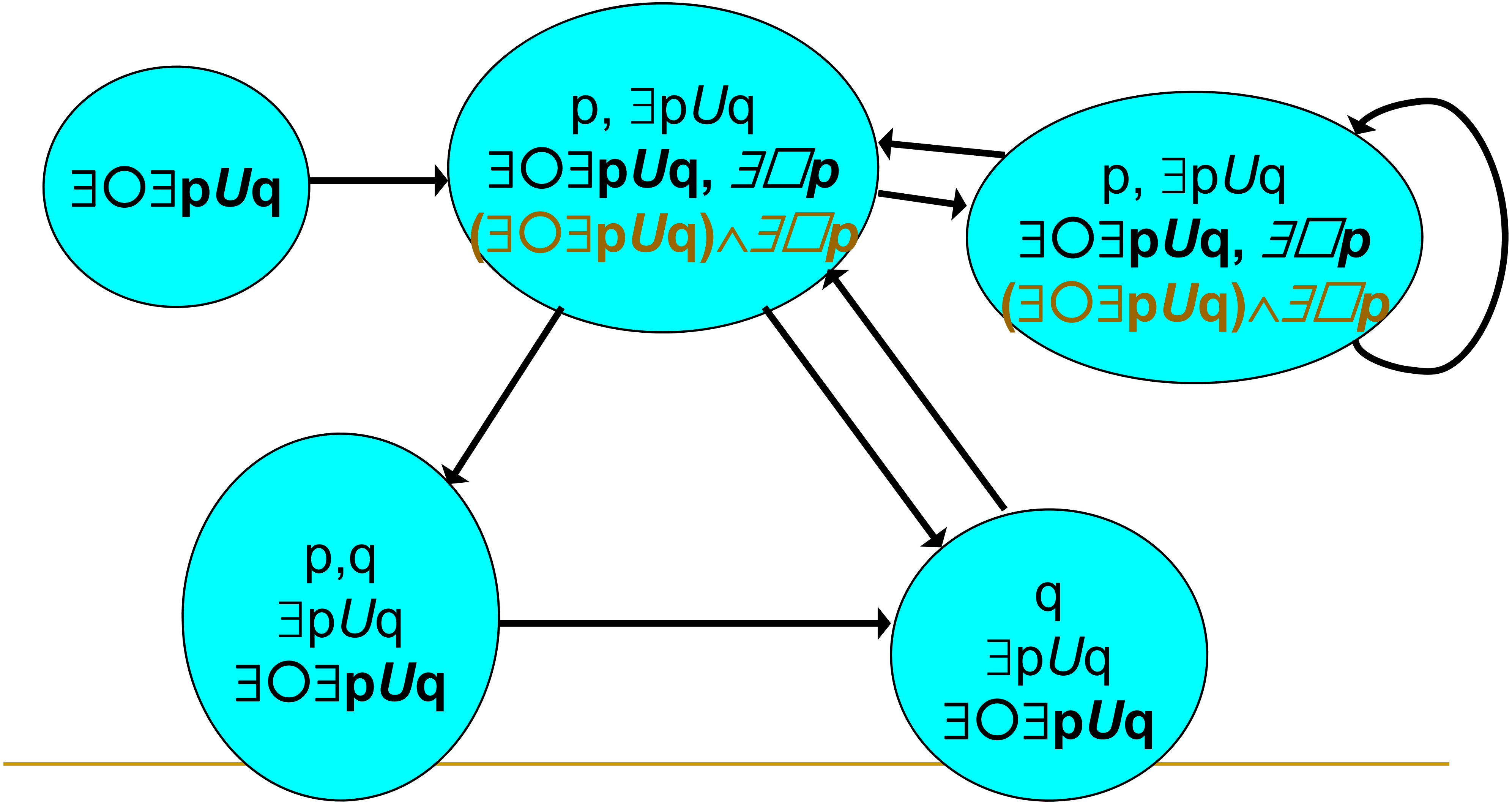
Evaluating $\exists\Box p$ using greatest fixpoint

Result:

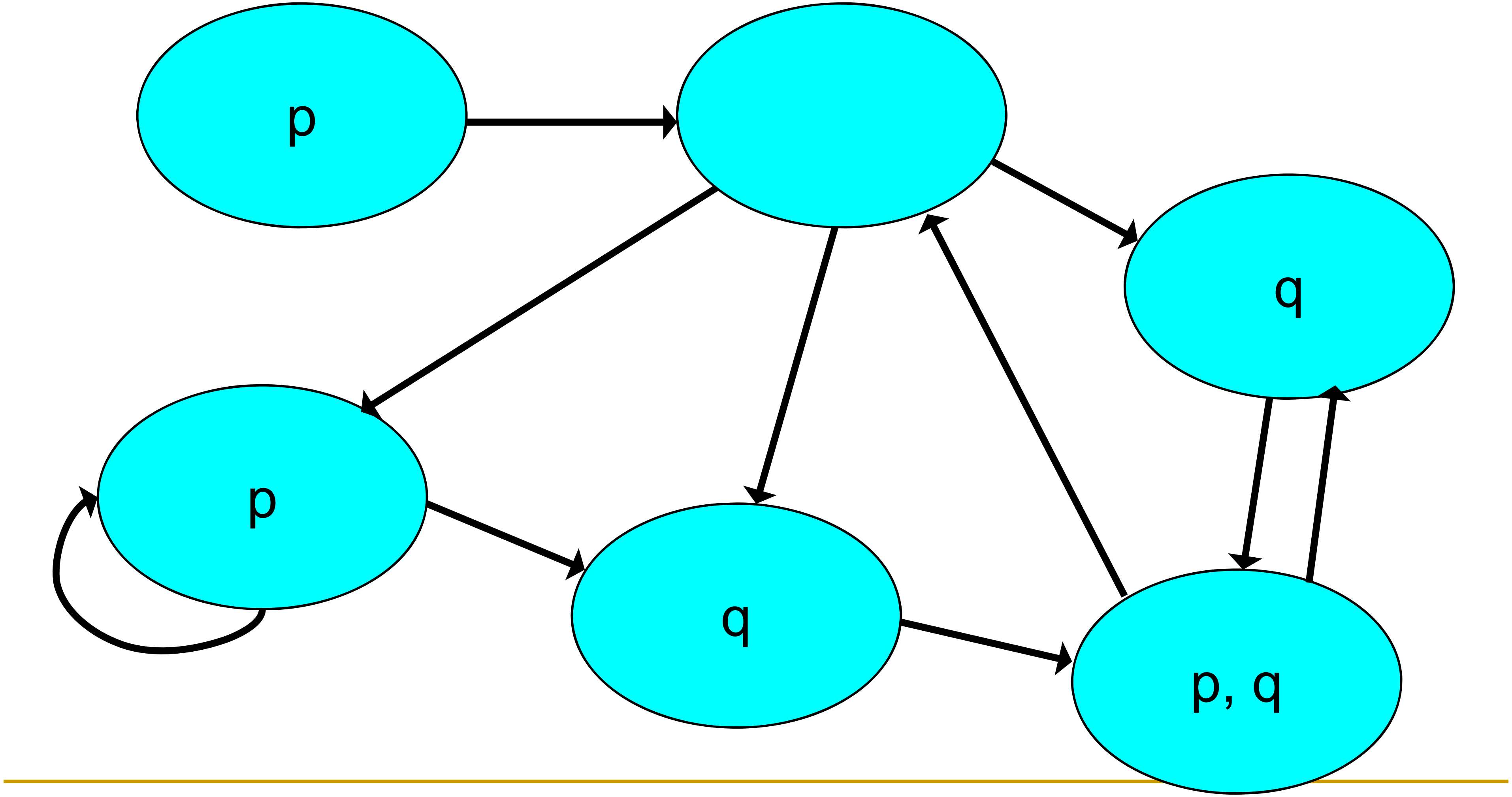


(∃○∃p*U*q) ∧ ∃□p

Finally, evaluating (∃○∃pUq) ∧ ∃□p



Workout: labelling $\exists\Diamond(p\wedge \exists\Box q)$



CTL
- model-checking problem complexities

- The PLTL model-checking problem is PSPACE-complete.
  - definition: Is there a run that satisfies the formula ?
- The PLTL without $\bigcirc$ （**modal operator “next”**） model-checking problem is NP-complete.
- The model-checking problem of CTL is PTIME-complete.
- The model-checking problem of CTL* is PSPACE-complete.

CTL
- symbolic model-checking with BDD

- System states are described with binary variables.

  ***n* binary variables** → **2^n states**

  $x_1, x_2, \ldots\ldots, x_n$

- we can use a BDD to describe legal states.

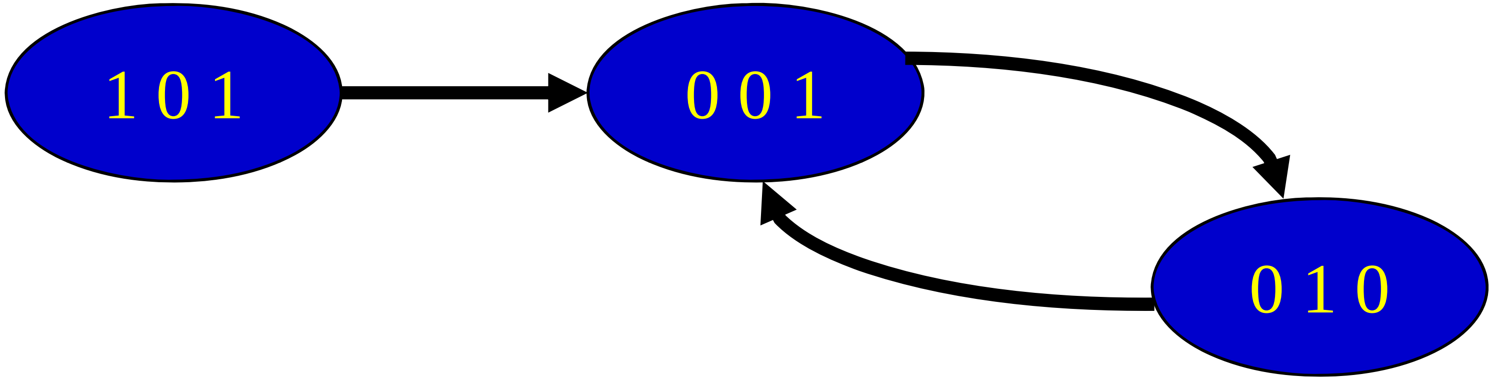
  a Boolean function with *n* binary variables

  state($x_1, x_2, \ldots\ldots, x_n$)

CTL - symbolic model-checking with Propositioal logics

Example:

x_1 x_2 x_3

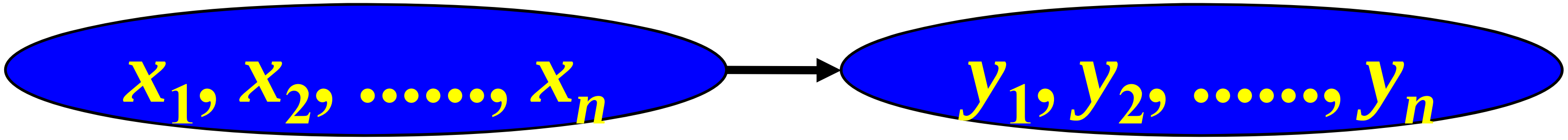


$$
\begin{aligned}
\text{state}(x_1, x_2, x_3) = \quad & (x_1 \land \neg x_2 \land x_3) \\
\lor \quad & (\neg x_1 \land \neg x_2 \land x_3) \\
\lor \quad & (\neg x_1 \land x_2 \land \neg\, x_3)
\end{aligned}
$$

CTL - symbolic model-checking with Propositioal logics

State transition relation as a logic funciton with 2*n* parameters

transition(x_1, x_2, ……, x_n, y_1, y_2, ……, y_n)

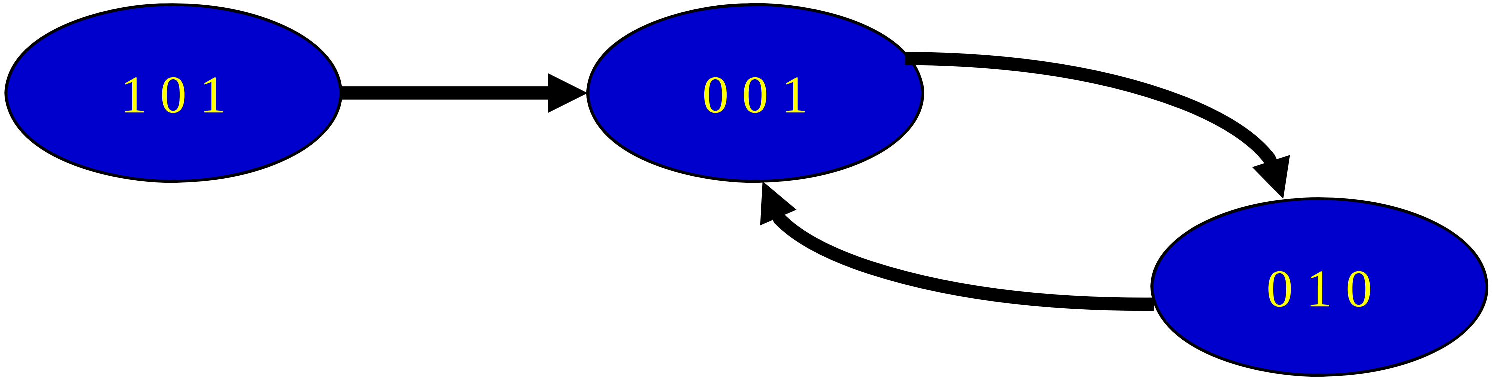


$x_1, x_2, \ldots\ldots, x_n$

$y_1, y_2, \ldots\ldots, y_n$

CTL - symbolic model-checking with Propositioal logics

$x_1 \quad x_2 \quad x_3 \quad y_1 \quad y_2 \quad y_3$



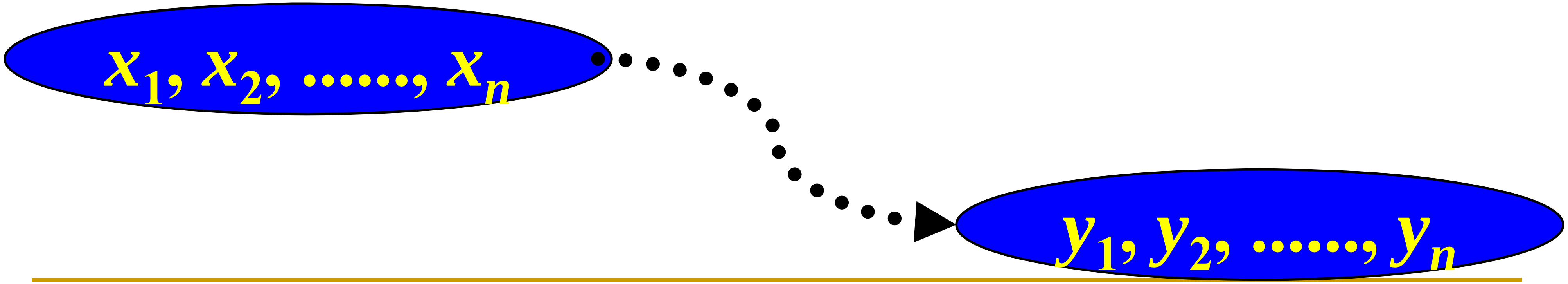
transition$(x_1, x_2, x_3, y_1, y_2, y_3) =$

$$(x_1 \wedge \neg x_2 \wedge x_3 \wedge \neg y_1 \wedge \neg y_2 \wedge y_3)$$
$$\vee \quad (\neg x_1 \wedge \neg x_2 \wedge x_3 \wedge \neg y_1 \wedge y_2 \wedge \neg y_3)$$
$$\vee \quad (\neg x_1 \wedge x_2 \wedge \neg x_3 \wedge \neg y_1 \wedge \neg y_2 \wedge y_3)$$

CTL - symbolic model-checking with Propositioal logics

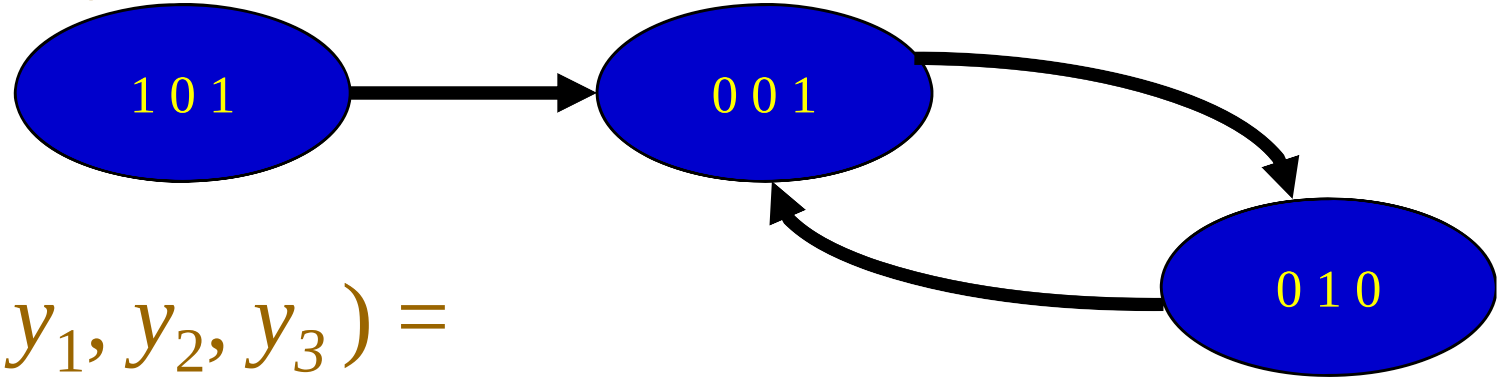
Path relation also as a logic funciton with <u>2*n* parameters</u>

reach(x_1, x_2, ……, x_n, y_1, y_2, ……, y_n)



CTL - symbolic model-checking with Propositioal logics

$x_1 \quad x_2 \quad x_3 \quad y_1 \quad y_2 \quad y_3$



$$
\begin{aligned}
\text{reach}(x_1, x_2, x_3, y_1, y_2, y_3) = \; & (x_1 \wedge \neg x_2 \wedge x_3 \wedge \neg y_1 \wedge \neg y_2 \wedge y_3) \\
\vee \; & (x_1 \wedge \neg x_2 \wedge x_3 \wedge \neg y_1 \wedge y_2 \wedge \neg y_3) \\
\vee \; & (\neg x_1 \wedge \neg x_2 \wedge x_3 \wedge \neg y_1 \wedge y_2 \wedge \neg y_3) \\
\vee \; & (\neg x_1 \wedge x_2 \wedge \neg x_3 \wedge \neg y_1 \wedge \neg y_2 \wedge y_3) \\
\vee \; & (\neg x_1 \wedge \neg x_2 \wedge x_3 \wedge \neg y_1 \wedge \neg y_2 \wedge y_3) \\
\vee \; & (\neg x_1 \wedge x_2 \wedge \neg x_3 \wedge \neg y_1 \wedge y_2 \wedge \neg y_3)
\end{aligned}
$$

Symbolic safety analysis

- I : initial condition with parameters

  $x, x_2, \ldots\ldots, x_n$

- η : safe condition with parameters

  $y_1, y_2, \ldots\ldots, y_n$

- If $I \wedge \neg\eta \wedge \text{reach}(x_1, x_2, \ldots\ldots, x_n, y_1, y_2, \ldots\ldots, y_n)$ is not false,
  - a risk state is reachable.
  - *the system is not safe.*

Symbolic safety analysis (backward)

Encode the states with variables $x_0,x_1,\ldots,x_n$.

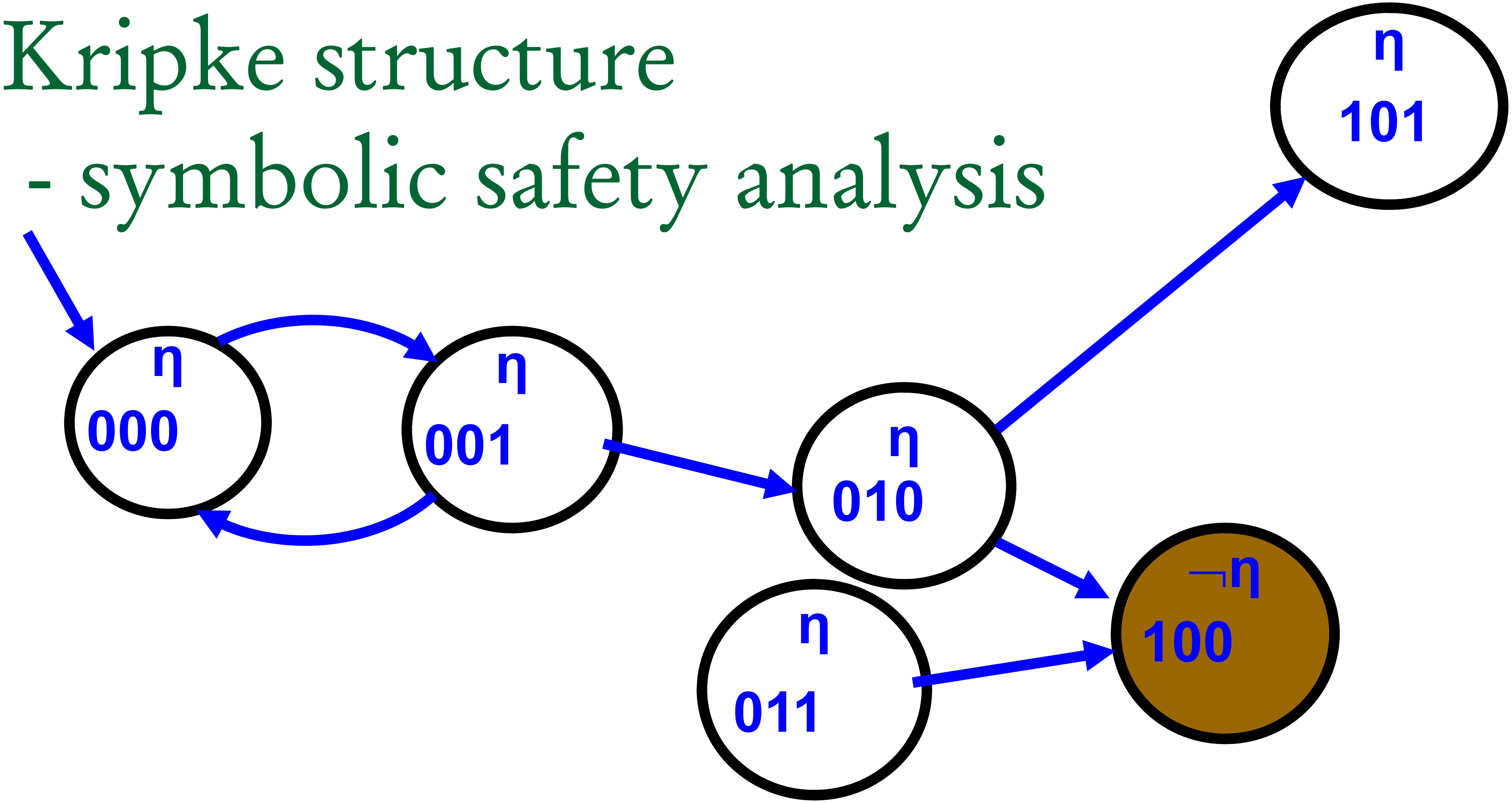
- the state set as a proposition formula: $s(x_0,x_1,\ldots,x_n)$
- the risk state set as r $(x_0,x_1,\ldots,x_n)$
- the initial state set as $i(x_0,x_1,\ldots,x_n)$
- the transition set as $t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n)$

change all umprimed variable in b_{k-1} to primed.

```
$b_0 = r(x_0,x_1,\ldots,x_n) \wedge s(x_0,x_1,\ldots,x_n)$; $k = 1$;
repeat
    $b_k = b_{k-1} \vee \exists x'_0 \exists x'_1 \ldots \exists x'_n (t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n) \wedge (b_{k-1}\uparrow))$;
    $k = k+1$;
  until $b_k \equiv b_{k-1}$;
if $(b_k \wedge i(x_0,x_1,\ldots,x_n)) \equiv$ false, return 'safe'; else return 'risky';
```

a least fixpoint procedure

Kripke structure
- symbolic safety analysis



states: $s(x,y,z) \equiv (\neg x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge \neg y \wedge z) \vee (\neg x \wedge y \wedge \neg z)$

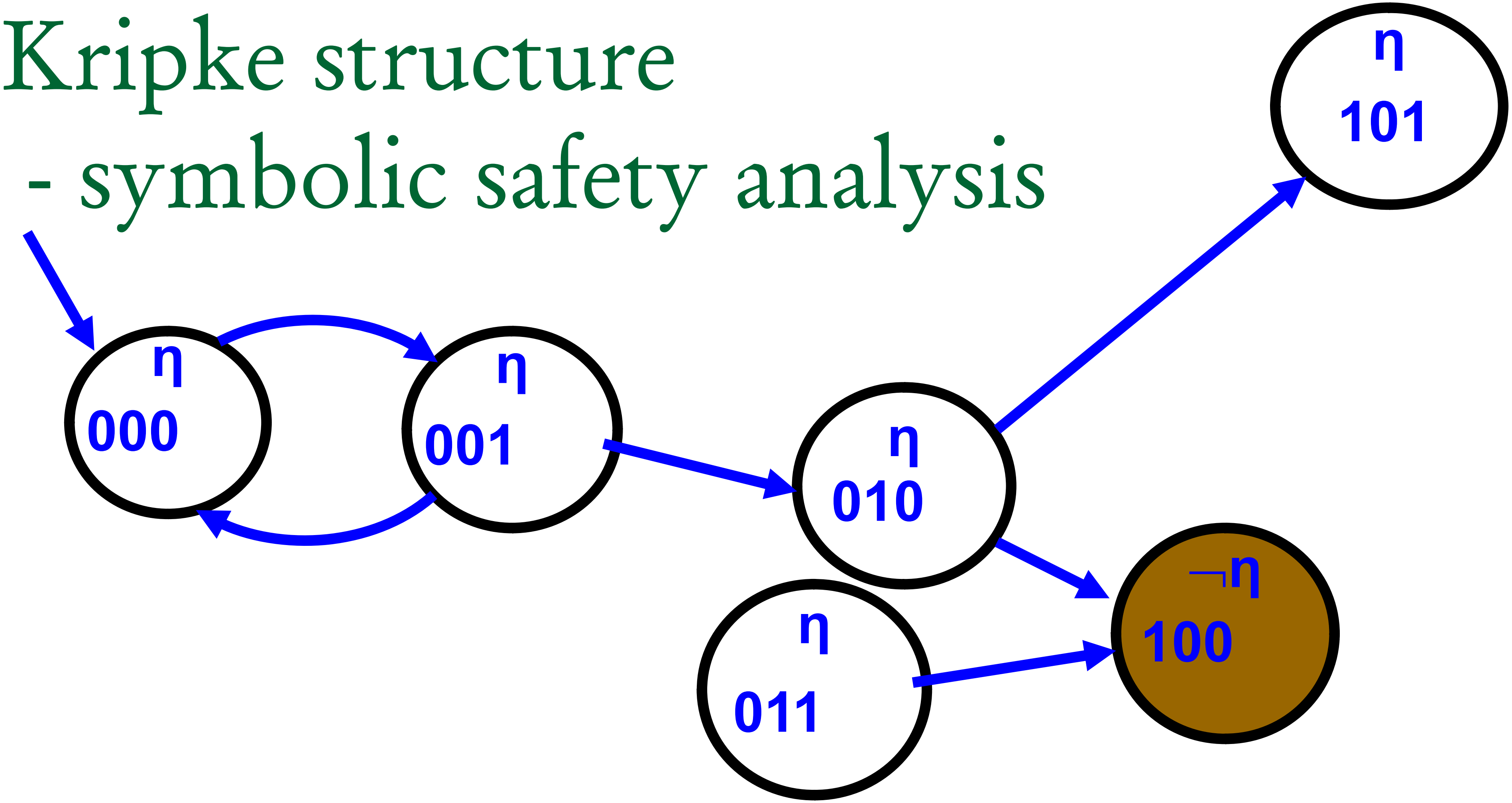
$\vee (\neg x \wedge y \wedge z) \vee (x \wedge \neg y \wedge \neg z) \vee (x \wedge \neg y \wedge z)$

$\equiv (\neg x) \vee (x \wedge \neg y)$

initial state: $i(x,y,z) \equiv \neg x \wedge \neg y \wedge \neg z$

risk state: $r(x,y,z) \equiv x \wedge \neg y \wedge \neg z$

Kripke structure - symbolic safety analysis



transitions: $T(x,y,z,x',y',z') \equiv$

$(\neg x \wedge \neg y \wedge \neg z \wedge \neg x' \wedge \neg y' \wedge z') \vee (\neg x \wedge \neg y \wedge z \wedge \neg x' \wedge \neg y' \wedge \neg z')$
$\vee (\neg x \wedge \neg y \wedge z \wedge \neg x' \wedge y' \wedge \neg z') \vee (\neg x \wedge y \wedge \neg z \wedge x' \wedge \neg y' \wedge \neg z')$
$\vee (\neg x \wedge y \wedge \neg z \wedge x' \wedge \neg y' \wedge z') \vee (\neg x \wedge y \wedge z \wedge x' \wedge \neg y' \wedge \neg z')$

Symbolic safety analysis (backward)

$b_0 = r(x,y,z) \equiv x \wedge \neg y \wedge \neg z$

$b_1 = b_0 \vee \exists x' \exists y' \exists z' (t(x,y,z,x',y',z') \wedge b_0 \uparrow)$

$= (x \wedge \neg y \wedge \neg z) \vee \exists x' \exists y' \exists z' (t(x,y,z,x',y',z') \wedge x' \wedge \neg y' \wedge \neg z')$

$= (x \wedge \neg y \wedge \neg z) \vee \exists x' \exists y' \exists z' (((\neg x \wedge y \wedge \neg z) \vee (\neg x \wedge y \wedge z)) \wedge x' \wedge \neg y' \wedge \neg z')$

$= (x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge y \wedge \neg z) \vee (\neg x \wedge y \wedge z)$

$b_2 = b_1 \vee \exists x' \exists y' \exists z' (t(x,y,z,x',y',z') \wedge b_1 \uparrow)$

$= (\neg x \wedge \neg y \wedge z) \vee (x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge y \wedge \neg z) \vee (\neg x \wedge y \wedge z)$

$b_3 = b_2 \vee \exists x' \exists y' \exists z' (t(x,y,z,x',y',z') \wedge b_2 \uparrow)$

$= (\neg x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge \neg y \wedge z) \vee (x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge y \wedge \neg z) \vee (\neg x \wedge y \wedge z)$

fixpoint

$b_4 = b_3 \vee \exists x' \exists y' \exists z' (t(x,y,z,x',y',z') \wedge b_3 \uparrow)$

$= (\neg x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge \neg y \wedge z) \vee (x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge y \wedge \neg z) \vee (\neg x \wedge y \wedge z)$

$b_4 \wedge i(x,y,z) = (\neg x \wedge \neg y \wedge \neg z)$

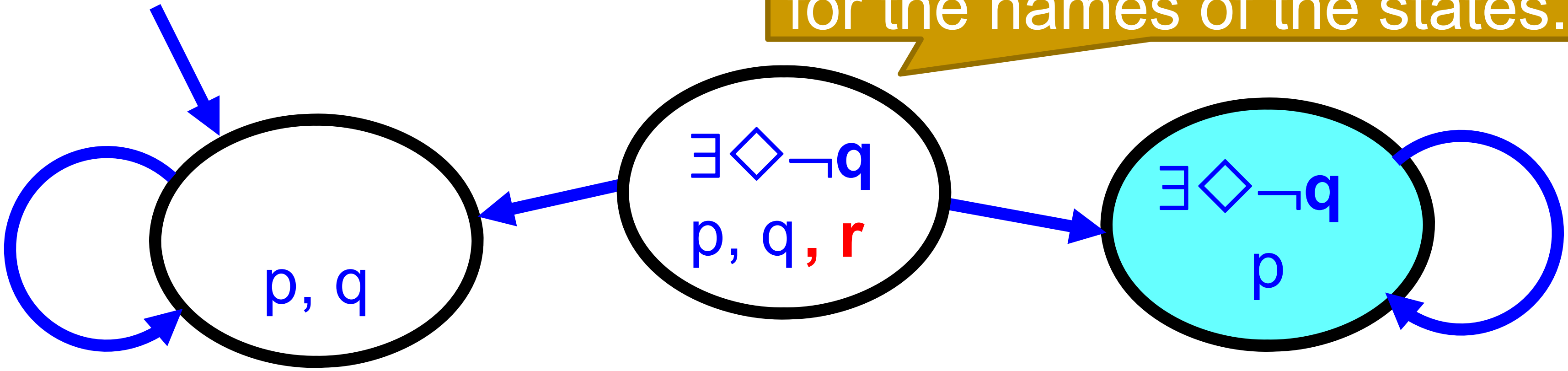
non-empty intersection with the initial condition → risk detected.

Symbolic safety analysis (backward)

One assumption for the correctness!

- Two states cannot be with the same proposition labeling.
- Otherwise, the collapsing of the states may cause problem.

may need a few propositions for the names of the states.



Symbolic safety analysis (forward)

Encode the states with variables $x_0,x_1,\ldots,x_n$.

- the state set as a proposition formula: $s(x_0,x_1,\ldots,x_n)$
- the risk state set as $r\ (x_0,x_1,\ldots,x_n)$
- the initial state set as $i(x_0,x_1,\ldots,x_n)$
- the transition set as $t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n)$

```
$f_0 = i(x_0,x_1,\ldots,x_n) \wedge s(x_0,x_1,\ldots,x_n)$; $k = 1$;
repeat
    $f_k = f_{k-1} \vee (\exists x_0 \exists x_1 \ldots \exists x_n (t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n) \wedge f_{k-1}))\downarrow$;
    $k = k + 1$;
  until $f_k \equiv f_{k-1}$;
if $(f_k \wedge r(x_0,x_1,\ldots,x_n)) \equiv$ false, return 'safe'; else return 'risky';
```

change all primed variable to umprimed.

Symbolic safety analysis (forward)

$$f_0 = i(x,y,z) \equiv \neg x \wedge \neg y \wedge \neg z$$

$$\begin{aligned}
f_1 &= f_0 \vee (\exists x \exists y \exists z(t(x,y,z,x',y',z') \wedge f_0))\downarrow \\
&= (\neg x \wedge \neg y \wedge \neg z) \vee (\exists x \exists y \exists z(t(x,y,z,x',y',z') \wedge \neg x \wedge \neg y \wedge \neg z))\downarrow \\
&= (\neg x \wedge \neg y \wedge \neg z) \vee (\exists x \exists y \exists z(\neg x' \wedge \neg y' \wedge z' \wedge \neg x \wedge \neg y \wedge \neg z))\downarrow \\
&= (\neg x \wedge \neg y \wedge \neg z) \vee (\neg x' \wedge \neg y' \wedge z')\downarrow \\
&= (\neg x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge \neg y \wedge z) = \neg x \wedge \neg y
\end{aligned}$$

$$f_2 = f_1 \vee (\exists x \exists y \exists z(t(x,y,z,x',y',z') \wedge f_1)\downarrow = (\neg x \wedge \neg y) \vee (\neg x \wedge y \wedge \neg z)$$

$$f_3 = f_2 \vee (\exists x \exists y \exists z(t(x,y,z,x',y',z') \wedge f_2)\downarrow = (\neg y) \vee (\neg x \wedge y \wedge \neg z)$$

$$f_4 = f_3 \vee (\exists x \exists y \exists z(t(x,y,z,x',y',z') \wedge f_3)\downarrow = (\neg y) \vee (\neg x \wedge y \wedge \neg z)$$

$$f_4 \wedge r(x,y,z) = ((\neg y) \vee (\neg x \wedge y \wedge \neg z)) \wedge (x \wedge \neg y \wedge \neg z) = (x \wedge \neg y \wedge \neg z)$$

fixpoint

non-empty intersection with the risk condition → risk detected.

Bounded model-checking

The value of x_n at state k.

Encode the states with variables $x_{0,k},x_{1,k},\ldots,x_{n,k}$.

- the state set as a proposition formula: $s(x_{0,k},x_{1,k},\ldots,x_{n,k})$
- the risk state set as $r(x_{0,k},x_{1,k},\ldots,x_{n,k})$
- the initial state set as $i(x_{0,0},x_{1,0},\ldots,x_{n,0})$
- the transition set as $t(x_{0,k-1},x_{1,k-1},\ldots,x_{n,k-1},x_{0,k},x_{1,k},\ldots,x_{n,k})$

```
$f_0 = i(x_{0,0},x_{1,0},\ldots,x_{n,0}) \wedge s(x_{0,0},x_{1,0},\ldots,x_{n,0})$; $k = 1$;
repeat
    $f_k = t(x_{0,k-1},x_{1,k-1},\ldots,x_{n,k-1},x_{0,k},x_{1,k},\ldots,x_{n,k}) \wedge f_{k-1}$;
    $k = k + 1$;
  until $f_k \wedge r(x_{0,k},x_{1,k},\ldots,x_{n,k}) \neq$ false
```

When to stop ?

1. diameter of the state graph
2. explosion up to tens of steps.

Bounded model-checking

$f_0 = i(x,y,z) \equiv \neg x_0 \wedge \neg y_0 \wedge \neg z_0$

$f_1 = t(x_0,y_0,z_0,x_1,y_1,z_1) \wedge f_0 = \neg x_0 \wedge \neg y_0 \wedge \neg z_0 \wedge \neg x_1 \wedge \neg y_1 \wedge z_1$

$f_2 = t(x_1,y_1,z_1,x_2,y_2,z_2) \wedge f_1$

$= \neg x_0 \wedge \neg y_0 \wedge \neg z_0 \wedge \neg x_1 \wedge \neg y_1 \wedge z_1 \wedge ((\neg x_2 \wedge \neg y_2 \wedge \neg z_2) \vee (\neg x_2 \wedge y_2 \wedge \neg z_2))$

$f_3 = t(x_2,y_2,z_2,x_3,y_3,z_3) \wedge f_2$

$= \neg x_0 \wedge \neg y_0 \wedge \neg z_0 \wedge \neg x_1 \wedge \neg y_1 \wedge z_1$

$\wedge(\ (\neg x_2 \wedge \neg y_2 \wedge \neg z_2 \wedge \neg x_3 \wedge \neg y_3 \wedge z_3)$

$\vee(\neg x_2 \wedge y_2 \wedge \neg z_2 \wedge ((x_3 \wedge \neg y_3 \wedge \neg z_3) \vee (x_3 \wedge \neg y_3 \wedge z_3)))$

$)$

$= \neg x_0 \wedge \neg y_0 \wedge \neg z_0 \wedge \neg x_1 \wedge \neg y_1 \wedge z_1$

$\wedge((\neg x_2 \wedge \neg y_2 \wedge \neg z_2 \wedge \neg x_3 \wedge \neg y_3 \wedge z_3) \vee (\neg x_2 \wedge y_2 \wedge \neg z_2 \wedge x_3 \wedge \neg y_3))$

$f_3 \wedge r(x_3,y_3,z_3) = (x_3 \wedge \neg y_3 \wedge \neg z_3)$

Symbolic liveness analysis

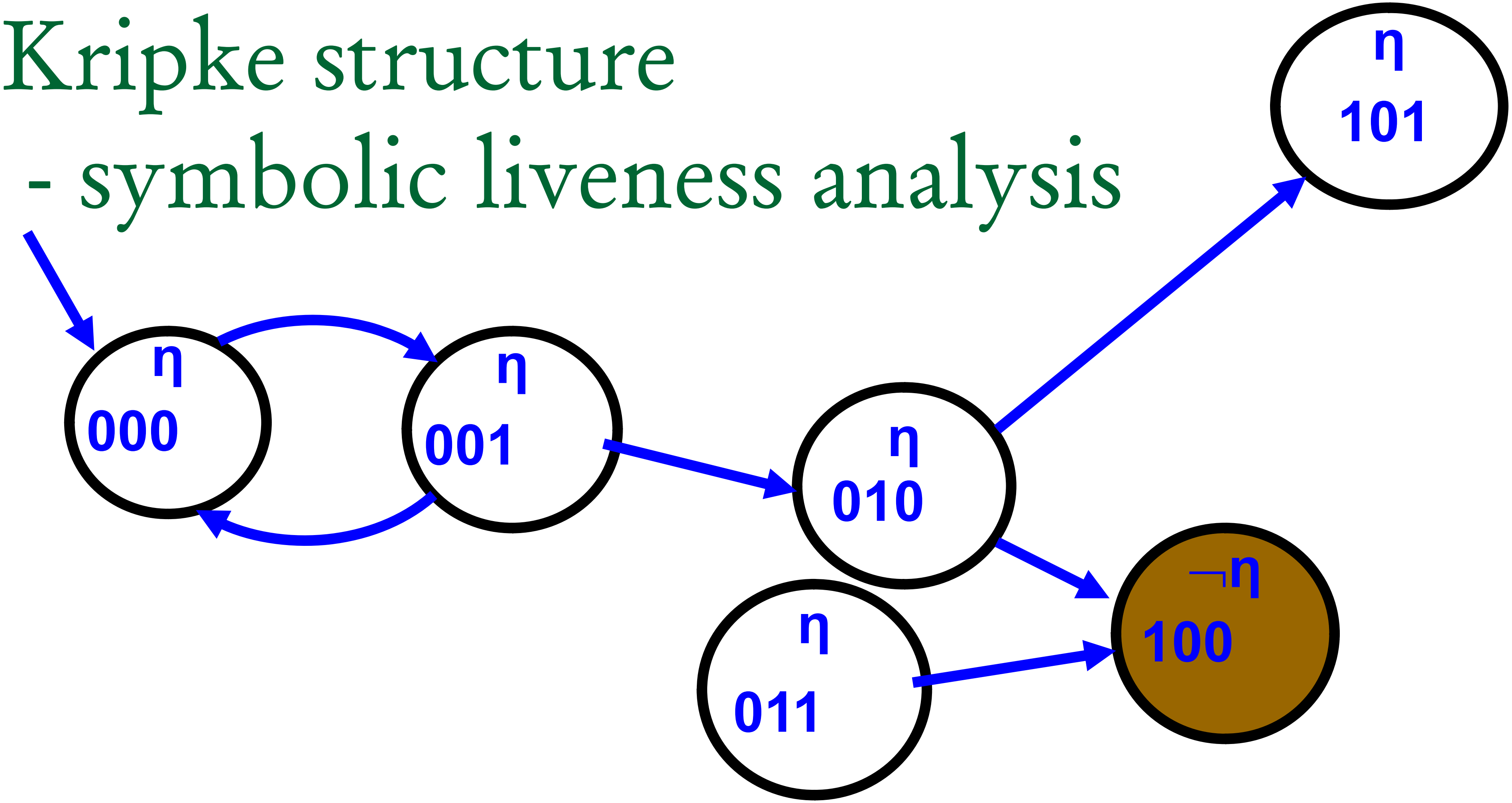
Encode the states with variables x0,x1,…,xn.

- the state set as a proposition formula: $s(x_0,x_1,\ldots,x_n)$
- the non-liveness state set as $b(x_0,x_1,\ldots,x_n)$
- the initial state set as $i(x_0,x_1,\ldots,x_n)$
- the transition set as $t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n)$

```
$b_0 = b(x_0,x_1,\ldots,x_n) \wedge s(x_0,x_1,\ldots,x_n)$; $k = 1$;
repeat
    $b_k = b_{k-1} \wedge \exists x'_0 \exists x'_1 \ldots \exists x'_n (t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n) \wedge b_{k-1}\uparrow)$;
    $k = k + 1$;
  until $b_k \equiv b_{k-1}$;
if $(b_k \wedge i(x_0,x_1,\ldots,x_n)) \equiv$ false, return 'live'; else return 'not live';
```

change all umprimed variable in b_{k-1} to primed.

Kripke structure
- symbolic liveness analysis



states: $s(x,y,z) \equiv (\neg x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge \neg y \wedge z) \vee (\neg x \wedge y \wedge \neg z)$

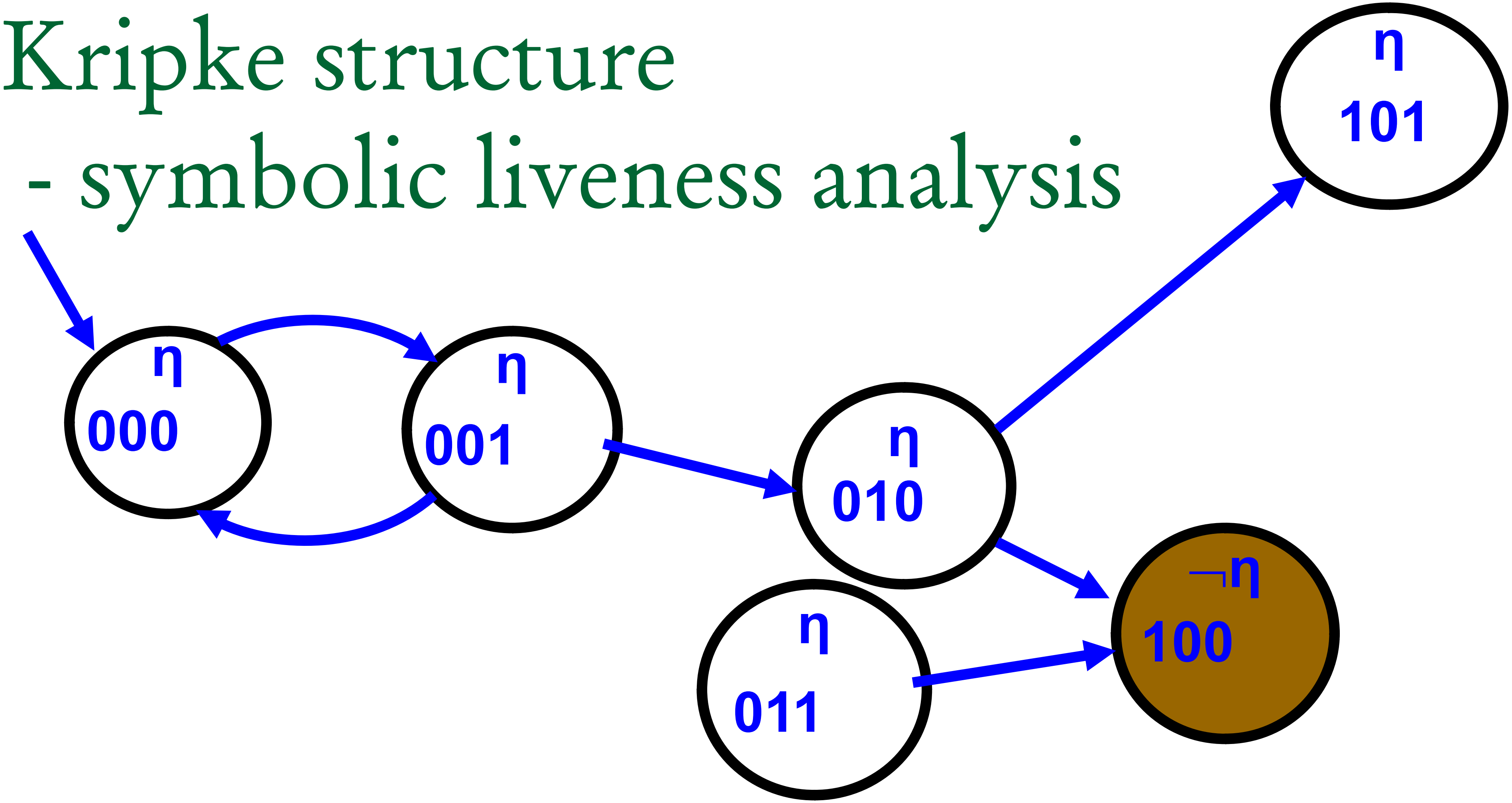
$\vee (\neg x \wedge y \wedge z) \vee (x \wedge \neg y \wedge \neg z) \vee (x \wedge \neg y \wedge z)$

$\equiv (\neg x) \vee (x \wedge \neg y)$

initial state: $i(x,y,z) \equiv \neg x \wedge \neg y \wedge \neg z$

non-liveness state: $b(x,y,z) \equiv (\neg x) \vee (x \wedge \neg y \wedge z)$

Kripke structure
- symbolic liveness analysis



transitions: $T(x,y,z,x',y',z') \equiv$

$(\neg x \wedge \neg y \wedge \neg z \wedge \neg x' \wedge \neg y' \wedge z') \vee (\neg x \wedge \neg y \wedge z \wedge \neg x' \wedge \neg y' \wedge \neg z')$

$\vee (\neg x \wedge \neg y \wedge z \wedge \neg x' \wedge y' \wedge \neg z') \vee (\neg x \wedge y \wedge \neg z \wedge x' \wedge \neg y' \wedge \neg z')$

$\vee (\neg x \wedge y \wedge \neg z \wedge x' \wedge \neg y' \wedge z') \vee (\neg x \wedge y \wedge z \wedge x' \wedge \neg y' \wedge \neg z')$

Symbolic liveness analysis

$b0 = b(x,y,z) \equiv (\neg x)\vee(x\wedge\neg y \wedge z)$

$b1 = b0 \wedge \exists x'\exists y'\exists z'(T(x,y,z,x',y',z') \wedge b0')$

$\quad = ((\neg x)\vee(x\wedge\neg y \wedge z))$

$\quad\quad \wedge\ \exists x'\exists y'\exists z'(T(x,y,z,x',y',z') \wedge ((\neg x')\vee(x'\wedge\neg y'\wedge z')))$

$\quad = ((\neg x)\vee(x\wedge\neg y \wedge z))\wedge$

$\quad \exists x'\exists y'\exists z'(\ \ ((\neg x\wedge\neg y\wedge\neg z)\vee(\neg x\wedge y\wedge\neg z)\vee(\neg x\wedge\neg y\wedge z))$

$\quad\quad\quad\quad \wedge((\neg x')\vee(x'\wedge\neg y'\wedge z')))$

$\quad = (\neg x\wedge\neg y\wedge\neg z)\vee(\neg x\wedge y\wedge\neg z)\vee(\neg x\wedge\neg y\wedge z)$

$b2 = b1 \wedge \exists x'\exists y'\exists z'(T(x,y,z,x',y',z') \wedge b1')$

$\quad = (\neg x\wedge\neg y\wedge\neg z)\vee(\neg x\wedge y\wedge\neg z)$

$b3 = b2 \wedge \exists x'\exists y'\exists z'(T(x,y,z,x',y',z') \wedge b2')$

$\quad = (\neg x\wedge\neg y\wedge\neg z)\vee(\neg x\wedge y\wedge\neg z)$

fixpoint

non-empty intersection with the initial condition → non-liveness detected.

CTL
- symbolic model-checking algorithm

Assume program with rules $r_1, r_2, \ldots, r_n$

```
label(φ ) {
case p, return p;
case $\neg\varphi$, return $\neg$label($\varphi$);
case $\varphi\vee\psi$, return label($\varphi$) $\vee$ label($\psi$),
case $\exists\bigcirc\varphi$, return $\vee_{i=n}$ *pred*($r_i$, label($\varphi$));
case $\exists\ \psi_1\ \mathbf{U}\psi_2$ , return lfp(l abel($\psi_1$), label($\psi_2$) );
case $\exists\Box\varphi$, return gfp(label($\varphi$));
}
```

Symbolic model-checking
- with real-world programs

Consider guarded commands with modes (GCM)

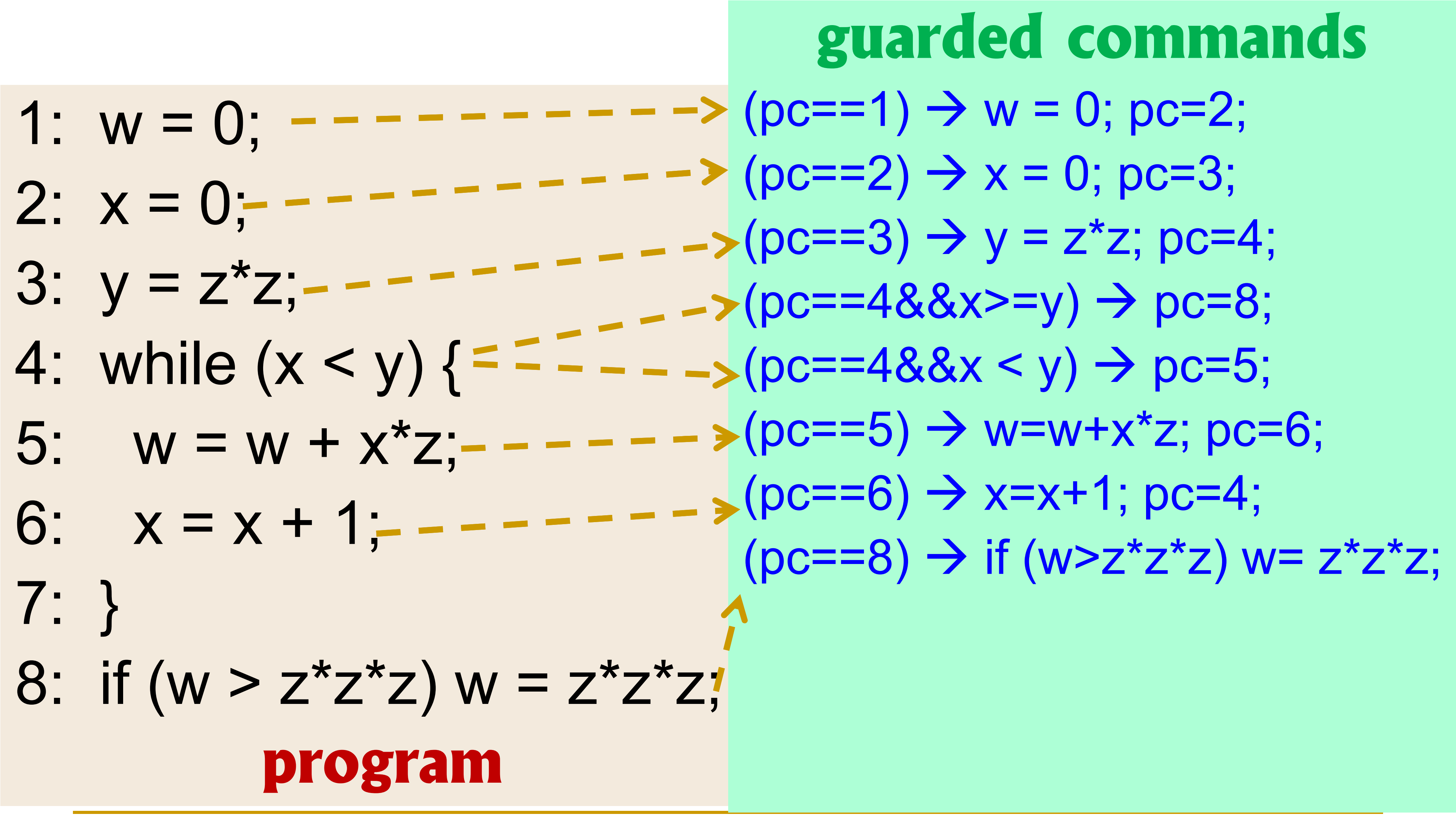
Guard → Actions

- Guard is a propositional formula of state variables.
- Actions is a command of the following syntax.

C ::= ACT | *{*C*}* | C C | *if (*B*)* C *else* C | *while (*B*)* C

ACT ::= *;* | *goto name;* | *x* = E *;*

Guarded commands with modes (GCM)



```
1:  w = 0;
2:  x = 0;
3:  y = z*z;
4:  while (x < y) {
5:     w = w + x*z;
6:     x = x + 1;
7:  }
8:  if (w > z*z*z) w = z*z*z;
```

program

guarded commands

```
(pc==1) → w = 0; pc=2;
(pc==2) → x = 0; pc=3;
(pc==3) → y = z*z; pc=4;
(pc==4&&x>=y) → pc=8;
(pc==4&&x < y) → pc=5;
(pc==5) → w=w+x*z; pc=6;
(pc==6) → x=x+1; pc=4;
(pc==8) → if (w>z*z*z) w= z*z*z;
```

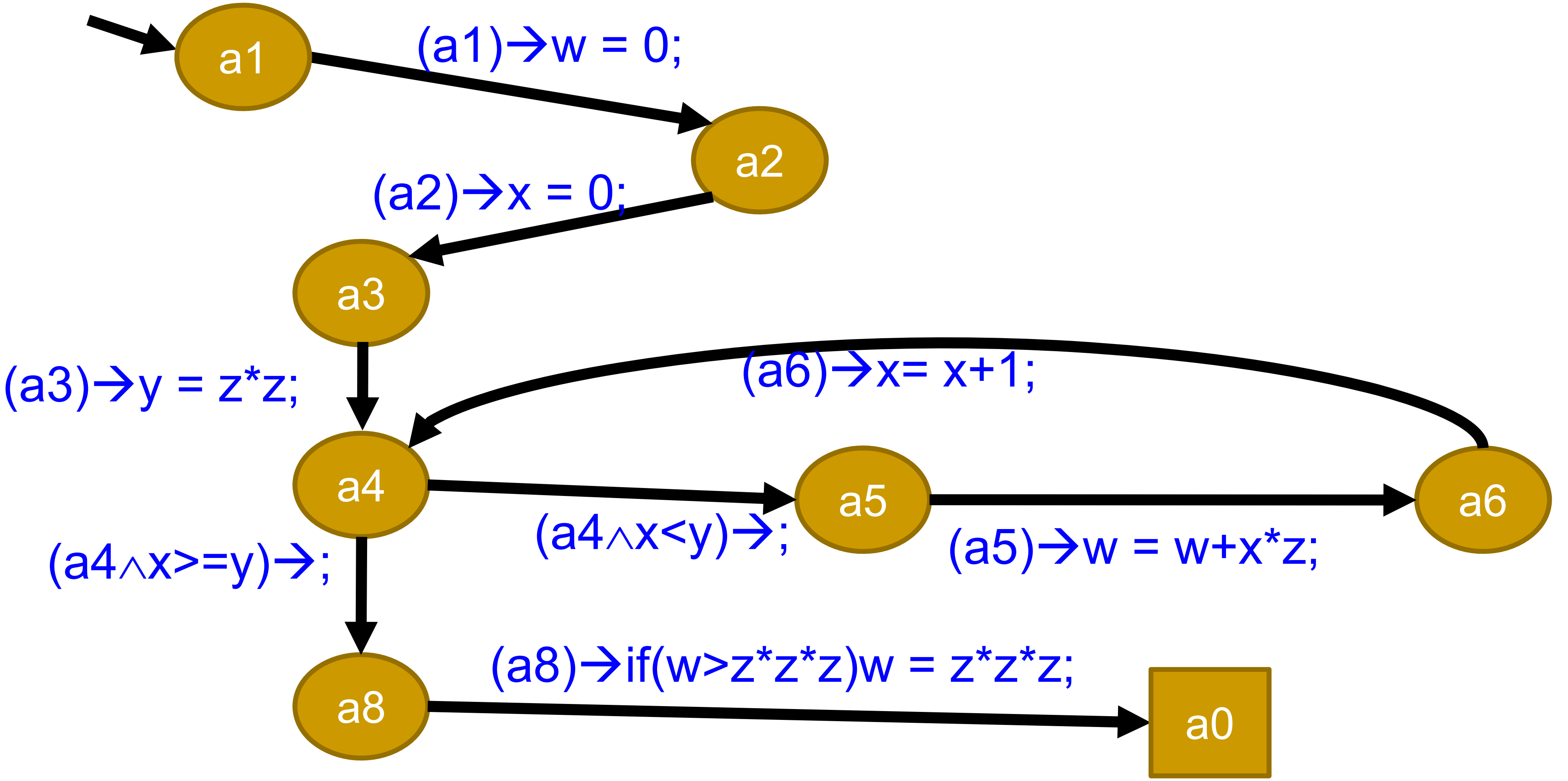
A state-transition
- represented as a GCM

8 rules in total:

```
(a1) → w = 0; goto a2;
(a2) → x = 0; goto a3;
(a3) → y = z*z; goto a4;
(a4&&x>=y) → goto a8;
(a4&&x < y) → goto a5;
(a5) → w=w+x*z; goto a6;
(a6) → x=x+1; goto a4;
(a8) → if (w>z*z*z) w= z*z*z; }
```

A state-transition

- represented as a GCM



Transition relation
- from state-transition graphs

Given a set of rules $r_1, r_2, \ldots, r_m$ of the form

$$r_k: \ (\tau_k) \rightarrow y_{k,0}=d_0;\ y_{k,1}=d_1;\ \ldots;\ y_{k,nk}=d_{nk};$$

$$t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n)$$

$$\equiv \bigvee_{k\in[1,m]} \Big(\ \tau_k \wedge y'_{k,0}==d_0 \wedge y'_{k,1}==d_1 \wedge \ldots \wedge y'_{k,nk}==d_{nk}$$

$$\wedge \bigwedge_{h\in[1,n]} \big(x_h \notin \{y_{k,0},\ y_{k,1},\ \ldots,\ y_{k,nk}\} \Rightarrow x_h == x'_h\big)$$

$$\Big)$$

Transition relation from GCM rules.

Given a set of rules for X={x,y,z}

r_1: (x<y&& y>2) → y=x+y; x=3;

r_2: (z>=2) → y=x+1; z=0;

r_3: (x<2) → x=0;

$t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n)$

$\equiv \quad (x<y \wedge y>2 \wedge y'==x+y \wedge x'==3 \wedge z'==z)$

$\vee(z>=2 \wedge y'==x+1 \wedge z'==0 \wedge x'==x)$

$\vee(x<2 \wedge x'==0 \wedge y'==y \wedge z'==z)$

Transition relation
- from state-transition graphs

In gneral, transition relation is expensive to construct.

Can we do the following state-space construction

$$\exists x'_0 \exists x'_1 \ldots \exists x'_n (t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n) \wedge (b_{k-1}\uparrow))$$

directly with the GCM rules ?

Yes, ***on-the-fly state space construction***.

On-the-fly precondition calculation with GCM rules.

Given a set of rules $r_1, r_2, \ldots, r_m$ of the form

$$r_k: \ (\tau_k) \rightarrow y_{k,0}=d_0;\ y_{k,1}=d_1;\ \ldots;\ y_{k,nk}=d_{nk};$$

$$\exists x'_0 \exists x'_1 \ldots \exists x'_n (t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n) \wedge (b\uparrow))$$

$$\equiv \bigvee_{k\in[1,m]} \Big(\tau_k \wedge \exists y_{k,0} \exists y_{k,1} \ldots \exists y_{k,nk} \Big(b \wedge \bigwedge_{h\in[0,nk]} y_{k,h}{==}d_h\Big)\Big)$$

However, GCM rules are more complex than that.

On-the-fly precondition calculation with GCM rules.

Given a set of rules for X={x,y,z}

r_1: (x<y&& y>2) → y=z; x=3;

r_2: (z>=2) → y=x+1; z=7;

r_3: (x<2) → z=0;

B

$\exists x'_0 \exists x'_1 \ldots \exists x'_n (t(x_0, x_1, \ldots, x_n, x'_0, x'_1, \ldots, x'_n) \wedge (x<4 \wedge z>5) \uparrow)$

$\equiv \quad (x<y \wedge y>2 \wedge \exists y \exists x(x<4 \wedge z>5 \wedge y==z \wedge x==3))$

$\vee(z>=2 \wedge \exists y \exists z(x<4 \wedge z>5 \wedge y==x+1 \wedge z==7))$

$\vee(x<2 \wedge \exists z(x<4 \wedge z>5 \wedge z==0))$

$\equiv \quad (x<y \wedge y>2 \wedge z>5) \vee(z>=2 \wedge x<4) \vee(x<2 \wedge \exists z(\mathrm{false}))$

$\equiv \quad (x<y \wedge y>2 \wedge z>5) \vee(z>=2 \wedge x<4)$

On-the-fly precondition calculation with GCM rules.

Given a set of rules r_1, r_2, …, r_m of the form

$$r_k:\ (\tau_k) \rightarrow s_k;$$

$$\exists x'_0 \exists x'_1 \ldots \exists x'_n (t(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x'_n) \wedge (b\uparrow))$$

$$\equiv \bigvee_{k\in[1,m]} \left(\tau_k \wedge \mathrm{pre}(s_k, b\ )\right)$$

precondition procedure

A general propositional formula

What is pre(s, b) ?

A GCM statement

On-the-fly precondition calculation with GCM rules.

Given a set of rules r_1, r_2, …, r_m of the form

$$r_k\text{:}\ \ (\tau_k) \rightarrow s_k;$$

What is pre(s, b) ?

new expression obtained from b by replacing every occurrence of x with E.

- pre(*x* = *E*;, b) ≡ b[x/E]

Ex 1. the precondition to x=x+z;
(x==y+2 ∧ x<4∧z>5) [x/x+z] ≡ x+z==y+2 ∧ x+z<4∧z>5

Ex 2. the precondition to x=5;
(x==y+2 ∧ x<4∧z>5) [x/x+z] ≡ 5==y+2 ∧ 5<4∧z>5

Ex 3. the precondition to x=2*x+1;
(x==y+2 ∧ x<4∧z>5) [x/x+z] ≡ 2*x+1==y+2 ∧ 2*x+1<4∧z>5

On-the-fly precondition calculation with GCM rules.

Given a set of rules $r_1, r_2, \ldots, r_m$ of the form

$$r_k: \ (\tau_k) \rightarrow s_k;$$

What is pre(s,b) ?

- pre(x = E;, b) ≡ b[x/E]
- pre($s_1 s_2$, b) ≡ pre(s_1, pre(s_2, b))
- pre(if (B) $s_1$else s_2) ≡ (B∧pre(s_1, b))∨(¬B∧pre(s_2,b))
- pre(while (B) s, b) ≡ ….

new expression obtained from b by replacing every occurrence of x with E.

Ex. the precondition to x=x+z;
(x==y+2 ∧ x<4∧z>5) [x/x+z]
≡ x+z==y+2 ∧ x+z<4∧z>5

On-the-fly precondition calculation with GCM rules.

Given a set of rules r_1, r_2, …, r_m of the form

$$r_k: \ (\tau_k) \rightarrow s_k;$$

What is pre(s,b) ?

pre(while (B) s, b) $\equiv$ formula $L_1 \vee L_2$ for

L_1: those states that reach $\neg B \wedge b$ with finite steps of s through states in B; and

L_2: those states that never leave B with steps of s.

On-the-fly precondition calculation with GCM rules.

L_1: those states that reach $\neg B \wedge b$ with finite steps of s through states in B

```
$w_0 = \neg B \wedge b$; $k = 1$;
repeat
    $w_k = w_{k-1} \vee (B \wedge pre(s, w_{k-1}))$;
    $k = k + 1$;
  until $w_k \equiv w_{k-1}$;
return $w_k$;
```

also a least fixpoint procedure

Precondition to b through while (B) s;

```
$w_0 = \neg B \wedge b$; k = 1;
repeat
    $w_k = w_{k-1} \vee (B \wedge pre(s, w_{k-1}))$;
    k = k +1;
  until $w_k \equiv w_{k-1}$;
return $w_k$;
```

Example: $b \equiv x==2 \wedge y == 3$

while (x < y) x = x+1;

L1 computation.

$w_0 \equiv x>=y \wedge x==2 \wedge y==3 \equiv \text{false}$; k = 1;

$w_1 \equiv \text{false} \vee (x<y \wedge pre(x=x+1, \text{false}))$;

$\equiv \text{false} \vee (x<y \wedge \text{false})$;

$\equiv$ false;

On-the-fly precondition calculation with GCM rules.

Given a set of rules r_1, r_2, …, r_m of the form

pre(while (B) s, b)

L_2: those states that never leave B with steps of s.

a *greatest* fixpoint procedure

```
$w_0$ = B; k = 1;
repeat
    $w_k = w_{k-1} \wedge pre(s, w_{k-1})$;
    k = k +1;
  until $w_k \equiv w_{k-1}$;
return $w_k$;
```

Precondition to b through while (B) s;

```
$w_0 = B$; $k = 1$;
repeat
    $w_k = w_{k-1} \wedge pre(s, w_{k-1})$;
    $k = k + 1$;
  until $w_k \equiv w_{k-1}$;
return $w_k$;
```

Example:

while (x<y && x>0) x = x+1;

L2 computation.

$w_0 \equiv x<y \wedge x>0$; $k = 1$;

$w_1 \equiv x<y \wedge x>0 \wedge pre(x=x+1, x<y \wedge x>0)$

$\equiv x<y \wedge x>0 \wedge x+1<y \wedge x+1>0 \equiv x>0 \wedge x+1<y$

$w_2 \equiv x+1<y \wedge x>0 \wedge pre(x=x+1, x+1<y \wedge x>0)$

$\equiv x+1<y \wedge x>0 \wedge x+2<y \wedge x+1>0 \equiv x>0 \wedge x+2<y$

non-terminating for algorithms and protocols!

Precondition to b through while (B) s;

```
$w_0$ = B; k = 1;
repeat
    $w_k = w_{k-1} \wedge pre(s, w_{k-1})$;
    k = k +1;
  until $w_k \equiv w_{k-1}$;
return $w_k$;
```

Example:

while (x>y && x>0) x = x+1;

L2 computation.

$w_0 \equiv x>y \wedge x>0$; k = 1;

$w_1 \equiv x>y \wedge x>0 \wedge pre(x=x+1, x>y \wedge x>0)$

$\equiv x>y \wedge x>0 \wedge x+1>y \wedge x+1>0 \equiv x>y \wedge x>0$

terminating for algorithms and protocols!

Precondition to b through while (B) s;

```
$w_0 = \neg B \wedge b$; $k = 1$;
repeat
    $w_k = w_{k-1} \vee (B \wedge pre(s, w_{k-1}))$;
    $k = k + 1$;
  until $w_k \equiv w_{k-1}$;
return $w_k$;
```

Example: $b \equiv x{==}2 \wedge y{==}3$

while (x>y && x>0) x = x+1;

L1 computation.

$w_0 \equiv (x{<=}y \vee x{<=}0) \wedge x{==}2 \wedge y{==}3 \equiv x{==}2 \wedge y{==}3;$

$w_1 \equiv (x{==}2 \wedge y{==}3) \vee (x>y \wedge x>0 \wedge pre(x=x+1, x{==}2 \wedge y{==}3));$

$\equiv (x{==}2 \wedge y{==}3) \vee (x>y \wedge x>0 \wedge x{==}1 \wedge y{==}3);$

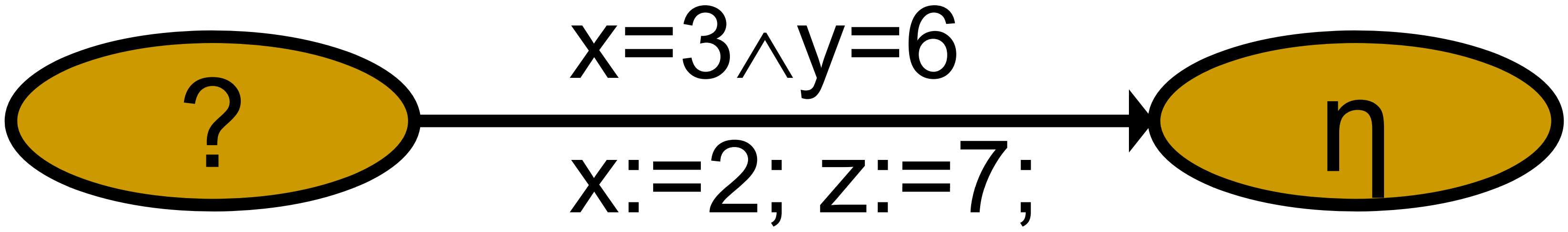
$\equiv (x{==}2 \wedge y{==}3) \vee \text{false}$

$\equiv x{==}2 \wedge y{==}3$

Symbolic weakest precondition

Assume program with rules

- $x=3 \wedge y=6$ → $x:=2; z:=7;$



?

$x=3 \wedge y=6$

$x:=2; z:=7;$

η

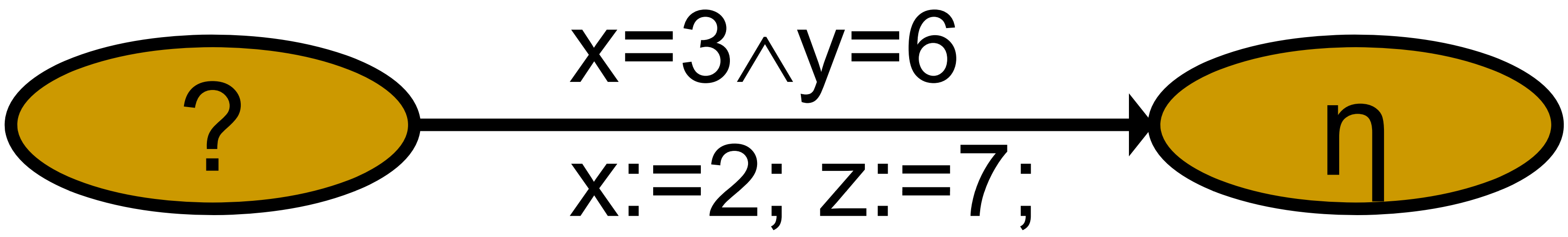
- x, y, z are discrete variables with range declarations

What is the weakest precondition of η *for those states before the transitions ?*

Symbolic weakest precondition

Assume program with rules

- r: $x=3 \wedge y=6 \rightarrow x:=2;\ z:=7;$



What is the weakest precondition of η *for those states before the transitions ?*

$$\mathit{pre}(r, \eta) \stackrel{\text{def}}{=} x=3 \wedge y=6 \wedge \exists x \exists z (x=2 \wedge z=7 \wedge \eta)$$

Symbolic safety analysis - with Kripke structures as programs

Assume program with rules $r_1, r_2, \ldots, r_n$

What charcterizes all states that can reach $\neg\eta$?

```
lfp (φ, ψ) /* for ∃φUψ*/ {
    Z' := false; Z:= ψ;
    while (Z ≠ Z') {
      Z' := Z;
      Z := Z ∨ (φ∧∨_{i=n} pred(r_i, Z));
    }
    return (Z);
}
```

$I \wedge \mathrm{lfp}(\mathrm{true}, \neg\eta) \neq \mathrm{false}$

risk predicate

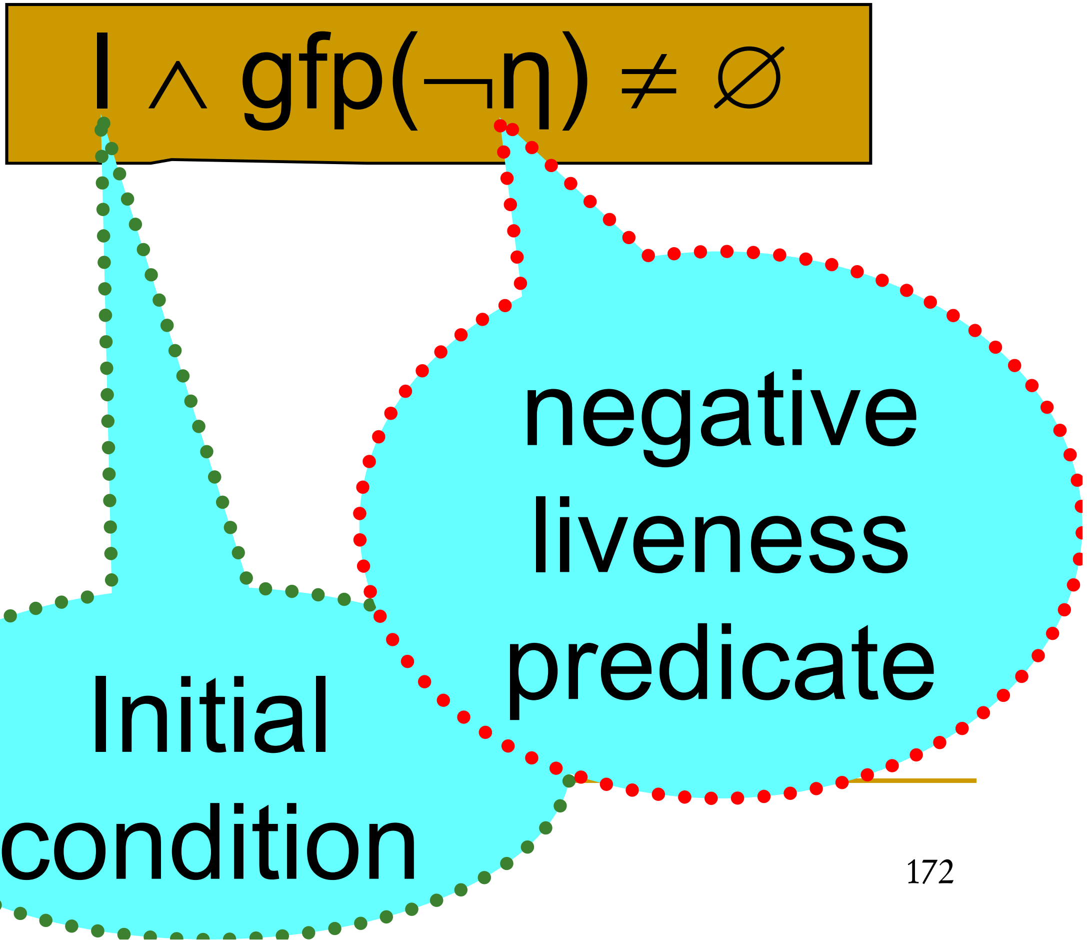
Initial condition

Symbolic liveness analysis - with Kripke structures as programs

Assume program with rules $r_1, r_2, \ldots, r_n$

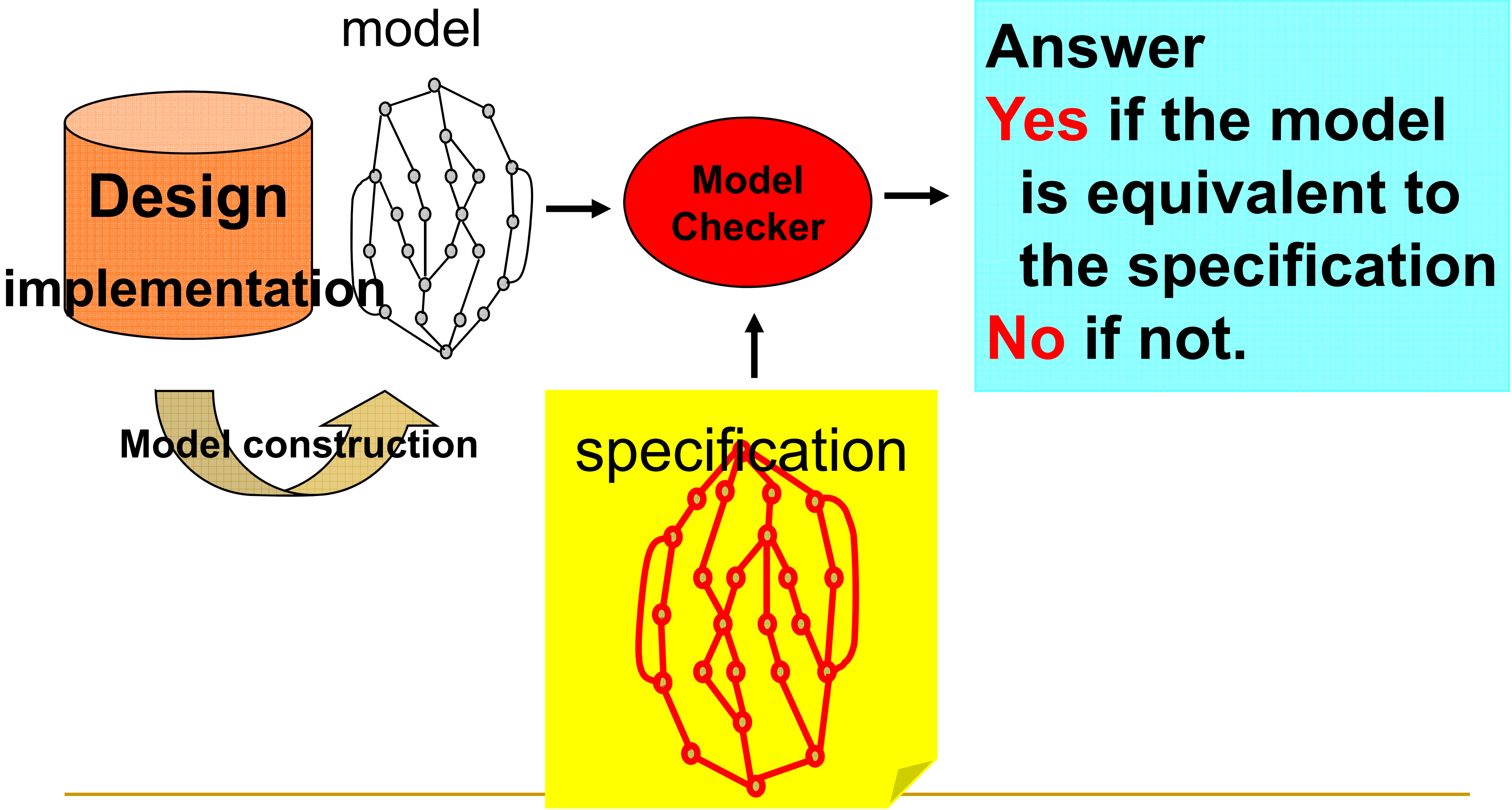
What is the charcterization of all states that may not reach η*?*

```
gfp ($\varphi$) /* for $\exists\Box\varphi$ */{
    Z' := false; Z:= $\varphi$;
    while ($Z \neq Z'$) {
      Z' := Z;
      Z:= $\varphi \wedge \vee_{i=n}$ pred($r_i$, Z);
    }
    return (Z);
}
```



$I \wedge gfp(\neg\eta) \neq \emptyset$

Bisimulation Framework



model

Design implementation

Model construction

Model Checker

specification

Answer

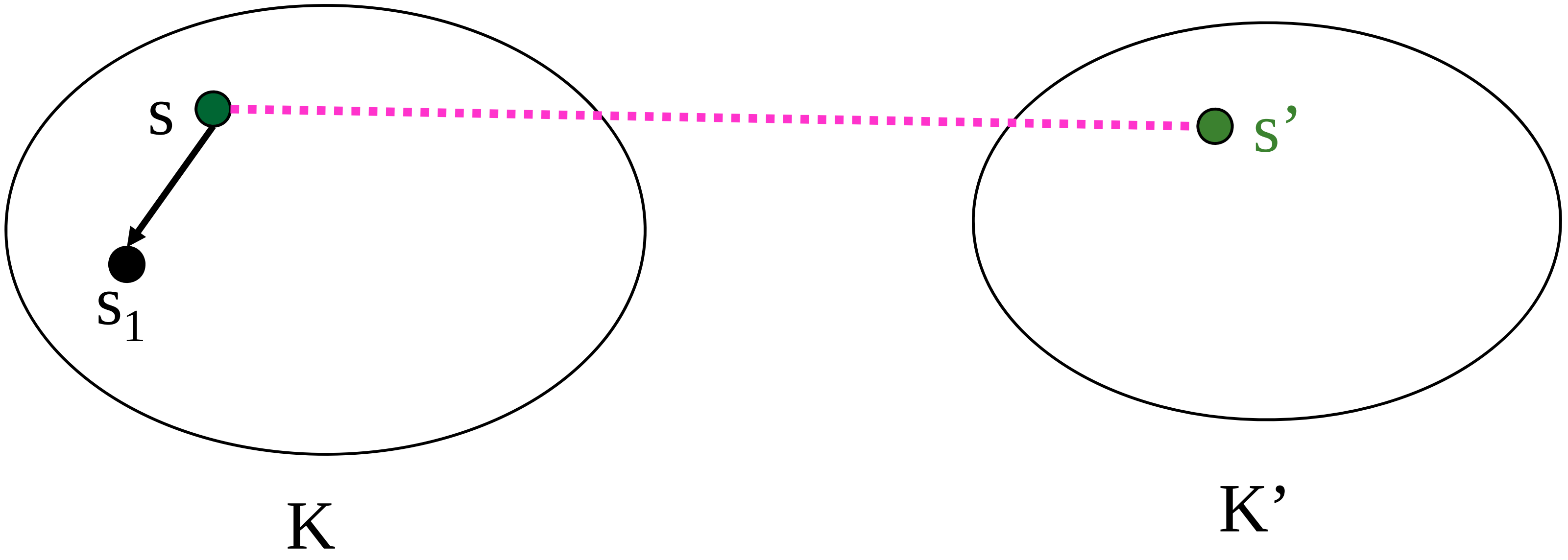
Yes if the model is equivalent to the specification

No if not.

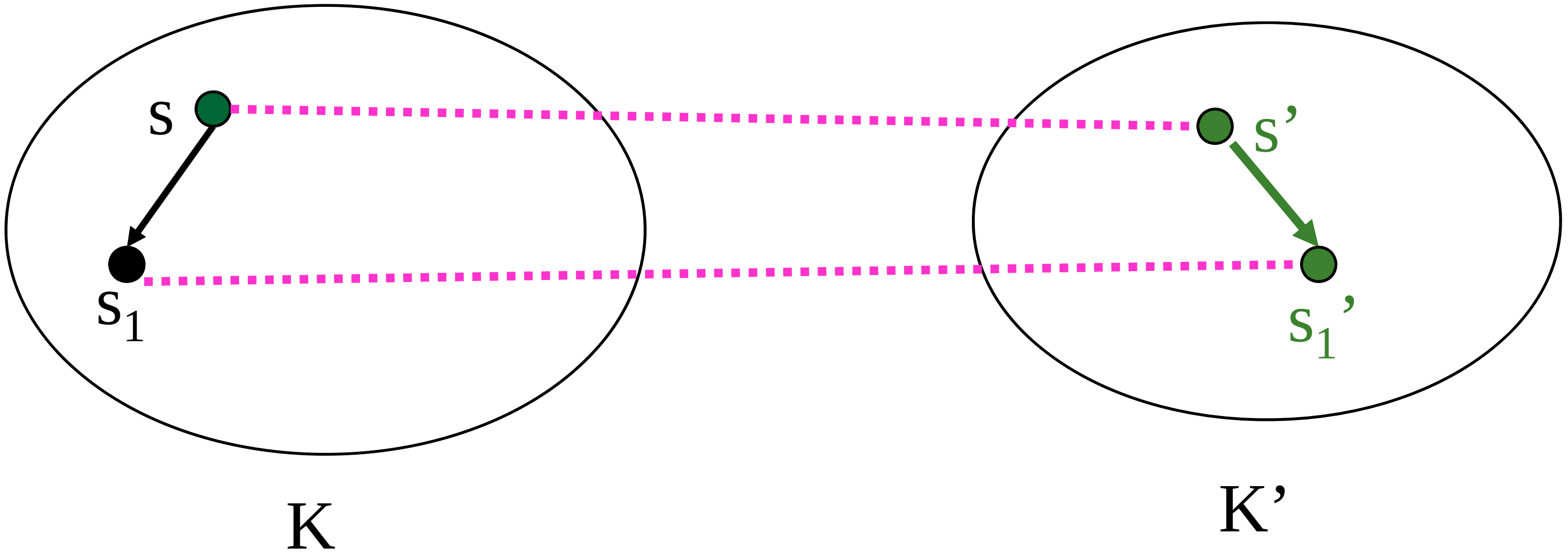
Bisimulation-checking

- K = (S, S_0, R, AP, L)
  K'= (S', S_0', R', AP, L')
- Note K and K' use the same set of atomic propositions AP.
- B∈S×S' is a bisimulation relation between K and K' iff for every B(s, s'):
  - L(s) = L'(s') (BSIM 1)
  - If R(s, s_1), then there exists s_1' such that R'(s', s_1') and B(s_1, s_1'). (BISIM 2)
  - If R(s', s_2'), then there exists s_2 such that R(s, s_2) and B(s_2, s_2'). (BISIM 3)

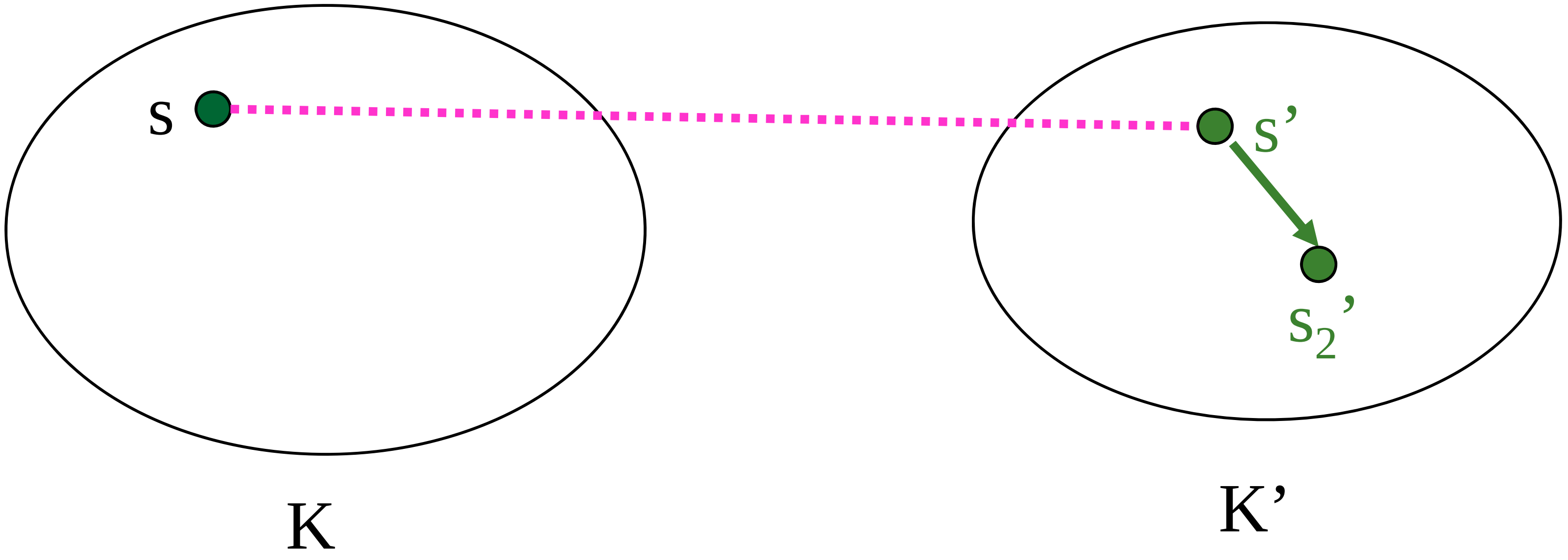
Bisimulations



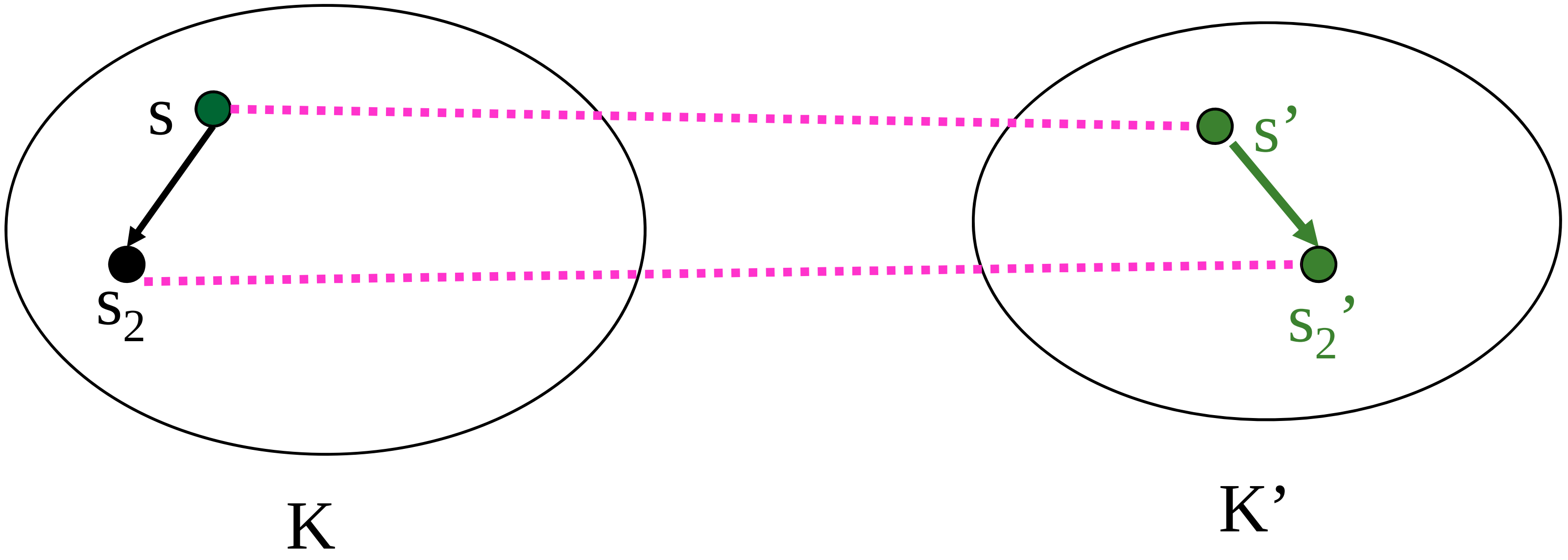
Bisimulations



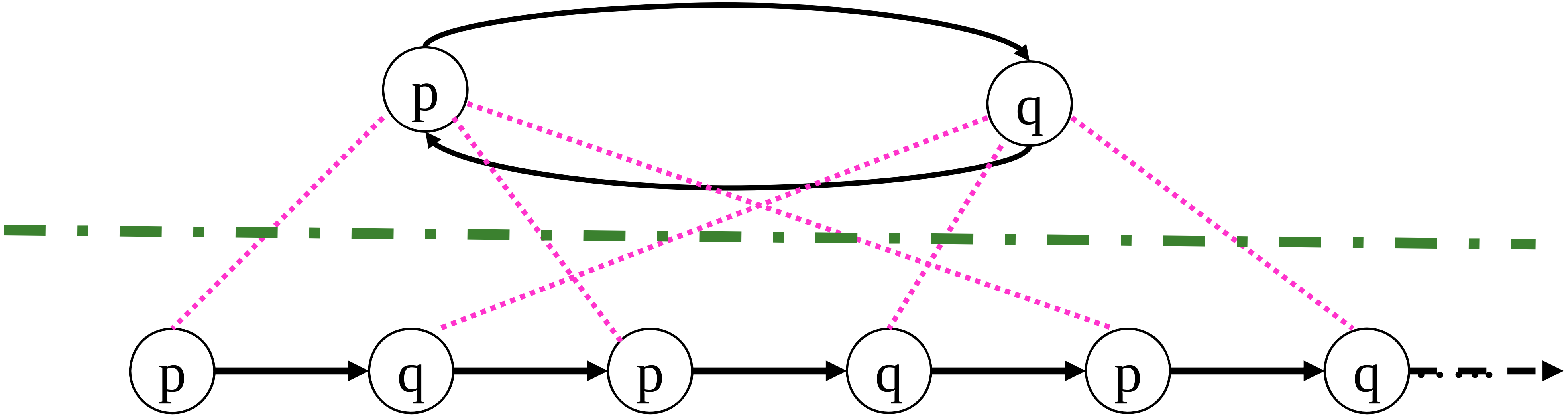
Bisimulations



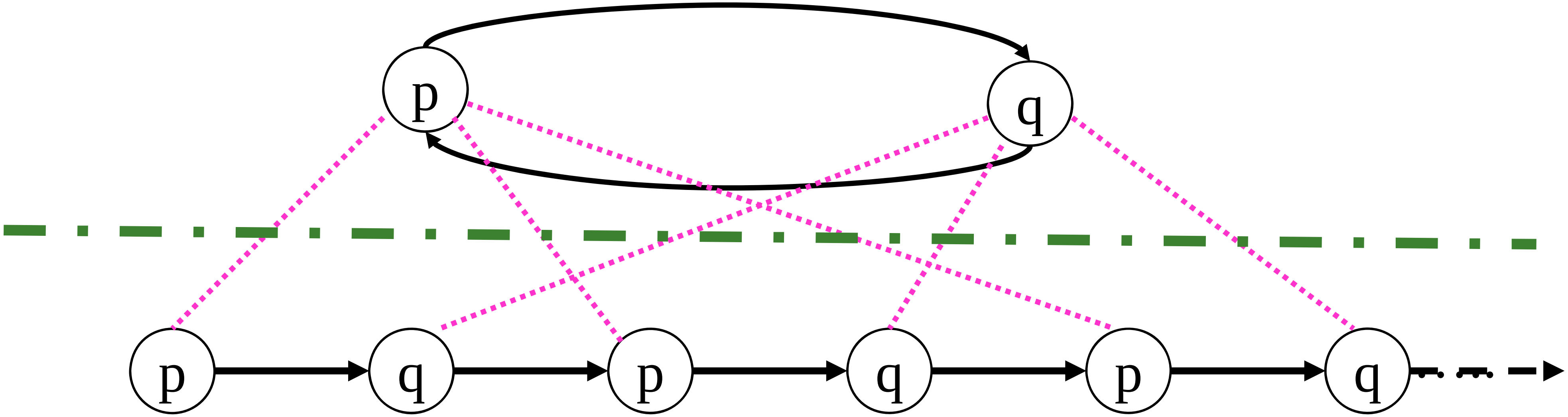
Bisimulations



Examples

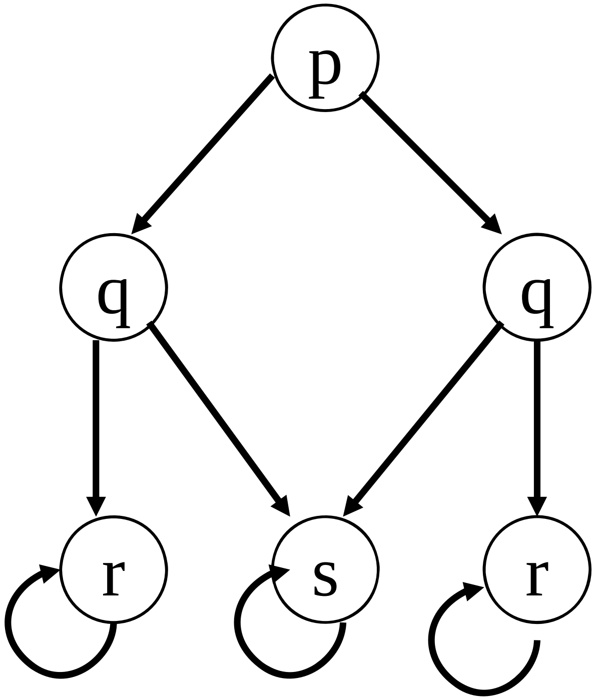


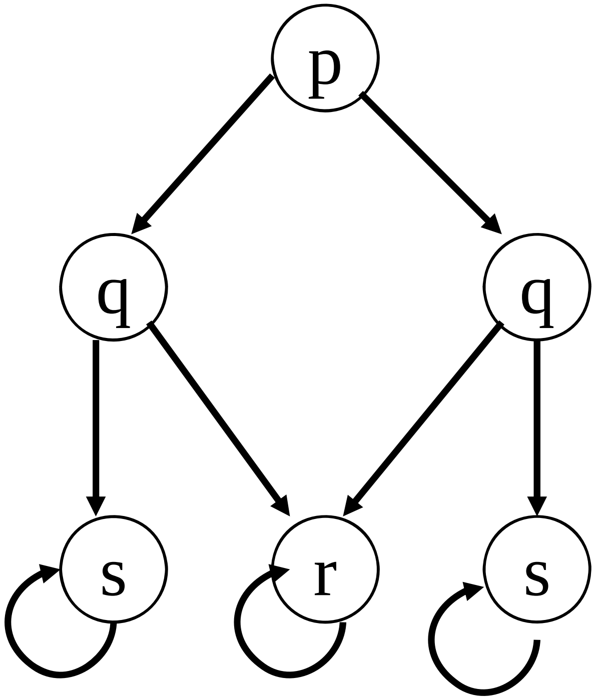
Examples



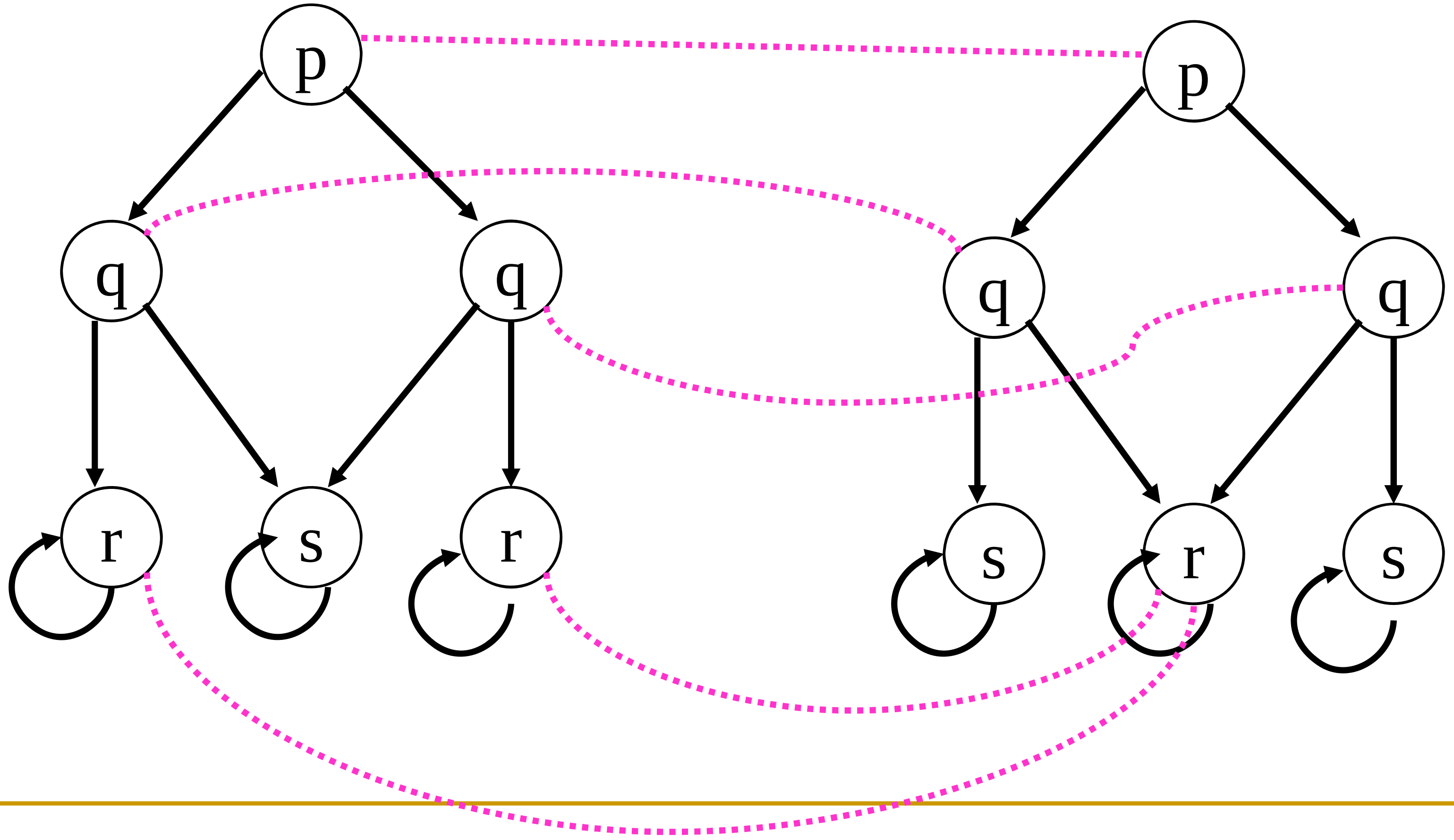
Unwinding preserves bisimulation

Examples

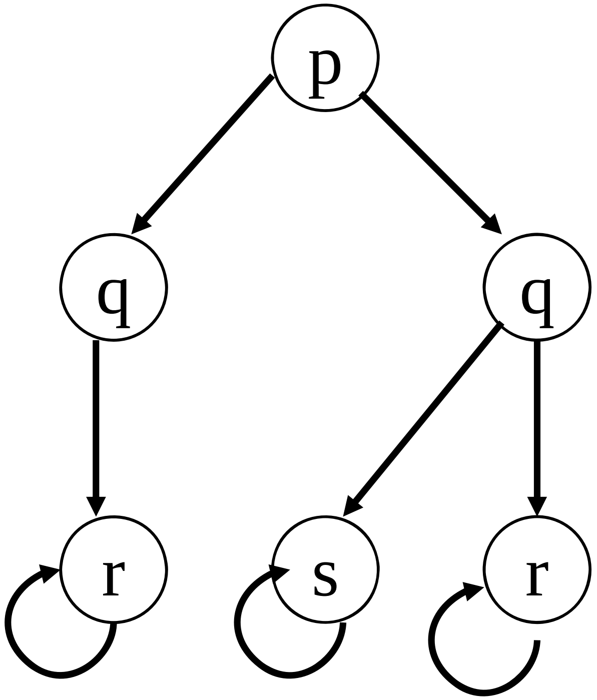


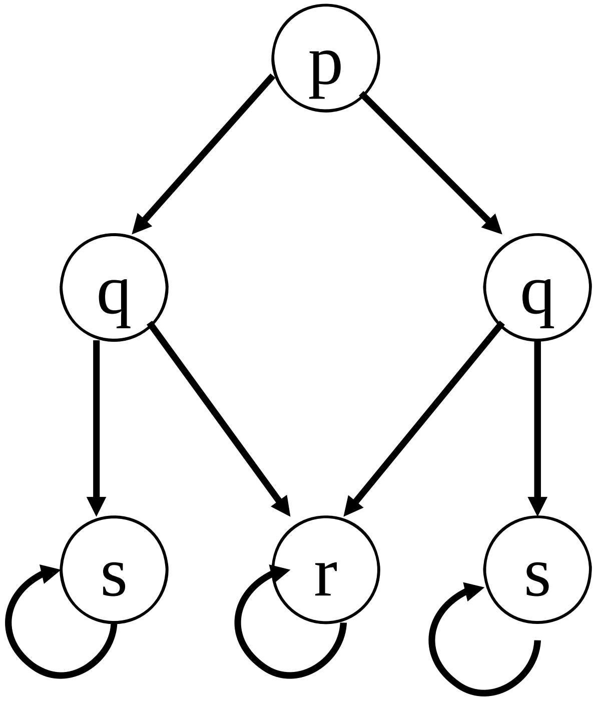


Examples

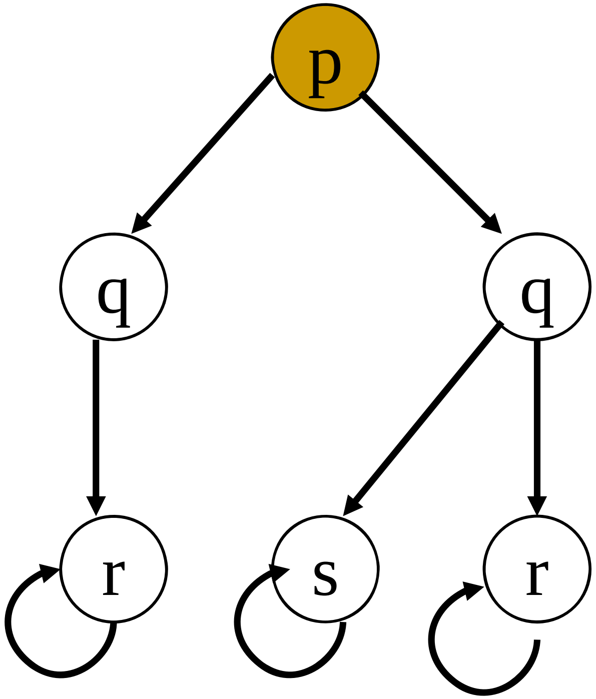


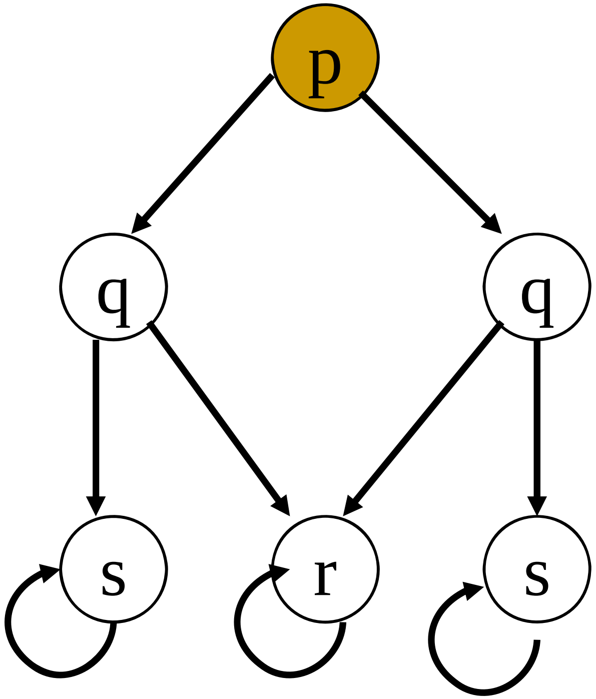
Examples



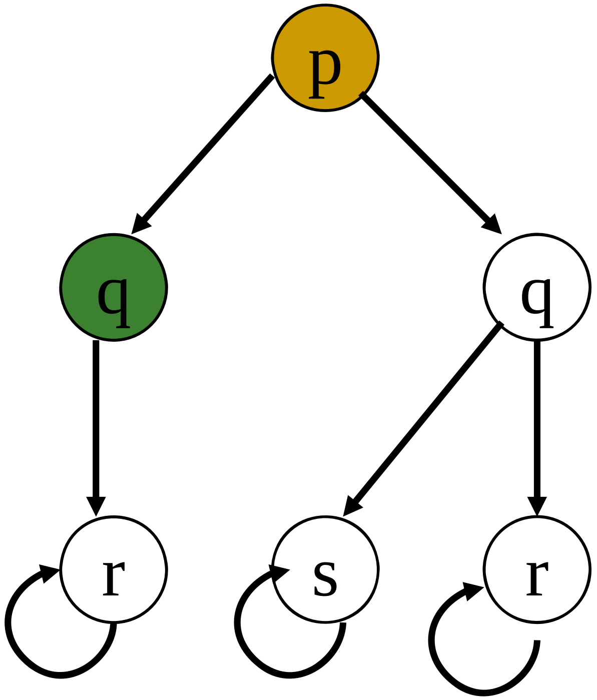


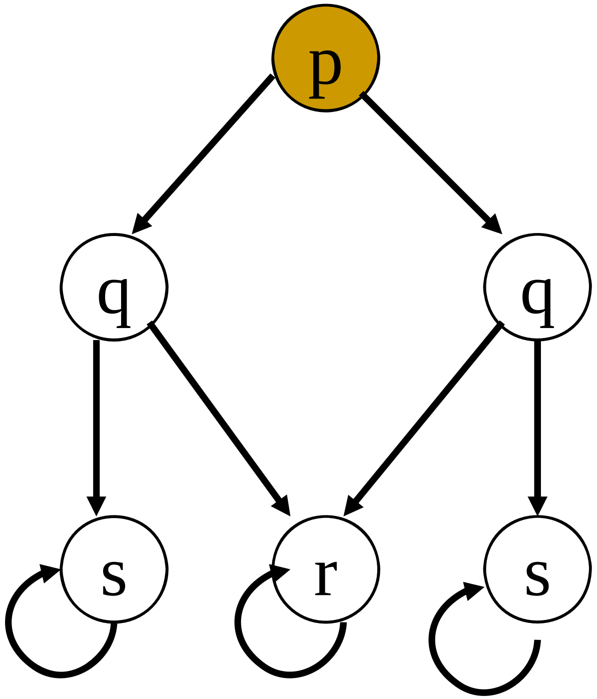
Examples



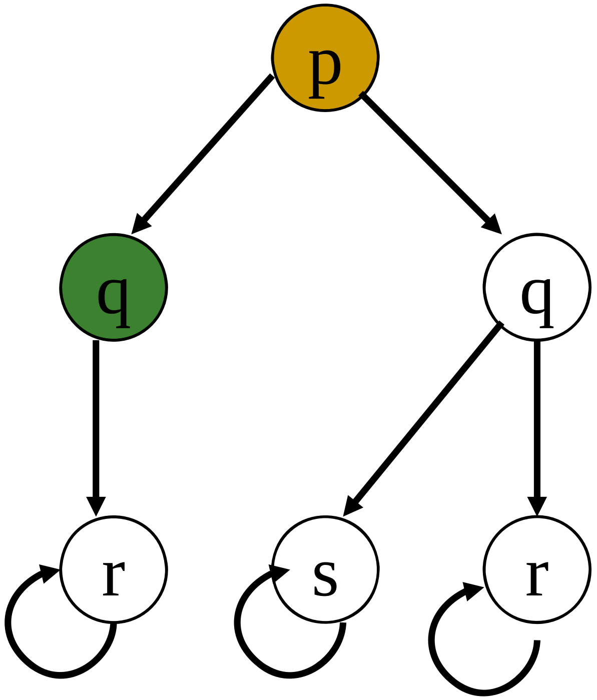


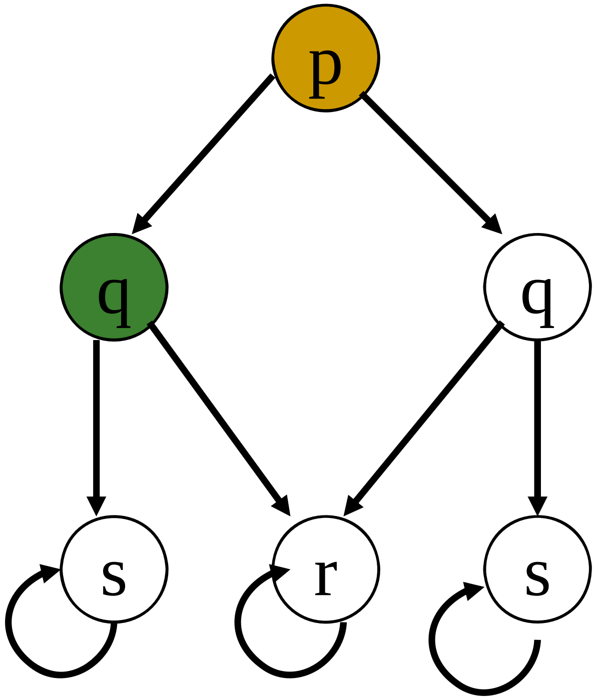
Examples



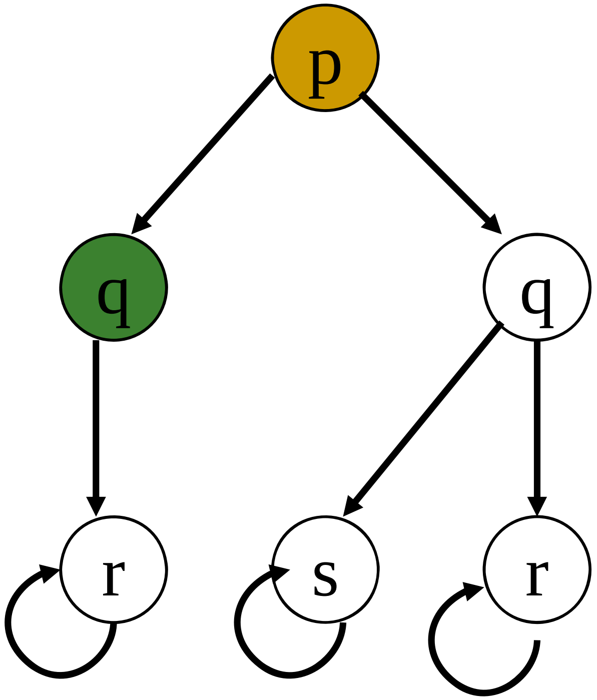


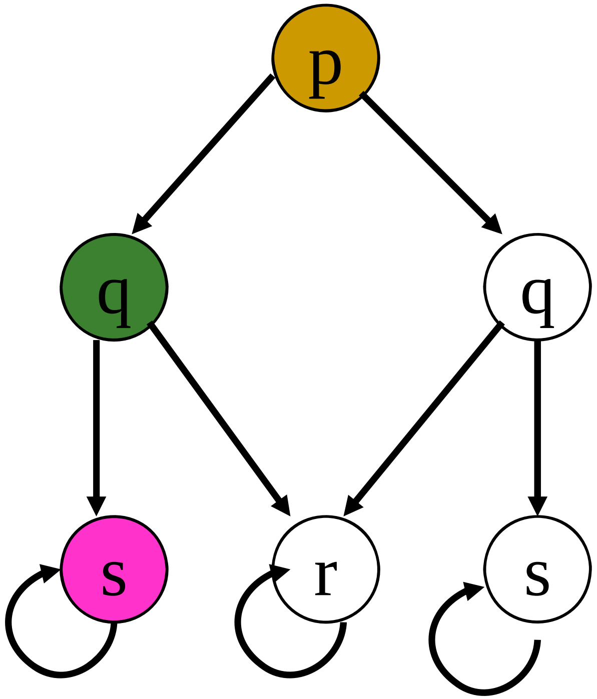
Examples





Examples





Bisimulations

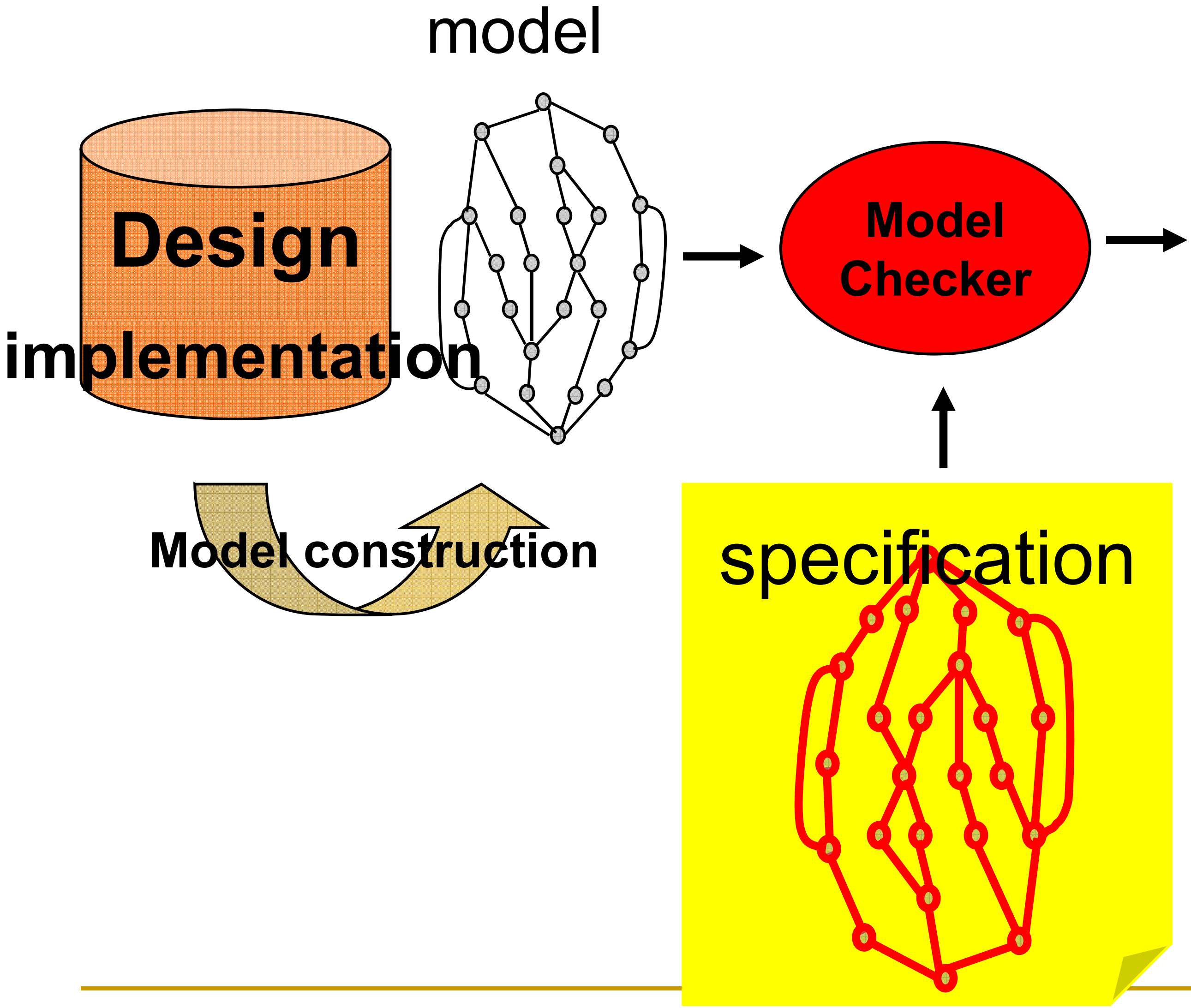
- K = (S, S_0, R, AP, L)
- K'= (S', S_0', R', AP, L')
- K and K' are bisimilar (bisimulation equivalent) iff there exists a bisimulation relation $\mathcal{B} \subseteq S \times S'$ between K and K' such that:
  - For each s_0 in S_0 there exists s_0' in S_0' such that $\mathcal{B}(s_0 , s_0')$.
  - For each s_0' in S_0' there exists s_0 in S_0 such that $\mathcal{B}(s_0 , s_0')$.

The Preservation Property.

- K = (S, S_0, R, AP, L)

  K’= (S’, S_0’, R’, AP, L’)
- $\mathbb{B} \subseteq S \times S'$, a bisimulation.
- Suppose $\mathbb{B}(s, s')$.
- FACT: For any CTL formula ψ (over AP), K,s$\models\psi$ iff K’,s’$\models\psi$.
- If K’ is smaller than K this is worth something.

Simulation Framework



model

Design implementation

Model construction

Model Checker

specification

Answer

Yes if the model satisfies the specification

No if not.

Simulation-checking

- $K = (S, S_0, R, AP, L)$
  $K' = (S', S_0', R', AP, L')$
- Note K and K' use the same set of atomic propositions AP.
- $B \in S \times S'$ is a simulation relation between K and K' iff for every $B(s, s')$:
  - $L(s) = L'(s')$ (BSIM 1)
  - If $R(s, s_1)$, then there exists s_1' such that $R'(s', s_1')$ and $B(s_1, s_1')$. (BISIM 2)

Simulations

- $K = (S, S_0, R, AP, L)$
- $K' = (S', S_0', R', AP, L')$
- K is simulated by (implements or refines) K' iff there exists a simulation relation $\mathcal{B} \subseteq S \times S'$ between K and K' such that for each s_0 in S_0 there exists s_0' in S_0' such that $\mathcal{B}(s_0 , s_0')$.

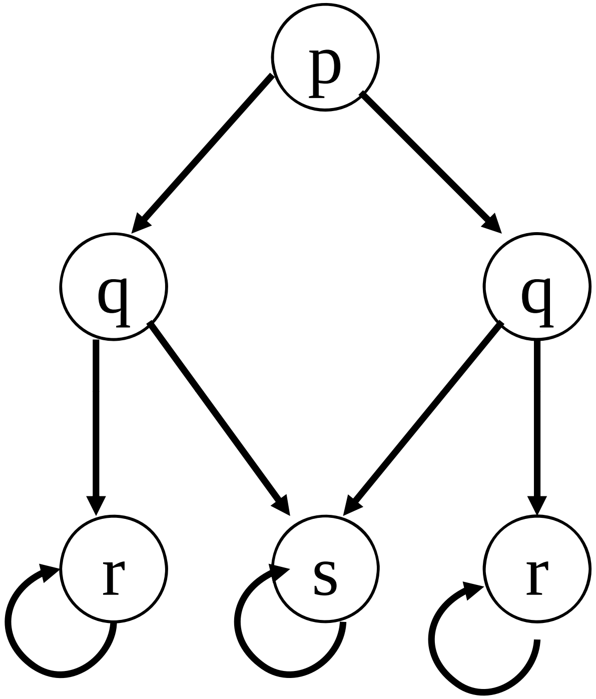
Bisimulation Quotients

- $K = (S, S_0, R, AP, L)$
- There is a maximal simulation $\mathrm{B} \subseteq S \times S$.
  - Let R be this bisimulation.
  - $[s] = \{s' \mid s \mathrm{R} s'\}$.
- R can be computed “easily”.
- K’ = K / R is the bisimulation quotient of K.

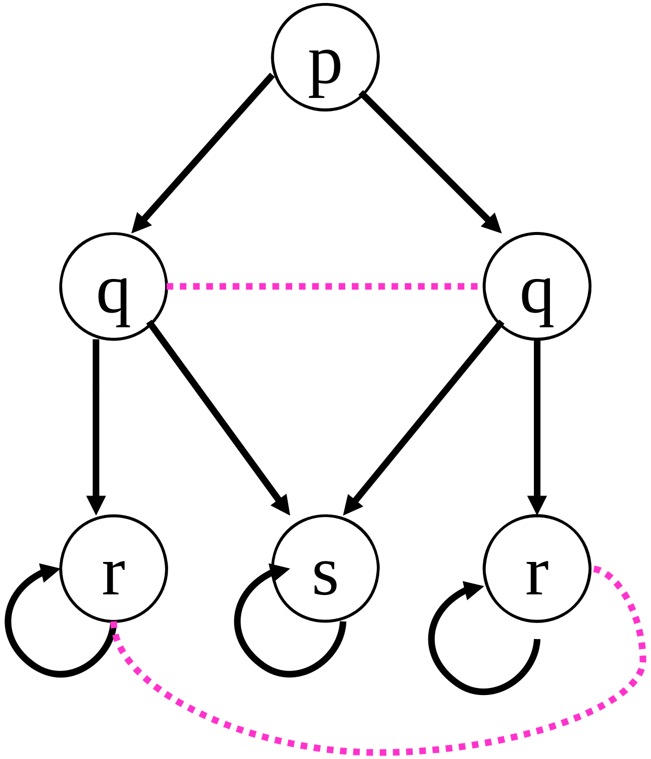
Bisimulation Quotient

- $K = (S, S_0, R, AP, L)$
- $[s] = \{s' \mid s \mathcal{R} s'\}$.
- $K' = K / \mathcal{R} = (S', S'_0, R', AP, L')$.
  - $S' = \{[s] \mid s \; 2 \; S\}$
  - $S'_0 = \{[s_0] \mid s_0 \; 2 \; S_0\}$
  - $R' = \{([s], [s']) \mid R(s_1, s_1') , s_1 \in [s], s_1' \in [s']\}$
  - $L'([s]) = L(s)$.

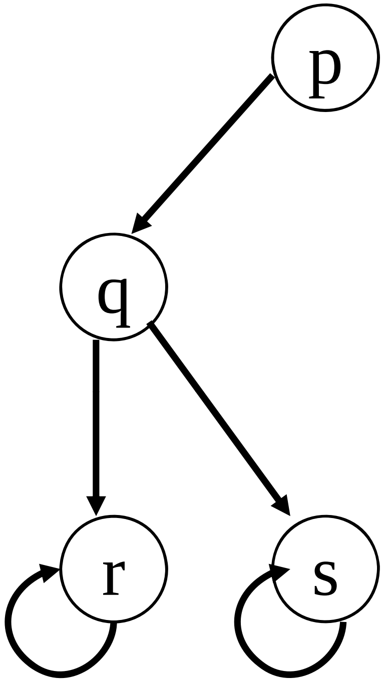
Examples



Examples



Examples



Facts About a (Bi)Simulation

- The empty set is always a (bi)simulation
- If R, R' are (bi)simulations, so is R U R'
- Hence, there always exists a *maximal* (bi)simulation:
  - Checking if DB_1=DB_2: compute the maximal bisimulation R, then test (root(DB_1),root(DB_2)) in R

Kripke structure
- simulation-checking

```
/* Given model A = (S, $S_0$, R, L), spec. A'=(S', $S'_0$, R', L') */
Simulation-checking(A,A') /* using greatest fixpoint algorithm */ {
    Let B={(s,s') | s$\in$S, s'$\in$S', L(s)=L'(s')} ;
    repeat {
        B = B - {(s,s') | (s,s')$\in$B, $\exists$(s,t)$\in$R$\forall$(s',t')$\in$R'((t,t')$\notin$B)};
    } until no more changes to B.
    if there is an $s_0 \in S_0$ with $\forall s'_0 \in S'_0((s_0,s'_0) \notin B)$,
        return 'no simulation,'
        else return 'simulation exists.'
}
```

The procedure terminates since B is finite in the Kripke structure.

Kripke structure - bisimulation-checking

```
/* Given model $A = (S, S_0, R, L)$, spec. $A'=(S', S'_0, R', L')$ */
Bisimulation-checking(A,A’) /* using greatest fixpoint algorithm */ {
    Let $B=\{(s,s') \mid s\in S, s'\in S', L(s)=L'(s')\}$ ;
    repeat {
        $B = B - \{(s,s') \mid (s,s')\in B, \exists(s,t)\in R\forall(s',t')\in R'((t,t')\notin B)\}$;
        $B = B - \{(s,s') \mid (s,s')\in B, \exists(s',t')\in R'\forall(s,t)\in R((t,t')\notin B)\}$;
    } until no more changes to B.
    if there is an $s_0\in S_0$ with $\forall s'_0\in S'_0((s_0,s'_0)\notin B)$,
        return ‘no simulation,’
    if there is an $s'_0\in S'_0$ with $\forall s_0\in S_0((s_0,s'_0)\notin B)$,
        return ‘no simulation,’
    else return ‘simulation exists.’
}
```

(Bi)Simulation - complexities

- Bisimulation: O((m+n)log(m+n))
- Simulation: O(m n)
- In contrast, finding a graph homeomorphism is NP-complete.

Symbolic simulation-checking

- Encode the states with variables
  - $x_0,x_1,\ldots,x_n$ (for the model) and
  - $y_0,y_1, \ldots, y_m$. (for the spec.)

  Usually there are shared variables

  between $\{x_0,x_1,\ldots,x_n\}$ and $\{y_0,y_1, \ldots, y_m\}$.

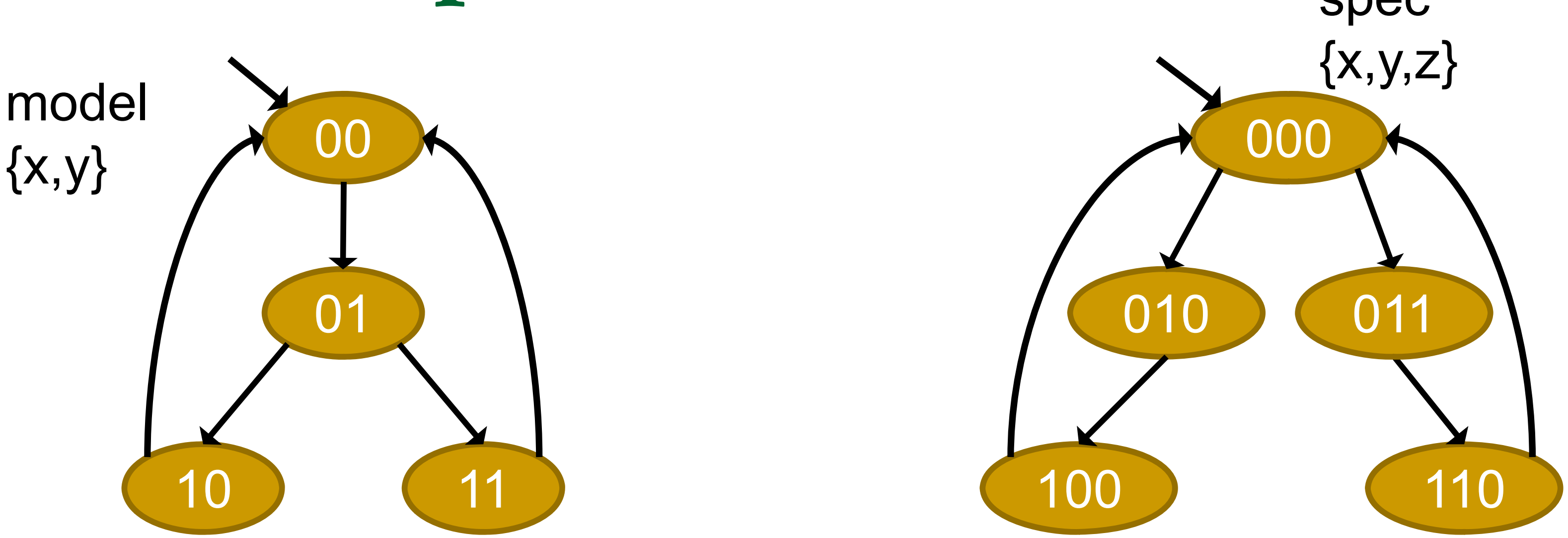
  L(s)=L'(s') means that the shared variables are of the same values.
- the state sets as proposition formulas:
  - $s(x_0,x_1,\ldots,x_n)$ & $s(y_0,y_1,\ldots,y_m)$
- the initial state set as
  - $i(x_0,x_1,\ldots,x_n)$ & $i'(y_0,y_1,\ldots,y_m)$
- the transition set as
  - $R(x_0,x_1,\ldots,x_n,x'_0,x'_1,\ldots,x_n)$ & $R'(y_0,y_1,\ldots,y_n,y'_0,y'_1,\ldots,y'_n)$

Symbolic simulation-checking

```
$B_0 = \bigwedge_{L(x0,x1,\ldots,xn) = L(y0,y1,\ldots,ym)} s(x_0,x_1,\ldots,x_n) \wedge s(y_0,y_1,\ldots,y_m)$;
for ($k = 1$, $B_1$= false; $B_k \neq B_{k-1}$; $k=k+1$)
    $B_k = B_{k-1} \wedge \neg\exists x'_0 \exists x'_1 \ldots \exists x'_n$ (
                $R(x_0,x_1,\ldots,x_n, x'_0,x'_1,\ldots,x'_n)$
              $\wedge \neg \exists y'_0 \exists y'_1 \ldots \exists y'_m$ (
                    $R'(y_0,y_1,\ldots,y_m, y'_0,y'_1,\ldots,y'_m) \wedge (B_{k-1} \uparrow)$
           )    ) ;
if ($i(x_0,x_1,\ldots,x_n) \neq \exists y_0 \exists y_1 \ldots \exists y_m (B_k)$),
    return 'no simulation';
else return 'a simulation exists';
```

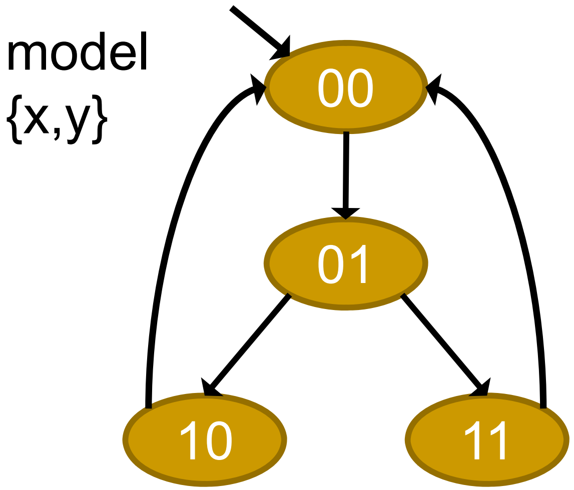
change all umprimed variable in B_{k-1} to primed.

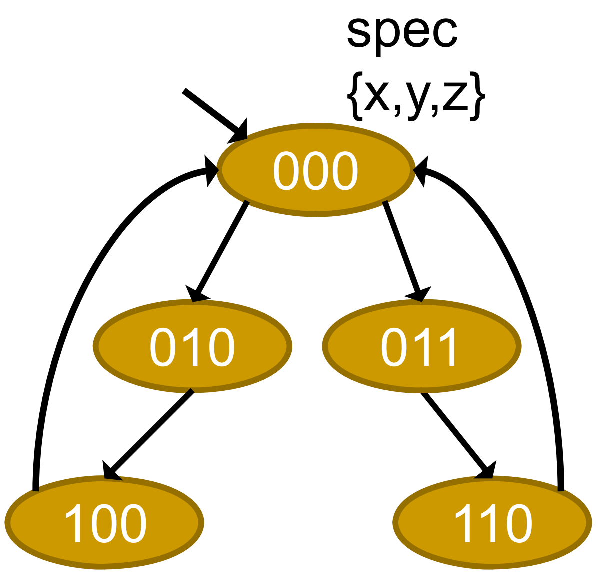
Symbolic simulation-checking - an example



- $s(x,y) \equiv true$, $s'(x,y,z) \equiv \neg z \vee (\neg x \wedge y \wedge z)$
- $i(x,y) \equiv \neg x \wedge \neg y$, $i'(x,y,z) \equiv \neg x \wedge \neg y \wedge \neg z$
- $R(x,y,x',y') \equiv \ldots\ldots$, $R'(x,y,z,x',y',z') \equiv \ldots\ldots$

Symbolic simulation-checking - an example





- $R(x,y,x',y') \equiv (\neg x \wedge \neg y \wedge \neg x' \wedge y') \vee (\neg x \wedge y \wedge x' \wedge \neg y')$
  $\vee (\neg x \wedge y \wedge x' \wedge y') \vee (x \wedge \neg y \wedge \neg x' \wedge \neg y') \vee (x \wedge y \wedge \neg x' \wedge \neg y')$
- $R'(x,y,z,x',y',z') \equiv (\neg x \wedge \neg y \wedge \neg z \wedge \neg x' \wedge y')$
  $\vee (\neg x \wedge y \wedge \neg z \wedge x' \wedge \neg y' \wedge \neg z') \vee (\neg x \wedge y \wedge z \wedge x' \wedge y' \wedge \neg z')$
  $\vee (x \wedge \neg y \wedge \neg z \wedge \neg x' \wedge \neg y' \wedge \neg z') \vee (x \wedge y \wedge \neg z \wedge \neg x' \wedge \neg y' \wedge \neg z')$

Symbolic simulation-checking - an example

$$B_0 = s(x,y) \wedge s'(x,y,z) = \neg z \vee (\neg x \wedge y \wedge z)$$

$$B_1 = (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg\, \exists x' \exists y' \,($$

$$((\neg x \wedge \neg y \wedge \neg x' \wedge y') \vee (\neg x \wedge y \wedge x' \wedge \neg y')$$

$$\vee\, (\neg x \wedge y \wedge x' \wedge y') \vee (x \wedge \neg y \wedge \neg x' \wedge \neg y') \vee (x \wedge y \wedge \neg x' \wedge \neg y')$$

$$)$$

$$\wedge\, \neg\, \exists x' \exists y' \exists z' \,($$

$$(\quad (\neg x \wedge \neg y \wedge \neg z \wedge \neg x' \wedge y')$$

$$\vee\, (\neg x \wedge y \wedge \neg z \wedge x' \wedge \neg y' \wedge \neg z') \vee (\neg x \wedge y \wedge z \wedge x' \wedge y' \wedge \neg z')$$

$$\vee\, (x \wedge \neg y \wedge \neg z \wedge \neg x' \wedge \neg y' \wedge \neg z') \vee (x \wedge y \wedge \neg z \wedge \neg x' \wedge \neg y' \wedge \neg z')$$

$$) \wedge (\neg z' \vee (\neg x' \wedge y' \wedge z'))\ ))$$

$$= (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg\, \exists x' \exists y' \,(((\neg x \wedge \neg y \wedge z \wedge \neg x' \wedge y') \vee (\neg x \wedge y \wedge x' \wedge y')$$

$$\vee (x \wedge \neg y \wedge z \wedge \neg x' \wedge \neg y') \vee (x \wedge y \wedge z \wedge \neg x' \wedge \neg y')))$$

$$= (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg((\neg x \wedge \neg y \wedge z) \vee (\neg x \wedge y) \vee (x \wedge \neg y \wedge z) \vee (x \wedge y \wedge z))$$

Symbolic simulation-checking - an example

$B_1 = (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg((\neg x \wedge \neg y \wedge z) \vee (\neg x \wedge y) \vee (x \wedge \neg y \wedge z) \vee (x \wedge y \wedge z))$

$= (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg((\neg x \wedge \neg y \wedge z) \vee (\neg x \wedge y) \vee (x \wedge \neg y \wedge z) \vee (x \wedge y \wedge z))$

$= (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg(z \vee (\neg x \wedge y \wedge \neg z))$

$= (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg(z) \wedge \neg(\neg x \wedge y \wedge \neg z)$

$= (\neg z \vee (\neg x \wedge y \wedge z)) \wedge \neg(z) \wedge \neg(\neg x \wedge y \wedge \neg z)$

$= (\neg x \wedge \neg y \wedge \neg z) \vee (x \wedge \neg y \wedge \neg z) \vee (x \wedge y \wedge \neg z)$

Symbolic simulation-checking - an example

$B_2 = ((\lnot x\land\lnot y\land\lnot z)\lor (x\land\lnot y\land\lnot z)\lor(x\land y \land \lnot z)) \land \lnot \exists x'\exists y' ($

$((\lnot x\land\lnot y\land\lnot x'\land y') \lor(\lnot x\land y\land x'\land\lnot y')$

$\lor (\lnot x\land y\land x'\land y') \lor(x\land\lnot y\land\lnot x'\land\lnot y') \lor(x\land y\land\lnot x'\land\lnot y')$

$)$

$\land \lnot \exists x'\exists y'\exists z' ($

$((\lnot x\land\lnot y\land\lnot z\land\lnot x'\land y')$

$\lor (\lnot x\land y \land\lnot z\land x'\land\lnot y'\land\lnot z') \lor(\lnot x\land y\land z\land x'\land y'\land\lnot z')$

$\lor (x\land\lnot y\land\lnot z \land\lnot x'\land\lnot y' \land\lnot z') \lor(x\land y \land\lnot z\land\lnot x'\land\lnot y'\land\lnot z')$

$) \land ((\lnot x'\land\lnot y'\land\lnot z')\lor (x'\land\lnot y'\land\lnot z')\lor(x'\land y' \land \lnot z'))))$

$= ((\lnot x\land\lnot y\land\lnot z)\lor (x\land\lnot y\land\lnot z)\lor(x\land y \land \lnot z)) \land \lnot \exists x'\exists y' ($

$((\lnot x\land\lnot y\land\lnot x'\land y') \lor(x\land\lnot y\land z\land\lnot x'\land\lnot y') \lor(x\land y\land z\land\lnot x'\land\lnot y')))$

$= ((\lnot x\land\lnot y\land\lnot z)\lor(x\land\lnot y\land\lnot z)\lor(x\land y \land\lnot z))\land\lnot((\lnot x\land\lnot y)\lor(x\land\lnot y\land z)\lor(x\land y\land z)))$

Symbolic simulation-checking - an example

B_2

$= ((\neg x \wedge \neg y \wedge \neg z) \vee (x \wedge \neg y \wedge \neg z) \vee (x \wedge y \wedge \neg z)) \wedge \neg((\neg x \wedge \neg y) \vee (x \wedge \neg y \wedge z) \vee (x \wedge y \wedge z)))$

$= (x \wedge \neg y \wedge \neg z) \vee (x \wedge y \wedge \neg z)$

Here, the initial statepair has been elimianted.